16

H8/38537 Group

Renesas 16-Bit Single-Chip Microcomputer H8 Family / H8/300H Super Low Power Series

H8/38532

H8/38533

H8/38534

H8/38535

H8/38536

H8/38537

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General Precautions in the Handling of MPU/MCU Products

The following usage notes are applicable to all MPU/MCU products from Renesas. For detailed usage notes on the products covered by this manual, refer to the relevant sections of the manual. If the descriptions under General Precautions in the Handling of MPU/MCU Products and in the body of the manual differ from each other, the description in the body of the manual takes precedence.

1. Handling of Unused Pins

Handle unused pins in accord with the directions given under Handling of Unused Pins in the manual.

The input pins of CMOS products are generally in the high-impedance state. In operation with an unused pin in the open-circuit state, extra electromagnetic noise is induced in the vicinity of LSI, an associated shoot-through current flows internally, and malfunctions occur due to the false recognition of the pin state as an input signal become possible. Unused pins should be handled as described under Handling of Unused Pins in the manual.

2. Processing at Power-on

The state of the product is undefined at the moment when power is supplied.

- The states of internal circuits in the LSI are indeterminate and the states of register settings and pins are undefined at the moment when power is supplied.
 - In a finished product where the reset signal is applied to the external reset pin, the states of pins are not guaranteed from the moment when power is supplied until the reset process is completed.

In a similar way, the states of pins in a product that is reset by an on-chip power-on reset function are not guaranteed from the moment when power is supplied until the power reaches the level at which resetting has been specified.

3. Prohibition of Access to Reserved Addresses

Access to reserved addresses is prohibited.

 The reserved addresses are provided for the possible future expansion of functions. Do not access these addresses; the correct operation of LSI is not guaranteed if they are accessed.

4. Clock Signals

After applying a reset, only release the reset line after the operating clock signal has become stable. When switching the clock signal during program execution, wait until the target clock signal has stabilized.

- When the clock signal is generated with an external resonator (or from an external oscillator) during a reset, ensure that the reset line is only released after full stabilization of the clock signal. Moreover, when switching to a clock signal produced with an external resonator (or by an external oscillator) while program execution is in progress, wait until the target clock signal is stable.
- Differences between Products

Before changing from one product to another, i.e. to one with a different type number, confirm that the change will not lead to problems.

The characteristics of MPU/MCU in the same group but having different type numbers may differ because of the differences in internal memory capacity and layout pattern. When changing to products of different type numbers, implement a system-evaluation test for each of the products.



How to Use This Manual

1. Objective and Target Users

This manual was written to explain the hardware functions and electrical characteristics of this LSI to the target users, i.e. those who will be using this LSI in the design of application systems. Target users are expected to understand the fundamentals of electrical circuits, logic circuits, and microcomputers.

This manual is organized in the following items: an overview of the product, descriptions of the CPU, system control functions, and peripheral functions, electrical characteristics of the device, and usage notes.

When designing an application system that includes this LSI, take all points to note into account. Points to note are given in their contexts and at the final part of each section, and in the section giving usage notes.

The list of revisions is a summary of major points of revision or addition for earlier versions. It does not cover all revised items. For details on the revised points, see the actual locations in the manual.

The following documents have been prepared for the H8/38537 Group. Before using any of the documents, please visit our web site to verify that you have the most up-to-date available version of the document.

Document Type	Contents	Document Title	Document No.
Data Sheet	Overview of hardware and electrical characteristics	_	_
Hardware Manual	Hardware specifications (pin assignments, memory maps, peripheral specifications, electrical characteristics, and timing charts) and descriptions of operation	H8/38537 Group Hardware Manual	This manual
Software Manual	Detailed descriptions of the CPU and instruction set	H8/300H Series Software Manual	REJ09B0213
Application Note	Examples of applications and sample programs	The latest versions are ava	ailable from our web
Renesas Technical Update	Preliminary report on the specifications of a product, document, etc.	_	

2. Description of Numbers and Symbols

Aspects of the notations for register names, bit names, numbers, and symbolic names in this manual are explained below.

(1) Overall notation

In descriptions involving the names of bits and bit fields within this manual, the modules and registers to which the bits belong may be clarified by giving the names in the forms "module name". "register name". "bit name" or "register name". "bit name".

(2) Register notation

The style "register name"_"instance number" is used in cases where there is more than one instance of the same function or similar functions.

[Example] CMCSR 0: Indicates the CMCSR register for the compare-match timer of channel 0.

(3) Number notation

Binary numbers are given as B'nnnn (B' may be omitted if the number is obviously binary), hexadecimal numbers are given as H'nnnn or 0xnnnn, and decimal numbers are given as nnnn.

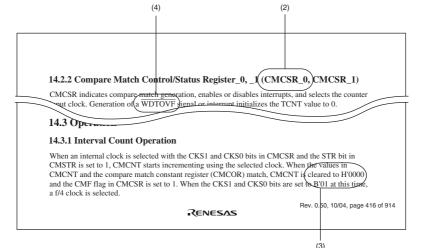
[Examples] Binary: B'11 or 11

Hexadecimal: H'EFA0 or 0xEFA0

Decimal: 1234

(4) Notation for active-low

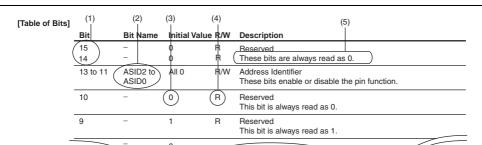
An overbar on the name indicates that a signal or pin is active-low. [Example] $\overline{\text{WDTOVF}}$



Note: The bit names and sentences in the above figure are examples and have nothing to do with the contents of this manual.

3. Description of Registers

Each register description includes a bit chart, illustrating the arrangement of bits, and a table of bits, describing the meanings of the bit settings. The standard format and notation for bit charts and tables are described below.



Note: The bit names and sentences in the above figure are examples, and have nothing to do with the contents of this manual

(1) Bit

Indicates the bit number or numbers.

In the case of a 32-bit register, the bits are arranged in order from 31 to 0. In the case of a 16-bit register, the bits are arranged in order from 15 to 0.

(2) Bit name

Indicates the name of the bit or bit field.

When the number of bits has to be clearly indicated in the field, appropriate notation is included (e.g., ASID[3:01).

A reserved bit is indicated by "-".

Certain kinds of bits, such as those of timer counters, are not assigned bit names. In such cases, the entry under Bit Name is blank.

(3) Initial value

Indicates the value of each bit immediately after a power-on reset, i.e., the initial value.

- 0: The initial value is 0
- 1: The initial value is 1
- -: The initial value is undefined

(4) R/W

For each bit and bit field, this entry indicates whether the bit or field is readable or writable, or both writing to and reading from the bit or field are impossible.

The notation is as follows:

R/W: The bit or field is readable and writable.

R/(W): The bit or field is readable and writable.

However, writing is only performed to flag clearing.

R: The bit or field is readable.

"R" is indicated for all reserved bits. When writing to the register, write the value under Initial Value in the bit chart to reserved bits or fields.

W: The bit or field is writable.

(5) Description

Describes the function of the bit or field and specifies the values for writing.

4. Description of Abbreviations

The abbreviations used in this manual are listed below.

• Abbreviations used in this manual

Abbreviation	Description
ACIA	Asynchronous communication interface adapter
bps	Bits per second
CRC	Cyclic redundancy check
DMA	Direct memory access
DMAC	Direct memory access controller
GSM	Global System for Mobile Communications
Hi-Z	High impedance
IEBus	Inter Equipment Bus (IEBus is a trademark of NEC Electronics Corporation.)
I/O	Input/output
IrDA	Infrared Data Association
LSB	Least significant bit
MSB	Most significant bit
NC	No connection
PLL	Phase-locked loop
PWM	Pulse width modulation
SFR	Special function register
SIM	Subscriber Identity Module
UART	Universal asynchronous receiver/transmitter
VCO	Voltage-controlled oscillator

5. List of Product Specifications

Operating temperature

Below is a table listing the product specifications for each group.

	2 2	_	H8/38537 Group
Item		Flash Memory	Mask ROM
Memory	ROM	32 K, 60 Kbytes	16 K, 24 K, 32 K, 40 K, 48 K, 60 Kbytes
	RAM	2 Kbytes	1 Kbyte, 2 Kbytes
Operating	4.5 to 5.5 V	16 MHz	16 MHz
voltage and operating	2.7 to 5.5 V	16 MHz	16 MHz
frequency	1.8 to 5.5 V	_	_
	2.7 to 3.6 V	_	_
	1.8 to 3.6 V	_	_
I/O ports	Input	9	9
	Output	_	_
	I/O	55	55
Timers	Clock (timer A)	1	1
	Reload (timer C)	1	1
	Compare (timer F)	1	1
	Capture (timer G)	1	1
	AEC	1	1
	WDT	_	-
	WDT (discrete)	1	1
SCI	UART/Synchronous	2 ch	2 ch
$\begin{array}{c} \text{A-D} \\ \text{(resolution} \times \end{array}$	input channels)	10 bit × 8 ch	10 bit × 8 ch
LCD	seg	32	32
	com	4	4
External inter (internal wak		13(8)	13(8)
Package		FP-80A	FP-80A
		TFP-80C	TFP-80C

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Standard specifications: -20 to 75°C, WTR: -40 to 85°C

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Section 1 Overview

1.1 Features

Microcontrollers of the H8/38537 Group are CISC (complex instruction set computer) microcontrollers whose core is an H8/300H CPU, which has an internal 32-bit architecture. The H8/300H CPU provides upward compatibility with the H8/300 CPUs of other Renesas Technology-original microcontrollers.

As peripheral functions, each LSI of this Group includes various timer functions that realize low-cost configurations for end systems. The power consumption of these modules can be kept down dynamically by power-down mode.

1.1.1 Application

Examples of the applications of this LSI include motor control, power meter, and health equipment.

1.1.2 Overview of Specifications

Table 1.1 lists the functions of H8/38537 Group products in outline.

Table 1.1 Overview of Functions

Classification	Module/ Function	Description
Memory	ROM	ROM lineup: Flash memory version and mask Rom version
		 ROM capacity: 16 K, 24 K, 32 K, 40 K, 48 K, and 60 Kbytes
	RAM	RAM capacity: 1 Kbyte and 2 Kbytes
CPU	CPU	H8/300H CPU (CISC type)
		Upward compatibility for H8/300 CPU at object level
		Sixteen 16-bit general registers
		Eight addressing modes
		64-Kbyte address space
		Program: 64 Kbytes available
		Data: 64 Kbytes available
		 62 basic instructions, classifiable as bit arithmetic and logic instructions, multiply and divide instructions, bit manipulation instructions, and others
		• Minimum instruction execution time: 400 ns (for an ADD instruction while system clock $\varphi=5$ MHz and $V_{\text{cc}}=2.7$ to 3.6 V)
		 On-chip multiplier (16 × 16 → 32 bits)
	Operating mode	Normal mode
	MCU	Mode: Single-chip mode
	operating mode	 Low power consumption state (transition driven by the SLEEP instruction)
Interrupt	Interrupt	Thirteen external interrupt pins (IRQ4 to IRQ0, WKP7 to WKP0)
(source)	controller (INTC)	23 internal interrupt sources
	(IIVIC)	Independent vector addresses

Classification	Module/ Function	Description
Clock	Clock pulse generator (CPG)	 Two clock generation circuits available Separate clock signals are provided for each of functional modules Includes frequency division circuit, so the operating frequency is selectable Seven low-power-consumption modes: Active (medium speed) mode, sleep (high speed or medium speed) mode, subactive mode, subsleep mode, standby mode, and watch mode
A/D converter	A/D converter (ADC)	 10-bit resolution × eight input channels Sample and hold function included Conversion time: 12.4 μs per channel (with φ at 5-MHz operation) A/D conversion can be started by external trigger input
Timer	14-bit PWM	 14 bits × one channel Four conversion periods selectable Pulse division method for less ripple
	Timer A	 8-bit timer Interval timer functionality: Eight internal clock sources are selectable Clock time base functionality: Four overflow periods are selectable Generates an interrupt upon overflow Timer output clock signals are selectable
	Timer C	 8-bit timer Eight clocks are selectable Auto-reload function supported Generates an interrupt upon overflow Up/down-counter switching is possible
	Timer F	 16-bit timer (also can be used as two independent 8-bit timers) Five clocks are selectable Output compare function supported Toggle output function supported Two interrupt sources: Compare match and overflow

Classification	Module/ Function	Description		
Timer	Timer G	8-bit timer		
		Four counter input clocks are selectable		
		Input capture functions supported (a built-in noise canceller)		
		Level detection at counter overflow is possible		
		Counter clearing option		
		Two interrupt sources: Input capture and overflow		
	Asynchron-	• 16-bit pulse timer (also can be used as 8 bits × two channels)		
	ous event	Can count asynchronously-input external events		
	counter (AEC)			
Watchdog timer	· ,	8 bits × one channel (selectable from two counter input clocks)		
vvateridog timer	timer (WDT)	o site / one onarmor (constitution the ocume) impart disorter,		
Serial interface	Serial communi-	Two serial communication interfaces, SCI3-1 and SCI3-2, that have identical functions		
	cation interface 3	 For both asynchronous and clock synchronous serial communications 		
(SCI3) •		Full-duplex communications capability		
		Select the desired bit rate		
		Six interrupt sources		
I/O ports		Nine CMOS input-only pins		
		55 CMOS input/output pins		
		32 pull-up resistors		
LCD (Liquid	LCD	A maximum of 32 segment pins and four common pins		
Crystal Display)		 Choice of four duty cycles (static, 1/2, 1/3, or 1/4) 		
drive	driver	 LCD RAM capacity: 8 bits x 32 bytes (256 bits) 		
		Word access to LCD RAM		
	•	 All eight segment output pins can be used individually as port pins 		
		• Common output pins not used because of the duty cycle can be used for common double-buffering (parallel connection)		
		Display possible in operating modes other than standby mode		
		Choice of 11 frame frequencies		
		 Built-in power supply split-resistance, supplying LCD drive power 		
		A or B waveform selectable by software		



Classification	Module/ Function	Description		
Internal power supply step-down circuit	Power supply circuit	• The internal power supply can be fixed at a constant level of approximately 3.0 V to 3.2 V, independently of the voltage of the power supply connected to the external V_{cc} pin		
		 It is also possible to use the same level of external power supply voltage and internal power supply voltage without using the internal power supply step-down circuit 		
Package		QFP-80: package code: FP-80A (package dimensions: 14 × 14 mm, pin pitch: 0.65 mm)		
		 TQFP-80: package code: TFP-80C (package dimensions: 12 × 12 mm, pin pitch: 0.50 mm) 		
Operating frequency/		Operating frequency: 2 to 16 MHz		
Power supply vo	oltage	Power supply voltage:		
		Vcc = 2.7 to 5.5 V, AVcc = 2.7 to 5.5 V		
		Supply current:		
		Flash memory version: 4.9 mA (typ.) (Vcc = 5.0 V, AVcc = 5.0 V, ϕ = 10 MHz)		
		Mask ROM version: 4.0 mA (typ.) $(Vcc = 5.0 \text{ V, AVcc} = 5.0 \text{ V, } \varphi = 10 \text{ MHz})$		
Operating peripl	heral	−20 to +75°C (regular specifications)		
temperature (°C)		 -40 to +85°C (wide-range specifications) 		

1.2 List of Products

Table 1.2 and figure 1.1 show the list of products and the structure of a product number, respectively.

Table 1.2 List of Products

Group	Product Type	ROM Size	RAM Size	Package	Remarks
H8/38537 Group	HD64F38537	60 Kbytes	2 Kbytes	FP-80A, TFP-80C	Flash memory version
	HD64338537	60 Kbytes	2 Kbytes		Mask ROM version
	HD64338536	48 Kbytes	2 Kbytes		Mask ROM version
	HD64338535	40 Kbytes	2 Kbytes		Mask ROM version
	HD64F38534	32 Kbytes	2 Kbytes		Flash memory version
	HD64338534	32 Kbytes	2 Kbytes		Mask ROM version
	HD64338533	24 Kbytes	1 Kbyte		Mask ROM version
	HD64338532	16 Kbytes	1 Kbyte		Mask ROM version

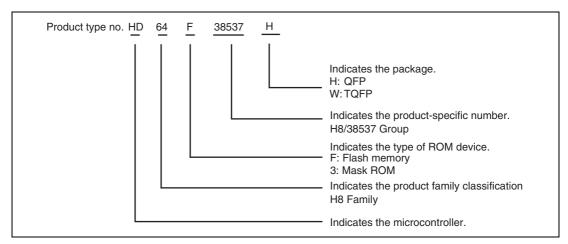


Figure 1.1 How to Read the Product Name Code

1.3 Block Diagram

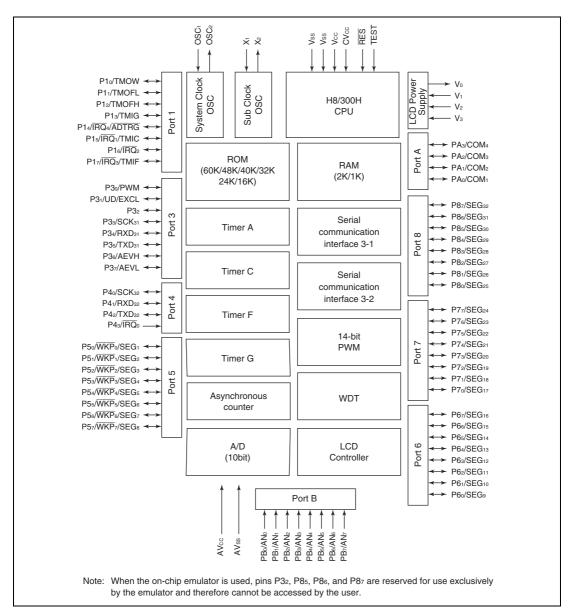


Figure 1.2 Block Diagram of H8/38537 Group

1.4 Pin Assignment

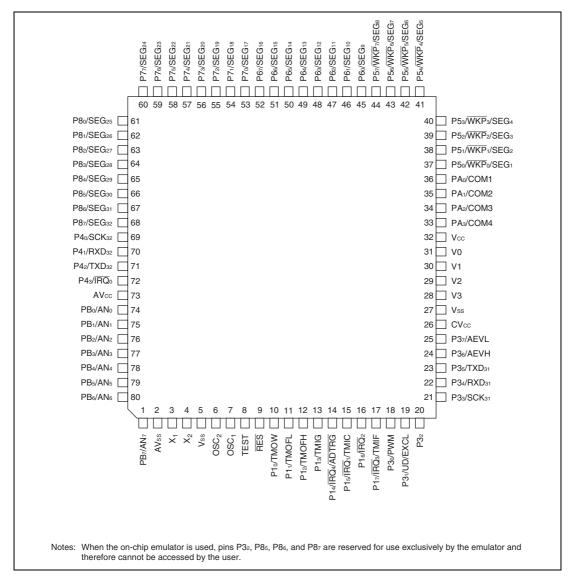


Figure 1.3 Pin Assignment of H8/38537 Group (FP-80A and TFP-80C)

1.5 Pin Functions

Table 1.3 Pin Functions

		Pin No.		
Туре	Symbol	FP-80A, TFP-80C	I/O	Name and Functions
Power source pins	V _{cc} CV _{cc}	32 26	Input	Power supply: All V_{cc} pins should be connected to the system power supply.
				See section 14, Power Supply Circuit, for a CV_cc pin.
	V _{ss}	5 27	Input	Ground: All V _{ss} pins should be connected to the system power supply (0 V).
	AV _{cc}	73	Input	Analog power supply: This is the power supply pin for the A/D converter. When the A/D converter is not used, connect this pin to the system power supply.
	AV _{ss}	2	Input	Analog ground: This is the A/D converter ground pin. It should be connected to the system power supply (0V).
	V0	31	Output	
	V1	30	Input	supply pins for the LCD controller/driver. They incorporate a power supply split-
	V2 V3	29 28		resistance, and are normally used with V0 and V1 shorted.
Clock pins	OSC1	7	Input	These pins connect to a crystal or ceramic
	OSC2	6	Output	oscillator, or can be used to input an external clock. See section 4, Clock Pulse Generators, for a typical connection diagram.
	X1	3	Input	These pins connect to a 32.768-kHz or
	X2	4	Output	38.4-kHz crystal oscillator.
				See section 4, Clock Pulse Generators, for a typical connection diagram.
	EXCL	19	Input	This pin connects to a 32.768-kHz or 38.4-kHz external clock. See section 4, Clock Pulse Generators, for typical connection diagram.

		Pin No.		
Туре	Symbol	FP-80A, TFP-80C	I/O	Name and Functions
System control	RES	9	Input	Reset: When this pin is driven low, the chip is reset
	TEST	8	Input	Test pin: This pin is reserved and cannot be used. It should be connected to V _{ss} .
Interrupt pins	IRQ0 IRQ1 IRQ2 IRQ3 IRQ4	72 15 16 17 14	Input	IRQ interrupt request 0 to 4: These are input pins for edge-sensitive external interrupts, with a selection of rising or falling edge
	WKP7 to WKP0	44 to 37	Input	Wakeup interrupt request 0 to 7: These are input pins for rising or falling-edgesensitive external interrupts.
Timer pins	TMOW	10	Output	Clock output: This is an output pin for waveforms generated by the timer A output circuit.
	AEVL AEVH	25 24	Input	Asynchronous event counter event input: This is an event input pin for input to the asynchronous event counter.
	TMIC	15	Input	Timer C event input: This is an event input pin for input to the timer C counter.
	UD	19	Input	Timer C up/down select: This pin selects up- or down-counting for the timer C counter. The counter operates as a down-counter when this pin is high, and as an up-counter when low.
	TMIF	17	Input	Timer F event input: This is an event input pin for input to the timer F counter.
	TMOFL	11	Output	Timer FL output: This is an output pin for waveforms generated by the timer FL output compare function.
	TMOFH	12	Output	Timer FH output: This is an output pin for waveforms generated by the timer FH output compare function.
	TMIG	13	Input	Timer G capture input: This is an input pin for timer G input capture.
14-bit PWM pin	PWM	18	Output	14-bit PWM output: This is an output pin for waveforms generated by the 14-bit PWM

		Pin No.		
Туре	Symbol	FP-80A, TFP-80C	I/O	Name and Functions
I/O ports	PB7 to PB0	1, 80 to 74	Input	Port B: This is an 8-bit input port.
	P43	72	Input	Port 4 (bit 3): This is a 1-bit input port.
	P42 to P40	71 to 69	I/O	Port 4 (bits 2 to 0): This is a 3-bit I/O port. Input or output can be designated for each bit by means of port control register 4 (PCR4).
	PA3 to PA0	33 to 36	I/O	Port A: This is a 4-bit I/O port. Input or output can be designated for each bit by means of port control register A (PCRA).
	P17 to P10	17 to 10	I/O	Port 1: This is an 8-bit I/O port. Input or output can be designated for each bit by means of port control register 1 (PCR1).
	P37 to P30	25 to 18	I/O	Port 3: This is an 8-bit I/O port. Input or output can be designated for each bit by means of port control register 3 (PCR3). When the on-chip emulator is used, pin P32 is reserved for use exclusively by the emulator and therefore cannot be accessed by the user. With the flash memory version, pull up pin P32 to high level to cancel a reset in the in the user mode.
	P57 to P50	44 to 37	I/O	Port 5: This is an 8-bit I/O port. Input or output can be designated for each bit by means of port control register 5 (PCR5).
	P67 to P60	52 to 45	I/O	Port 6: This is an 8-bit I/O port. Input or output can be designated for each bit by means of port control register 6 (PCR6).
	P77 to P70	60 to 53	I/O	Port 7: This is an 8-bit I/O port. Input or output can be designated for each bit by means of port control register 7 (PCR7).
	P87 to P80	68 to 61	I/O	Port 8: This is an 8-bit I/O port. Input or output can be designated for each bit by means of port control register 8 (PCR8). When the on-chip emulator is used, pins P85, P86, and P87 are reserved for use exclusively by the emulator and therefore cannot be accessed by the user.



		Pin No.		
Туре	Symbol	FP-80A, TFP-80C	1/0	Name and Functions
Serial communi-	RXD31	22	Input	SCI31 receive data input: This is the SCI31 data input pin.
cation interface (SCI)	TXD31	23	Output	SCI31 transmit data output: This is the SCI31 data output pin.
(001)	SCK31	21	I/O	SCI31 clock I/O: This is the SCI31 clock I/O pin.
	RXD32	70	Input	SCI32 receive data input: This is the SCI32 data input pin.
	TXD32	71	Output	SCI32 transmit data output: This is the SCI32 data output pin.
	SCK32	69	I/O	SCI32 clock I/O: This is the SCI32 clock I/O pin.
A/D converter	AN7 to AN0	1 80 to 74	Input	Analog input channels 7 to 0: These are analog data input channels to the A/D converter
	ADTRG	14	Input	A/D converter trigger input: This is the external trigger input pin to the A/D converter
LCD controller/	COM4 to COM1	33 to 36	Output	LCD common output: These are the LCD common output pins.
driver	SEG32 to SEG1	68 to 37	Output	LCD segment output: These are the LCD segment output pins.

Section 2 CPU

This LSI has an H8/300H CPU with an internal 32-bit architecture that is upward-compatible with the H8/300 CPU, and supports only normal mode, which has a 64-Kbyte address space.

• Upward-compatible with H8/300 CPUs

Can execute H8/300 CPUs object programs

Additional eight 16-bit extended registers

32-bit transfer and arithmetic and logic instructions are added

Signed multiply and divide instructions are added.

• General-register architecture

Sixteen 16-bit general registers also usable as sixteen 8-bit registers and eight 16-bit registers, or eight 32-bit registers

• Sixty-two basic instructions

8/16/32-bit data transfer and arithmetic and logic instructions

Multiply and divide instructions

Powerful bit-manipulation instructions

• Eight addressing modes

Register direct [Rn]

Register indirect [@ERn]

Register indirect with displacement [@(d:16,ERn) or @(d:24,ERn)]

Register indirect with post-increment or pre-decrement [@ERn+ or @-ERn]

Absolute address [@aa:8, @aa:16, @aa:24]

Immediate [#xx:8, #xx:16, or #xx:32]

Program-counter relative [@(d:8,PC) or @(d:16,PC)]

Memory indirect [@@aa:8]

- 64-Kbyte address space
- High-speed operation

All frequently-used instructions execute in one or two states

8/16/32-bit register-register add/subtract : 2 state 8×8 -bit register-register multiply : 14 states $16 \div 8$ -bit register-register divide : 14 states 16×16 -bit register-register multiply : 22 states $32 \div 16$ -bit register-register divide : 22 states

Power-down state
 Transition to power-down state by SLEEP instruction

2.1 Address Space and Memory Map

The memory map of the H8/38537 is shown in figure 2.1(1), that of the H8/38536 in figure 2.1(2), that of the H8/38535 in figure 2.1(3), that of the H8/38534 in figure 2.1(4), that of the H8/38533 in figure 2.1(5), and that of the H8/38532 in figure 2.1(6).



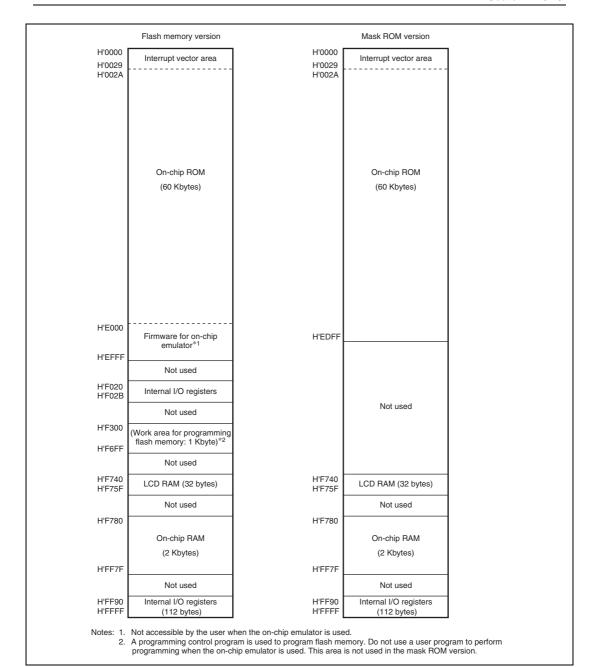


Figure 2.1 (1) H8/38537 Memory Map

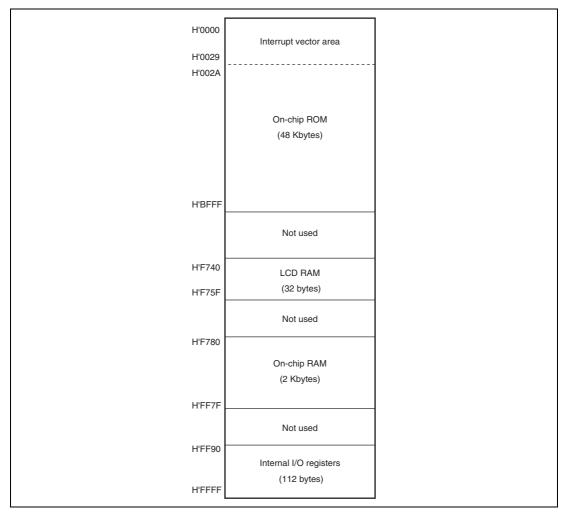


Figure 2.1 (2) H8/38536 Memory Map

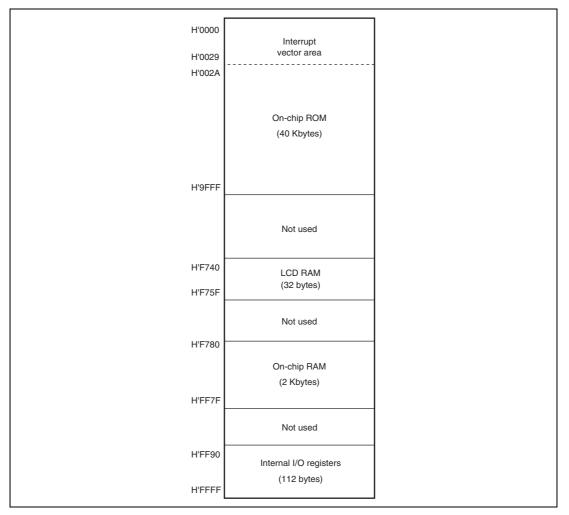


Figure 2.1 (3) H8/38535 Memory Map

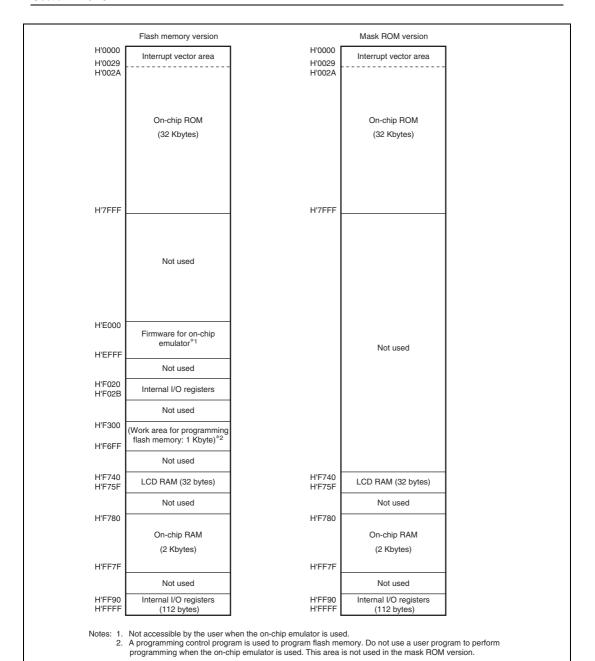


Figure 2.1 (4) H8/38534 Memory Map

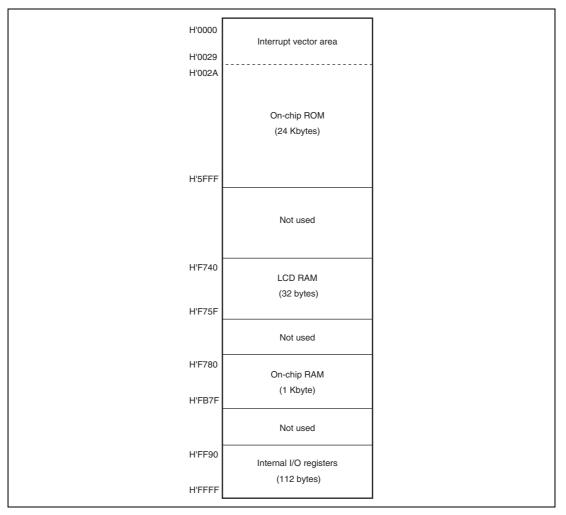


Figure 2.1 (5) H8/38533 Memory Map

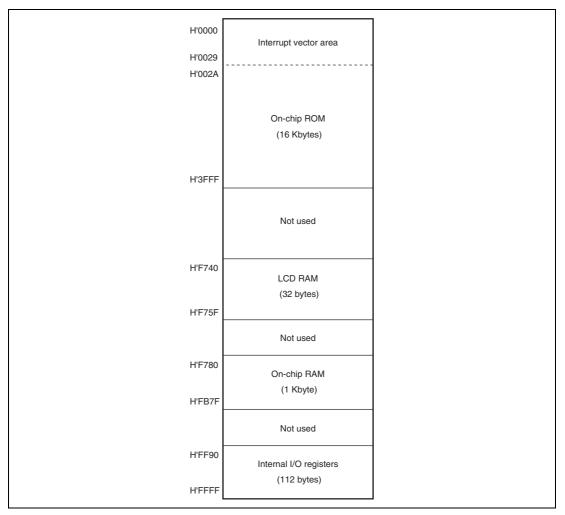


Figure 2.1 (6) H8/38532 Memory Map

2.2 Register Configuration

The H8/300H CPU has the internal registers shown in figure 2.2. There are two types of registers; general registers and control registers. The control registers are a 24-bit program counter (PC), and an 8-bit condition-code register (CCR).

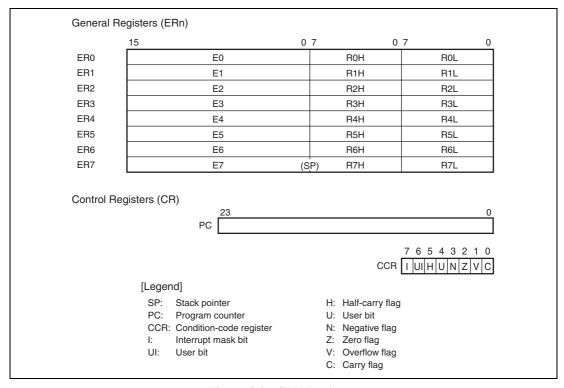


Figure 2.2 CPU Registers

2.2.1 General Registers

The H8/300H CPU has eight 32-bit general registers. These general registers are all functionally identical and can be used as both address registers and data registers. When a general register is used as a data register, it can be accessed as a 32-bit, 16-bit, or 8-bit register. Figure 2.3 illustrates the usage of the general registers. When the general registers are used as 32-bit registers or address registers, they are designated by the letters ER (ER0 to ER7).

The ER registers divide into 16-bit general registers designated by the letters E (E0 to E7) and R (R0 to R7). These registers are functionally equivalent, providing a maximum of sixteen 16-bit registers. The E registers (E0 to E7) are also referred to as extended registers.

The R registers divide into 8-bit registers designated by the letters RH (R0H to R7H) and RL (R0L to R7L). These registers are functionally equivalent, providing a maximum of sixteen 8-bit registers.

The usage of each register can be selected independently.

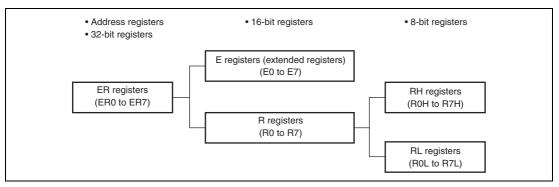


Figure 2.3 Usage of General Registers

General register ER7 has the function of the stack pointer (SP) in addition to its general-register function, and is used implicitly in exception handling and subroutine calls. Figure 2.4 shows the relationship between the stack pointer and the stack area.

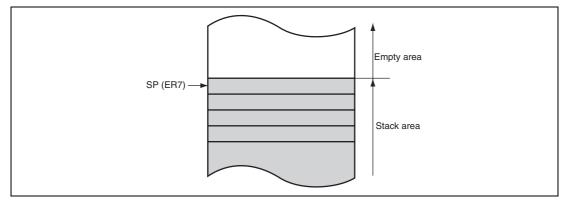


Figure 2.4 Relationship between Stack Pointer and Stack Area

2.2.2 Program Counter (PC)

This 24-bit counter indicates the address of the next instruction the CPU will execute. The length of all CPU instructions is 2 bytes (one word), so the least significant PC bit is ignored. (When an instruction is fetched, the least significant PC bit is regarded as 0). The PC is initialized when the start address is loaded by the vector address generated during reset exception-handling sequence.

2.2.3 Condition-Code Register (CCR)

This 8-bit register contains internal CPU status information, including an interrupt mask bit (I) and half-carry (H), negative (N), zero (Z), overflow (V), and carry (C) flags. The I bit is initialized to 1 by reset exception-handling sequence, but other bits are not initialized.

Some instructions leave flag bits unchanged. Operations can be performed on the CCR bits by the LDC, STC, ANDC, ORC, and XORC instructions. The N, Z, V, and C flags are used as branching conditions for conditional branch (Bcc) instructions.

For the action of each instruction on the flag bits, see appendix A.1, Instruction List.

Bit	Bit Name	Initial Value	R/W	Description
7	I	1	R/W	Interrupt Mask Bit
				Masks interrupts other than NMI when set to 1. NMI is accepted regardless of the I bit setting. The I bit is set to 1 at the start of an exception-handling sequence.
6	UI	Undefined	R/W	User Bit
				Can be written and read by software using the LDC, STC, ANDC, ORC, and XORC instructions.
5	Н	Undefined	R/W	Half-Carry Flag
				When the ADD.B, ADDX.B, SUB.B, SUBX.B, CMP.B, or NEG.B instruction is executed, this flag is set to 1 if there is a carry or borrow at bit 3, and cleared to 0 otherwise. When the ADD.W, SUB.W, CMP.W, or NEG.W instruction is executed, the H flag is set to 1 if there is a carry or borrow at bit 11, and cleared to 0 otherwise. When the ADD.L, SUB.L, CMP.L, or NEG.L instruction is executed, the H flag is set to 1 if there is a carry or borrow at bit 27, and cleared to 0 otherwise.
4	U	Undefined	R/W	User Bit
				Can be written and read by software using the LDC, STC, ANDC, ORC, and XORC instructions.
3	N	Undefined	R/W	Negative Flag
				Stores the value of the most significant bit of data as a sign bit.
2	Z	Undefined	R/W	Zero Flag
				Set to 1 to indicate zero data, and cleared to 0 to indicate non-zero data.
1	V	Undefined	R/W	Overflow Flag
				Set to 1 when an arithmetic overflow occurs, and cleared to 0 at other times.
0	С	Undefined	R/W	Carry Flag
				Set to 1 when a carry occurs, and cleared to 0 otherwise. Used by:
				Add instructions, to indicate a carry
				Subtract instructions, to indicate a borrow
				Shift and rotate instructions, to indicate a carry
				The carry flag is also used as a bit accumulator by bit manipulation instructions.



2.3 Data Formats

The H8/300H CPU can process 1-bit, 4-bit (BCD), 8-bit (byte), 16-bit (word), and 32-bit (longword) data. Bit-manipulation instructions operate on 1-bit data by accessing bit n (n = 0, 1, 2, ..., 7) of byte operand data. The DAA and DAS decimal-adjust instructions treat byte data as two digits of 4-bit BCD data.

2.3.1 General Register Data Formats

Figure 2.5 shows the data formats in general registers.

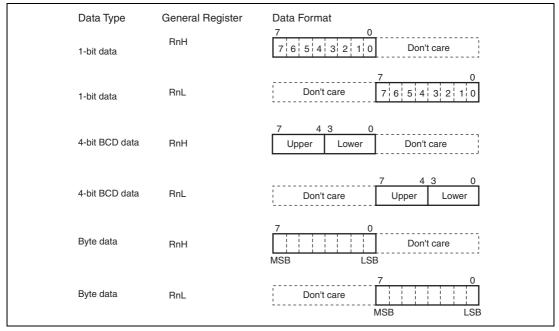


Figure 2.5 General Register Data Formats (1)

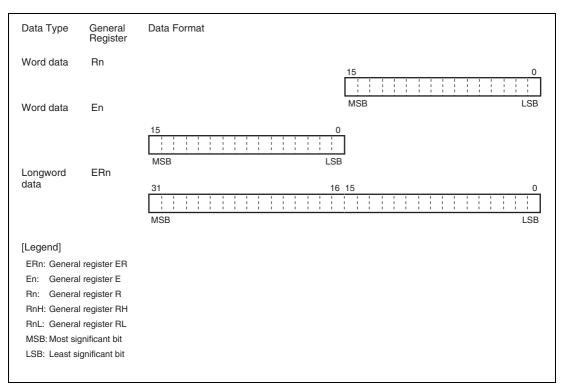


Figure 2.5 General Register Data Formats (2)

2.3.2 Memory Data Formats

Figure 2.6 shows the data formats in memory. The H8/300H CPU can access word data and longword data in memory, however word or longword data must begin at an even address. If an attempt is made to access word or longword data at an odd address, an address error does not occur, however the least significant bit of the address is regarded as 0, so access begins the preceding address. This also applies to instruction fetches.

When ER7 (SP) is used as an address register to access the stack area, the operand size should be word or longword.

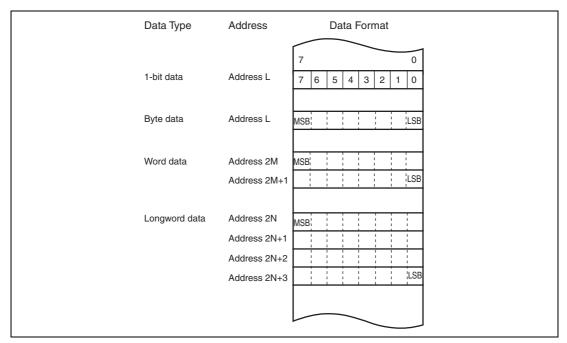


Figure 2.6 Memory Data Formats

2.4 Instruction Set

2.4.1 Table of Instructions Classified by Function

The H8/300H CPU has 62 instructions. Tables 2.2 to 2.9 summarize the instructions in each functional category. The notation used in tables 2.2 to 2.9 is defined in table 2.1.

Table 2.1 Operation Notation

Symbol	Description
Rd	General register (destination)*
Rs	General register (source)*
Rn	General register*
ERn	General register (32-bit register or address register)
(EAd)	Destination operand
(EAs)	Source operand
CCR	Condition-code register
N	N (negative) flag in CCR
Z	Z (zero) flag in CCR
V	V (overflow) flag in CCR
С	C (carry) flag in CCR
PC	Program counter
SP	Stack pointer
#IMM	Immediate data
disp	Displacement
+	Addition
_	Subtraction
×	Multiplication
÷	Division
٨	Logical AND
V	Logical OR
\oplus	Logical XOR
\rightarrow	Move
7	NOT (logical complement)

Symbol	Description
:3/:8/:16/:24	3-, 8-, 16-, or 24-bit length

Note: * General registers include 8-bit registers (R0H to R7H, R0L to R7L), 16-bit registers (R0 to R7, E0 to E7), and 32-bit registers/address register (ER0 to ER7).

Table 2.2 Data Transfer Instructions

Instruction	Size*	Function
MOV	B/W/L	$(EAs) \rightarrow Rd, Rs \rightarrow (EAd)$ Moves data between two general registers or between a general register and memory, or moves immediate data to a general register.
MOVFPE	В	$(EAs) \rightarrow Rd$ Cannot be used in this LSI.
MOVTPE	В	$Rs \rightarrow (EAs)$ Cannot be used in this LSI.
POP	W/L	@SP+ \rightarrow Rn Pops a general register from the stack. POP.W Rn is identical to MOV.W @SP+, Rn. POP.L ERn is identical to MOV.L @SP+, ERn.
PUSH	W/L	$Rn \rightarrow @-SP$ Pushes a general register onto the stack. PUSH.W Rn is identical to MOV.W Rn, @-SP. PUSH.L ERn is identical to MOV.L ERn, @-SP.

Note: * Refers to the operand size.

Table 2.3 Arithmetic Operations Instructions (1)

Instruction	Size*	Function
ADD SUB	B/W/L	Rd \pm Rs \rightarrow Rd, Rd \pm #IMM \rightarrow Rd Performs addition or subtraction on data in two general registers, or on immediate data and data in a general register (immediate byte data cannot be subtracted from byte data in a general register. Use the SUBX or ADD instruction.)
ADDX SUBX	В	Rd \pm Rs \pm C \rightarrow Rd, Rd \pm #IMM \pm C \rightarrow Rd Performs addition or subtraction with carry on byte data in two general registers, or on immediate data and data in a general register.
INC DEC	B/W/L	$Rd \pm 1 \rightarrow Rd$, $Rd \pm 2 \rightarrow Rd$ Increments or decrements a general register by 1 or 2. (Byte operands can be incremented or decremented by 1 only.)
ADDS SUBS	L	$Rd \pm 1 \rightarrow Rd$, $Rd \pm 2 \rightarrow Rd$, $Rd \pm 4 \rightarrow Rd$ Adds or subtracts the value 1, 2, or 4 to or from data in a 32-bit register.
DAA DAS	В	Rd (decimal adjust) \rightarrow Rd Decimal-adjusts an addition or subtraction result in a general register by referring to the CCR to produce 4-bit BCD data.
MULXU	B/W	$Rd \times Rs \rightarrow Rd$ Performs unsigned multiplication on data in two general registers: either 8 bits \times 8 bits \rightarrow 16 bits or 16 bits \times 16 bits \rightarrow 32 bits.
MULXS	B/W	$Rd \times Rs \rightarrow Rd$ Performs signed multiplication on data in two general registers: either 8 bits \times 8 bits \rightarrow 16 bits or 16 bits \times 16 bits \rightarrow 32 bits.
DIVXU	B/W	Rd \div Rs \rightarrow Rd Performs unsigned division on data in two general registers: either 16 bits \div 8 bits \rightarrow 8-bit quotient and 8-bit remainder or 32 bits \div 16 bits \rightarrow 16-bit quotient and 16-bit remainder.



Table 2.3 Arithmetic Operations Instructions (2)

Instruction	Size*	Function
DIVXS	B/W	Rd \div Rs \rightarrow Rd Performs signed division on data in two general registers: either 16 bits \div 8 bits \rightarrow 8-bit quotient and 8-bit remainder or 32 bits \div 16 bits \rightarrow 16-bit quotient and 16-bit remainder.
CMP	B/W/L	Rd – Rs, Rd – #IMM Compares data in a general register with data in another general register or with immediate data, and sets CCR bits according to the result.
NEG	B/W/L	$0-\text{Rd}\to\text{Rd}$ Takes the two's complement (arithmetic complement) of data in a general register.
EXTU	W/L	Rd (zero extension) → Rd Extends the lower 8 bits of a 16-bit register to word size, or the lower 16 bits of a 32-bit register to longword size, by padding with zeros on the left.
EXTS	W/L	Rd (sign extension) → Rd Extends the lower 8 bits of a 16-bit register to word size, or the lower 16 bits of a 32-bit register to longword size, by extending the sign bit.

Table 2.4 Logic Operations Instructions

Instruction	Size*	Function
AND	B/W/L	$Rd \wedge Rs \rightarrow Rd$, $Rd \wedge \#IMM \rightarrow Rd$ Performs a logical AND operation on a general register and another general register or immediate data.
OR	B/W/L	$Rd \lor Rs \to Rd, Rd \lor \#IMM \to Rd$ Performs a logical OR operation on a general register and another general register or immediate data.
XOR	B/W/L	$Rd \oplus Rs \rightarrow Rd$, $Rd \oplus \#IMM \rightarrow Rd$ Performs a logical exclusive OR operation on a general register and another general register or immediate data.
NOT	B/W/L	\neg (Rd) \rightarrow (Rd) Takes the one's complement (logical complement) of general register contents.

B: Byte W: Word L: Longword

Table 2.5 Shift Instructions

Instruction	Size*	Function
SHAL SHAR	B/W/L	Rd (shift) $\rightarrow Rd$ Performs an arithmetic shift on general register contents.
SHLL SHLR	B/W/L	Rd (shift) $\rightarrow Rd$ Performs a logical shift on general register contents.
ROTL ROTR	B/W/L	Rd (rotate) $\rightarrow Rd$ Rotates general register contents.
ROTXL ROTXR	B/W/L	Rd (rotate) \rightarrow Rd Rotates general register contents through the carry flag.

Note: * Refers to the operand size.



Table 2.6 Bit Manipulation Instructions (1)

Instruction	Size*	Function
BSET	В	1 → (<bit-no.> of <ead>) Sets a specified bit in a general register or memory operand to 1. The bit number is specified by 3-bit immediate data or the lower three bits of a general register.</ead></bit-no.>
BCLR	В	$0 \rightarrow$ (bit-No.> of <ead>) Clears a specified bit in a general register or memory operand to 0. The bit number is specified by 3-bit immediate data or the lower three bits of a general register.</ead>
BNOT	В	¬ (<bit-no.> of <ead>) → (<bit-no.> of <ead>) Inverts a specified bit in a general register or memory operand. The bit number is specified by 3-bit immediate data or the lower three bits of a general register.</ead></bit-no.></ead></bit-no.>
BTST	В	\neg (<bit-no.> of <ead>) \rightarrow Z Tests a specified bit in a general register or memory operand and sets or clears the Z flag accordingly. The bit number is specified by 3-bit immediate data or the lower three bits of a general register.</ead></bit-no.>
BAND	В	$C \wedge (\text{-bit-No}) \text{ of -cEAd}) \rightarrow C$ ANDs the carry flag with a specified bit in a general register or memory operand and stores the result in the carry flag.
BIAND	В	$C \land \neg$ (<bit-no.> of <ead>) $\rightarrow C$ ANDs the carry flag with the inverse of a specified bit in a general register or memory operand and stores the result in the carry flag.</ead></bit-no.>
		The bit number is specified by 3-bit immediate data.
BOR	В	$C \lor (\text{-}\text{bit-No.}\text{-}\text{of} < \text{EAd-}\text{-}) \to C$ ORs the carry flag with a specified bit in a general register or memory operand and stores the result in the carry flag.
BIOR	В	$C \lor \neg$ (<bit-no.> of <ead>) $\to C$ ORs the carry flag with the inverse of a specified bit in a general register or memory operand and stores the result in the carry flag.</ead></bit-no.>
		The bit number is specified by 3-bit immediate data.

B: Byte

Table 2.6 Bit Manipulation Instructions (2)

Instruction	Size*	Function
BXOR	В	$C \oplus (\text{-bit-No}) \text{ of -EAd}) \to C$ XORs the carry flag with a specified bit in a general register or memory operand and stores the result in the carry flag.
BIXOR	В	$C \oplus \neg$ (<bit-no.> of <ead>) $\rightarrow C$ XORs the carry flag with the inverse of a specified bit in a general register or memory operand and stores the result in the carry flag.</ead></bit-no.>
		The bit number is specified by 3-bit immediate data.
BLD	В	(<bit-no.> of <ead>) \rightarrow C Transfers a specified bit in a general register or memory operand to the carry flag.</ead></bit-no.>
BILD	В	\neg (<bit-no.> of <ead>) \rightarrow C Transfers the inverse of a specified bit in a general register or memory operand to the carry flag.</ead></bit-no.>
		The bit number is specified by 3-bit immediate data.
BST	В	C → (<bit-no.> of <ead>) Transfers the carry flag value to a specified bit in a general register or memory operand.</ead></bit-no.>
BIST	В	\neg C \rightarrow (<bit-no.> of <ead>) Transfers the inverse of the carry flag value to a specified bit in a general register or memory operand.</ead></bit-no.>
		The bit number is specified by 3-bit immediate data.

B: Byte

Table 2.7 Branch Instructions

Instruction	Size	Function			
Bcc*	_	Branches to a specified address if a specified condition is true. The branching conditions are listed below.			
		Mnemonic	Description	Condition	
		BRA(BT)	Always (true)	Always	
		BRN(BF)	Never (false)	Never	
		BHI	High	C ∨ Z = 0	
		BLS	Low or same	C ∨ Z = 1	
		BCC(BHS)	Carry clear (high or same)	C = 0	
		BCS(BLO)	Carry set (low)	C = 1	
		BNE	Not equal	Z = 0	
		BEQ	Equal	Z = 1	
		BVC	Overflow clear	V = 0	
		BVS	Overflow set	V = 1	
		BPL	Plus	N = 0	
		BMI	Minus	N = 1	
		BGE	Greater or equal	$N \oplus V = 0$	
		BLT	Less than	N ⊕ V = 1	
		BGT	Greater than	$Z_{\vee}(N \oplus V) = 0$	
		BLE	Less or equal	Z∨(N ⊕ V) = 1	
JMP		Branches unco	nditionally to a specified	d address.	
BSR	_	Branches to a	subroutine at a specified	d address.	
JSR	_	Branches to a	subroutine at a specified	d address.	
RTS	_	Returns from a	subroutine		
Note: *	Bcc is the ge	eneral name for co	onditional branch instruc	tions.	

Note: * Bcc is the general name for conditional branch instructions.

Table 2.8 System Control Instructions

Instruction	Size*	Function
RTE	_	Returns from an exception-handling routine.
SLEEP	_	Causes a transition to a power-down state.
LDC	B/W	(EAs) → CCR Moves the source operand contents to the CCR. The CCR size is one byte, but in transfer from memory, data is read by word access.
STC	B/W	$CCR \to (EAd)$ Transfers the CCR contents to a destination location. The condition code register size is one byte, but in transfer to memory, data is written by word access.
ANDC	В	CCR ∧ #IMM → CCR Logically ANDs the CCR with immediate data.
ORC	В	CCR ∨ #IMM → CCR Logically ORs the CCR with immediate data.
XORC	В	$CCR \oplus \#IMM \rightarrow CCR$ Logically XORs the CCR with immediate data.
NOP	—	$PC + 2 \rightarrow PC$ Only increments the program counter.

B: Byte W: Word



Table 2.9 Block Data Transfer Instructions

Instruction	Size	Function
EEPMOV.B	_	if R4L \neq 0 then Repeat @ER5+ \rightarrow @ER6+, R4L-1 \rightarrow R4L Until R4L = 0 else next;
EEPMOV.W	_	if R4 \neq 0 then Repeat @ER5+ \rightarrow @ER6+, R4-1 \rightarrow R4 Until R4 = 0 else next;
		Transfers a data block. Starting from the address set in ER5, transfers data for the number of bytes set in R4L or R4 to the address location set in ER6.
		Execution of the next instruction begins as soon as the transfer is completed.

2.4.2 Basic Instruction Formats

H8/300H CPU instructions consist of 2-byte (1-word) units. An instruction consists of an operation field (op), a register field (r), an effective address extension (EA), and a condition field (cc).

Figure 2.7 shows examples of instruction formats.

(1) Operation Field

Indicates the function of the instruction, the addressing mode, and the operation to be carried out on the operand. The operation field always includes the first four bits of the instruction. Some instructions have two operation fields.

(2) Register Field

Specifies a general register. Address registers are specified by 3 bits, and data registers by 3 bits or 4 bits. Some instructions have two register fields. Some have no register field.

(3) Effective Address Extension

8, 16, or 32 bits specifying immediate data, an absolute address, or a displacement. A24-bit address or displacement is treated as a 32-bit data in which the first 8 bits are 0 (H'00).

(4) Condition Field

Specifies the branching condition of Bcc instructions.

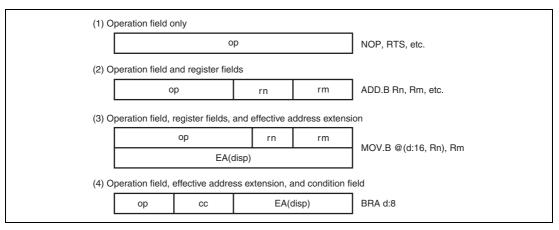


Figure 2.7 Instruction Formats

2.5 Addressing Modes and Effective Address Calculation

The following describes the H8/300H CPU. In this LSI, the upper eight bits are ignored in the generated 24-bit address, so the effective address is 16 bits.

2.5.1 Addressing Modes

The H8/300H CPU supports the eight addressing modes listed in table 2.10. Each instruction uses a subset of these addressing modes. Addressing modes that can be used differ depending on the instruction. For details, refer to appendix A.4, Combinations of Instructions and Addressing Modes.

Arithmetic and logic instructions can use the register direct and immediate modes. Data transfer instructions can use all addressing modes except program-counter relative and memory indirect. Bit-manipulation instructions use register direct, register indirect, or the absolute addressing mode (@aa:8) to specify an operand, and register direct (BSET, BCLR, BNOT, and BTST instructions) or immediate (3-bit) addressing mode to specify a bit number in the operand.

Table 2.10 Addressing Modes

No.	Addressing Mode	Symbol
1	Register direct	Rn
2	Register indirect	@ERn
3	Register indirect with displacement	@(d:16,ERn)/@(d:24,ERn)
4	Register indirect with post-increment Register indirect with pre-decrement	@ERn+ @-ERn
5	Absolute address	@ aa:8/@ aa:16/@ aa:24
6	Immediate	#xx:8/#xx:16/#xx:32
7	Program-counter relative	@(d:8,PC)/@(d:16,PC)
8	Memory indirect	@ @ aa:8

(1) Register Direct—Rn

The register field of the instruction specifies an 8-, 16-, or 32-bit general register containing the operand. R0H to R7H and R0L to R7L can be specified as 8-bit registers. R0 to R7 and E0 to E7 can be specified as 16-bit registers. ER0 to ER7 can be specified as 32-bit registers.



(2) Register Indirect—@ERn

The register field of the instruction code specifies an address register (ERn), the lower 24 bits of which contain the address of the operand on memory.

(3) Register Indirect with Displacement—@(d:16, ERn) or @(d:24, ERn)

A 16-bit or 24-bit displacement contained in the instruction is added to an address register (ERn) specified by the register field of the instruction, and the lower 24 bits of the sum the address of a memory operand. A 16-bit displacement is sign-extended when added.

(4) Register Indirect with Post-Increment or Pre-Decrement—@ERn+ or @-ERn

- Register indirect with post-increment—@ERn+
 - The register field of the instruction code specifies an address register (ERn) the lower 24 bits of which contains the address of a memory operand. After the operand is accessed, 1, 2, or 4 is added to the address register contents (32 bits) and the sum is stored in the address register. The value added is 1 for byte access, 2 for word access, or 4 for longword access. For the word or longword access, the register value should be even.
- Register indirect with pre-decrement—@-ERn
 The value 1, 2, or 4 is subtracted from an address register (ERn) specified by the register field in the instruction code, and the lower 24 bits of the result is the address of a memory operand. The result is also stored in the address register. The value subtracted is 1 for byte access, 2 for word access, or 4 for longword access. For the word or longword access, the register value should be even.

(5) Absolute Address—@aa:8, @aa:16, @aa:24

The instruction code contains the absolute address of a memory operand. The absolute address may be 8 bits long (@aa:8), 16 bits long (@aa:16), 24 bits long (@aa:24)

For an 8-bit absolute address, the upper 16 bits are all assumed to be 1 (H'FFFF). For a 16-bit absolute address the upper 8 bits are a sign extension. A 24-bit absolute address can access the entire address space.

The access ranges of absolute addresses for this LSI are those shown in table 2.11, because the upper 8 bits are ignored.



Table 2.11 Absolute Address Access Ranges

Absolute Address	Access Range	
8 bits (@aa:8)	H'FF00 to H'FFFF	
16 bits (@aa:16)	H'0000 to H'FFFF	
24 bits (@aa:24)	H'0000 to H'FFFF	

(6) Immediate—#xx:8, #xx:16, or #xx:32

The instruction contains 8-bit (#xx:8), 16-bit (#xx:16), or 32-bit (#xx:32) immediate data as an operand.

The ADDS, SUBS, INC, and DEC instructions contain immediate data implicitly. Some bit manipulation instructions contain 3-bit immediate data in the instruction code, specifying a bit number.

(7) Program-Counter Relative—@(d:8, PC) or @(d:16, PC)

This mode is used in the BSR instruction. An 8-bit or 16-bit displacement contained in the instruction is sign-extended and added to the 24-bit PC contents to generate a branch address. The PC value to which the displacement is added is the address of the first byte of the next instruction, so the possible branching range is -126 to +128 bytes (-63 to +64 words) or -32766 to +32768 bytes (-16383 to +16384 words) from the branch instruction. The resulting value should be an even number

(8) Memory Indirect—@@aa:8

This mode can be used by the JMP and JSR instructions. The instruction code contains an 8-bit absolute address specifying a memory operand. This memory operand contains a branch address. The memory operand is accessed in words, generating a 16-bit branch address. Figure 2.8 shows how to specify branch address for in memory indirect mode. The upper bits of the absolute address are all assumed to be 0, so the address range is 0 to 255 (H'0000 to H'00FF).

Note that the first part of the address range is also the exception vector area.



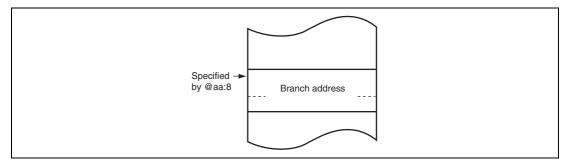


Figure 2.8 Branch Address Specification in Memory Indirect Mode

2.5.2 Effective Address Calculation

Table 2.12 indicates how effective addresses are calculated in each addressing mode. In this LSI the upper 8 bits of the effective address are ignored in order to generate a 16-bit effective address.

Table 2.12 Effective Address Calculation (1)

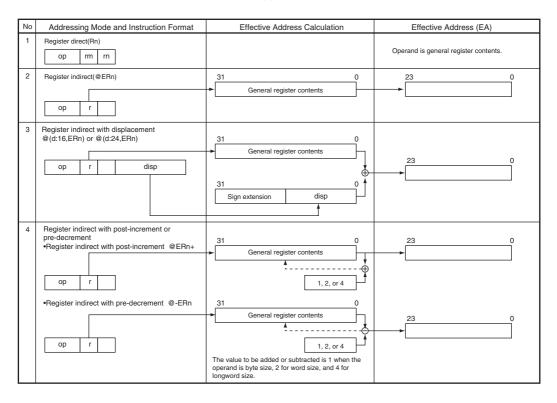
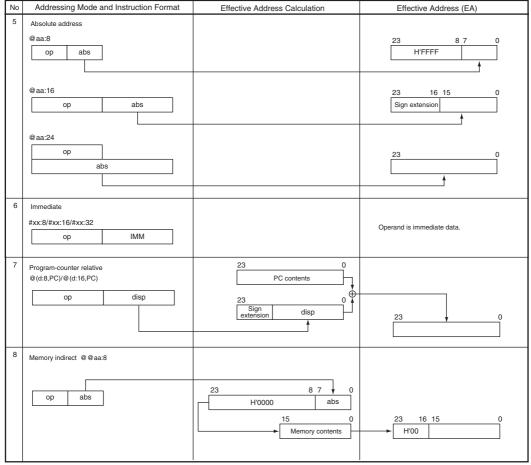


Table 2.12 Effective Address Calculation (2)



[Legend]

r, rm,rn: Register field
op: Operation field
disp: Displacement
IMM: Immediate data
abs: Absolute address

2.6 Basic Bus Cycle

CPU operation is synchronized by a system clock (ϕ) or a subclock (ϕ_{SUB}). The period from a rising edge of ϕ or ϕ_{SUB} to the next rising edge is called one state. A bus cycle consists of two states or three states. The cycle differs depending on whether access is to on-chip memory or to on-chip peripheral modules.

2.6.1 Access to On-Chip Memory (RAM, ROM)

Access to on-chip memory takes place in two states. The data bus width is 16 bits, allowing access in byte or word size. Figure 2.9 shows the on-chip memory access cycle.

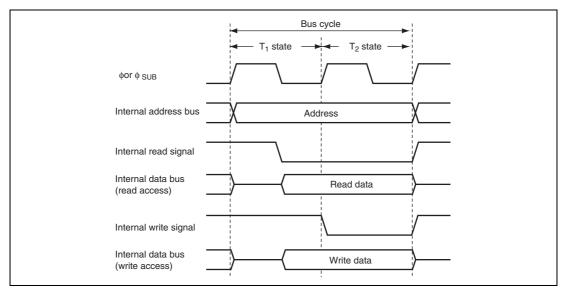


Figure 2.9 On-Chip Memory Access Cycle

2.6.2 On-Chip Peripheral Modules

On-chip peripheral modules are accessed in two states or three states. The data bus width is 8 bits or 16 bits depending on the register. For description on the data bus width and number of accessing states of each register, refer to section 15.1, Register Addresses (Address Order). Registers with 16-bit data bus width can be accessed by word size only. Registers with 8-bit data bus width can be accessed by byte or word size. When a register with 8-bit data bus width is accessed by word size, a bus cycle occurs twice. In two-state access, the operation timing is the same as that for on-chip memory.

Figure 2.10 shows the operation timing in the case of three-state access to an on-chip peripheral module.

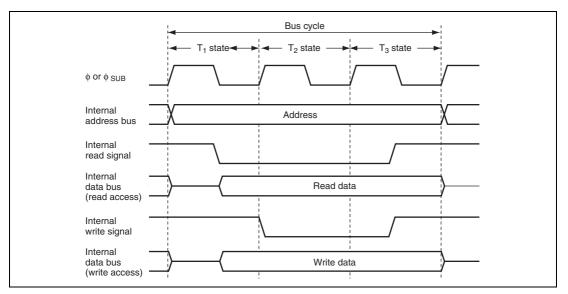


Figure 2.10 On-Chip Peripheral Module Access Cycle (3-State Access)

2.7 CPU States

There are four CPU states: the reset state, program execution state, program halt state, and exception-handling state. The program execution state includes active (high-speed or medium-speed) mode and subactive mode. For the program halt state, there are sleep (high-speed or medium-speed) mode, standby mode, watch mode, and subsleep mode. These states are shown in figure 2.11. Figure 2.12 shows the state transitions. For details on program execution state and program halt state, refer to section 5, Power-Down Modes. For details on exception handling, refer to section 3, Exception Handling.

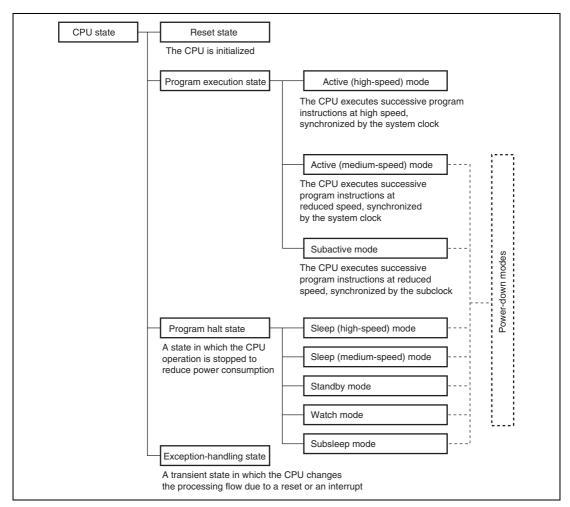


Figure 2.11 CPU Operating States

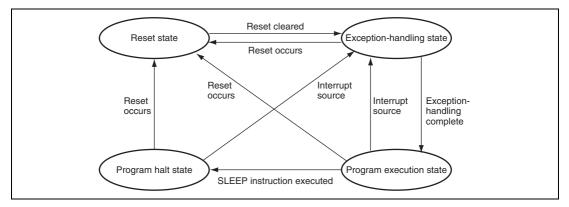


Figure 2.12 State Transitions



2.8 **Usage Notes**

2.8.1 **Notes on Data Access to Empty Areas**

The address space of this LSI includes empty areas in addition to the ROM, RAM, and on-chip I/O registers areas available to the user. When data is transferred from CPU to empty areas, the transferred data will be lost. This action may also cause the CPU to malfunction. When data is transferred from an empty area to CPU, the contents of the data cannot be guaranteed.

2.8.2 **EEPMOV Instruction**

EEPMOV is a block-transfer instruction and transfers the byte size of data indicated by R4L, which starts from the address indicated by R5, to the address indicated by R6. Set R4L and R6 so that the end address of the destination address (value of R6 + R4L) does not exceed H'FFFF (the value of R6 must not change from H'FFFF to H'0000 during execution).

2.8.3 **Bit-Manipulation Instruction**

The BSET, BCLR, BNOT, BST, and BIST instructions read data from the specified address in byte units, manipulate the data of the target bit, and write data to the same address again in byte units. Special care is required when using these instructions in cases where two registers are assigned to the same address, or when a bit is directly manipulated for a port or a register containing a write-only bit, because this may rewrite data of a bit other than the bit to be manipulated.

(1) Bit manipulation for two registers assigned to the same address

Example 1: Bit manipulation for the timer load register and timer counter

Figure 2.13 shows an example of a timer in which two timer registers are assigned to the same address. When a bit-manipulation instruction accesses the timer load register and timer counter of a reloadable timer, since these two registers share the same address, the following operations takes place.

- 1. Data is read in byte units.
- 2. The CPU sets or resets the bit to be manipulated with the bit-manipulation instruction.
- 3. The written data is written again in byte units to the timer load register.



The timer is counting, so the value read is not necessarily the same as the value in the timer load register. As a result, bits other than the intended bit in the timer counter may be modified and the modified value may be written to the timer load register.

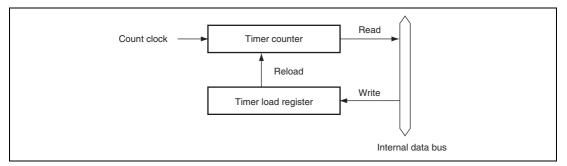


Figure 2.13 Example of Timer Configuration with Two Registers Allocated to Same Address

Example 2: When the BSET instruction is executed for port 5

P57 and P56 are input pins, with a low-level signal input at P57 and a high-level signal input at P56. P55 to P50 are output pins and output low-level signals. An example to output a high-level signal at P50 with a BSET instruction is shown below.

Prior to executing BSET instruction

	P57	P56	P55	P54	P53	P52	P51	P50
Input/output	Input	Input	Output	Output	Output	Output	Output	Output
Pin state	Low level	High level	Low level	Low level	Low level	Low level	Low level	Low level
PCR5	0	0	1	1	1	1	1	1
PDR5	1	0	0	0	0	0	0	0

BSET instruction executed

BSET #0, @PDR5

The BSET instruction is executed for port 5.



 After executing BSET instruction 	on
--	----

	P57	P56	P55	P54	P53	P52	P51	P50
Input/output	Input	Input	Output	Output	Output	Output	Output	Output
Pin state	Low level	High level	Low level	Low level	Low level	Low level	Low level	High level
PCR5	0	0	1	1	1	1	1	1
PDR5	0	1	0	0	0	0	0	1

- Description on operation
- 1. When the BSET instruction is executed, first the CPU reads port 5.

Since P57 and P56 are input pins, the CPU reads the pin states (low-level and high-level input).

P55 to P50 are output pins, so the CPU reads the value in PDR5. In this example PDR5 has a value of H'80, but the value read by the CPU is H'40.

- 2. Next, the CPU sets bit 0 of the read data to 1, changing the PDR5 data to H'41.
- 3. Finally, the CPU writes H'41 to PDR5, completing execution of BSET instruction.

 As a result of the BSET instruction, bit 0 in PDR5 becomes 1, and P50 outputs a high-level signal. However, bits 7 and 6 of PDR5 end up with different values. To prevent this problem,

store a copy of the PDR5 data in a work area in memory. Perform the bit manipulation on the data in the work area, then write this data to PDR5.

Prior to executing BSET instruction

MOV.B #H'80, R0L MOV.B R0L, @RAM0 MOV.B R0L, @PDR5 The PDR5 value (H'80) is written to a work area in memory (RAM0) as well as to PDR5.

	P57	P56	P55	P54	P53	P52	P51	P50
Input/output	Input	Input	Output	Output	Output	Output	Output	Output
Pin state	Low level	High level	Low level	Low level	Low level	Low level	Low level	Low level
PCR5	0	0	1	1	1	1	1	1
PDR5	1	0	0	0	0	0	0	0
RAM0	1	0	0	0	0	0	0	0

BSET instruction executed

BSET #0, @RAM0

The BSET instruction is executed designating the PDR5 work area (RAM0).

After executing BSET instruction

MOV.B @RAM0, R0L MOV.B R0L, @PDR5 The work area (RAM0) value is written to PDR5.

	P57	P56	P55	P54	P53	P52	P51	P50
Input/output	Input	Input	Output	Output	Output	Output	Output	Output
Pin state	Low level	High level	Low level	Low level	Low level	Low level	Low level	High level
PCR5	0	0	1	1	1	1	1	1
PDR5	1	0	0	0	0	0	0	1
RAM0	1	0	0	0	0	0	0	1

(2) Bit Manipulation in a Register Containing a Write-Only Bit

Example 3: BCLR instruction executed designating port 5 control register PCR5

P57 and P56 are input pins, with a low-level signal input at P57 and a high-level signal input at P56. P55 to P50 are output pins that output low-level signals. An example of setting the P50 pin as an input pin by the BCLR instruction is shown below. It is assumed that a high-level signal will be input to this input pin.

Prior to executing BCLR instruction

	P57	P56	P55	P54	P53	P52	P51	P50
Input/output	Input	Input	Output	Output	Output	Output	Output	Output
Pin state	Low level	High level	Low level	Low level	Low level	Low level	Low level	Low level
PCR5	0	0	1	1	1	1	1	1
PDR5	1	0	0	0	0	0	0	0



BCLR instruction executed

BCLR #0, @PCR5

The BCLR instruction is executed for PCR5.

• After executing BCLR instruction

	P57	P56	P55	P54	P53	P52	P51	P50
Input/output	Output	Output	Output	Output	Output	Output	Output	Input
Pin state	Low level	High level	Low level	Low level	Low level	Low level	Low level	High level
PCR5	1	1	1	1	1	1	1	0
PDR5	1	0	0	0	0	0	0	0

- Description on operation
- 1. When the BCLR instruction is executed, first the CPU reads PCR5. Since PCR5 is a write-only register, the CPU reads a value of H'FF, even though the PCR5 value is actually H'3F.
- 2. Next, the CPU clears bit 0 in the read data to 0, changing the data to H'FE.
- 3. Finally, H'FE is written to PCR5 and BCLR instruction execution ends.

 As a result of this operation, bit 0 in PCR5 becomes 0, making P50 an input port. However, bits 7 and 6 in PCR5 change to 1, so that P57 and P56 change from input pins to output pins. To prevent this problem, store a copy of the PDR5 data in a work area in memory and manipulate data of the bit in the work area, then write this data to PDR5.
- Prior to executing BCLR instruction

MOV.B #H'3F, R0L MOV.B R0L, @RAM0 MOV.B R0L, @PCR5

The PCR5 value (H'3F) is written to a work area in memory (RAM0) as well as to PCR5.

	P57	P56	P55	P54	P53	P52	P51	P50
Input/output	Input	Input	Output	Output	Output	Output	Output	Output
Pin state	Low level	High level	Low level	Low level	Low level	Low level	Low level	Low level
PCR5	0	0	1	1	1	1	1	1
PDR5	1	0	0	0	0	0	0	0
RAM0	0	0	1	1	1	1	1	1

BCLR instruction executed

BCLR #0, @RAM0

The BCLR instructions executed for the PCR5 work area (RAM0).

• After executing BCLR instruction

MOV.B @RAM0, R0L MOV.B R0L, @PCR5 The work area (RAM0) value is written to PCR5.

	P57	P56	P55	P54	P53	P52	P51	P50
Input/output	Input	Input	Output	Output	Output	Output	Output	Output
Pin state	Low level	High level	Low level	Low level	Low level	Low level	Low level	High level
PCR5	0	0	1	1	1	1	1	0
PDR5	1	0	0	0	0	0	0	0
RAM0	0	0	1	1	1	1	1	0

Table 2.13 lists the pairs of registers that share identical addresses. Table 2.14 lists the registers that contain write-only bits.

Table 2.13 Registers with Shared Addresses

Register Name	Abbr.	Address
Timer counter and timer load register C	TCC/TLC	H'FFB5
Port data register 1*	PDR1	H'FFD4
Port data register 3*	PDR3	H'FFD6
Port data register 4*	PDR4	H'FFD7
Port data register 5*	PDR5	H'FFD8
Port data register 6*	PDR6	H'FFD9
Port data register 7*	PDR7	H'FFDA
Port data register 8*	PDR8	H'FFDB
Port data register A*	PDRA	H'FFDD

Note: * Port data registers have the same addresses as input pins.



Table 2.14 Registers with Write-Only Bits

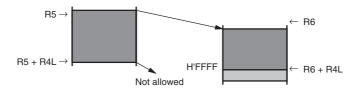
Register Name	Abbr.	Address	
Port control register 1	PCR1	H'FFE4	
Port control register 3	PCR3	H'FFE6	
Port control register 4	PCR4	H'FFE7	
Port control register 5	PCR5	H'FFE8	
Port control register 6	PCR6	H'FFE9	
Port control register 7	PCR7	H'FFEA	
Port control register 8	PCR8	H'FFEB	
Port control register A	PCRA	H'FFED	
Timer control register F	TCRF	H'FFB6	
PWM control register	PWCR	H'FFD0	
PWM data register U	PWDRU	H'FFD1	
PWM data register L	PWDRL	H'FFD2	

2.8.4 Notes on Use of the EEPMOV Instruction

• The EEPMOV instruction is a block data transfer instruction. It moves the number of bytes specified by R4L from the address specified by R5 to the address specified by R6.



When setting R4L and R6, make sure that the final destination address (R6 + R4L) does not
exceed H'FFFF. The value in R6 must not change from H'FFFF to H'0000 during execution of
the instruction.



Section 3 Exception Handling

3.1 Overview

Exception handling is performed in this LSI when a reset or interrupt occurs. Table 3.1 shows the priorities of these two types of exception handling.

Table 3.1 Exception Handling Types and Priorities

Priority	Exception Source	Time of Start of Exception Handling
High	Reset	Exception handling starts as soon as the reset state is cleared
Low	Interrupt	When an interrupt is requested, exception handling starts after execution of the present instruction or the exception handling in progress is completed

3.2 Reset

3.2.1 Overview

A reset is the highest-priority exception. The internal state of the CPU and the registers of the onchip peripheral modules are initialized.

3.2.2 Reset Sequence

As soon as the RES pin goes low, all processing is stopped and the chip enters the reset state.

To make sure the chip is reset properly, observe the following precautions.

- At power on: Hold the RES pin low until the clock pulse generator output stabilizes.
- Resetting during operation: Hold the RES pin low for at least 10 system clock cycles.

Reset exception handling takes place as follows.

- The CPU internal state and the registers of on-chip peripheral modules are initialized, with the I bit of the condition code register (CCR) set to 1.
- The PC is loaded from the reset exception handling vector address (H'0000 to H'0001), after which the program starts executing from the address indicated in PC.

When system power is turned on or off, the \overline{RES} pin should be held low.

Figure 3.1 shows the reset sequence starting from \overline{RES} input.

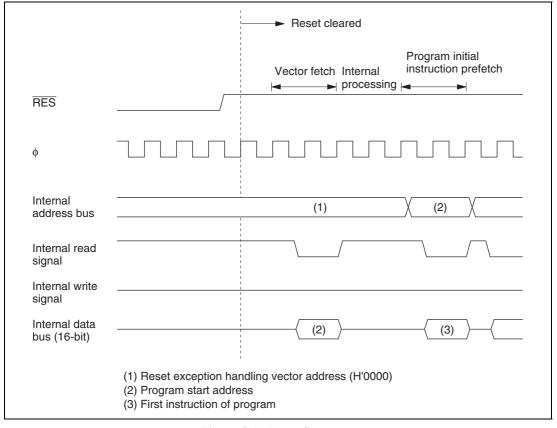


Figure 3.1 Reset Sequence

3.2.3 Interrupt Immediately after Reset

After a reset, if an interrupt were to be accepted before the stack pointer (SP: R7) was initialized, PC and CCR would not be pushed onto the stack correctly, resulting in program runaway. To prevent this, immediately after reset exception handling all interrupts are masked. For this reason, the initial program instruction is always executed immediately after a reset. This instruction should initialize the stack pointer (e.g. MOV.W #xx: 16, SP).

3.3 Interrupts

3.3.1 Overview

The interrupt sources include 13 external interrupts (IRQ $_4$ to IRQ $_6$, WKP $_7$ to WKP $_0$) and 23 internal interrupts from on-chip peripheral modules. Table 3.2 shows the interrupt sources, their priorities, and their vector addresses. When more than one interrupt is requested, the interrupt with the highest priority is processed.

The interrupts have the following features:

- Internal and external interrupts can be masked by the I bit in CCR. When the I bit is set to 1, interrupt request flags can be set but the interrupts are not accepted.
- IRQ₄ to IRQ₀ and WKP₇ to WKP₀ can be set to either rising edge sensing or falling edge sensing.

Table 3.2 Interrupt Sources and Their Priorities

Interrupt Source	Interrupt	Vector Number	Vector Address	Priority
RES	Reset	0	H'0000 to H'0001	High
Watchdog timer				_ 1
ĪRQ₀	$IRQ_{_{0}}$	4	H'0008 to H'0009	_
ĪRQ ₁	IRQ,	5	H'000A to H'000B	
ĪRQ ₂	IRQ ₂	6	H'000C to H'000D	
ĪRQ₃	IRQ ₃	7	H'000E to H'000F	_
ĪRQ₄	IRQ ₄	8	H'0010 to H'0011	_
WKP ₀ WKP ₁ WKP ₂ WKP ₃ WKP ₄ WKP ₅ WKP ₆	WKP ₀ WKP ₁ WKP ₂ WKP ₃ WKP ₄ WKP ₅ WKP ₆	9	H'0012 to H'0013	
WKP,	WKP ₇			_
Timer A	Timer A overflow	11	H'0016 to H'0017	_
Asynchronous counter	Asynchronous counter overflow	12	H'0018 to H'0019	_
Timer C	Timer C overflow or underflow	13	H'001A to H'001B	_
Timer FL	Timer FL compare match Timer FL overflow	14	H'001C to H'001D	
Timer FH	Timer FH compare match Timer FH overflow	15	H'001E to H'001F	_
Timer G	Timer G input capture Timer G overflow	16	H'0020 to H'0021	_
SCI3-1	SCI3-1 transmit end SCI3-1 transmit data empty SCI3-1 receive data full SCI3-1 overrrun error SCI3-1 framing error SCI3-1 parity error	17	H'0022 to H'0023	
SCI3-2	SCI3-2 transmit end SCI3-2 transmit data empty SCI3-2 receive data full SCI3-2 overrun error SCI3-2 framing error SCI3-2 parity error	18	H'0024 to H'0025	
A/D	A/D conversion end	19	H'0026 to H'0027	_
(SLEEP instruction executed)	Direct transfer	20	H'0028 to H'0029	Low

Note: Vector addresses H'0002 to H'0007 and H'0014 to H'0015 are reserved and cannot be used.



3.3.2 Interrupt Control Registers

Table 3.3 lists the registers that control interrupts.

Table 3.3 Interrupt Control Registers

Name	Abbreviation	R/W	Initial Value	Address
IRQ edge select register	IEGR	R/W	H'E0	H'FFF2
Interrupt enable register 1	IENR1	R/W	H'00	H'FFF3
Interrupt enable register 2	IENR2	R/W	H'00	H'FFF4
Interrupt request register 1	IRR1	R/W*	H'20	H'FFF6
Interrupt request register 2	IRR2	R/W*	H'00	H'FFF7
Wakeup interrupt request register	IWPR	R/W*	H'00	H'FFF9
Wakeup edge select register	WEGR	R/W	H'00	H'FF90

Note: * Write is enabled only for writing of 0 to clear a flag.

(1) IRQ Edge Select Register (IEGR)

Bit	7	6	5	4	3	2	1	0
	_	_	_	IEG4	IEG3	IEG2	IEG1	IEG0
Initial value	1	1	1	0	0	0	0	0
Read/Write	_	_	_	R/W	R/W	R/W	R/W	R/W

IEGR is an 8-bit read/write register used to designate whether pins \overline{IRQ}_4 to \overline{IRQ}_0 are set to rising edge sensing or falling edge sensing.

Bits 7 to 5: Reserved bits

Bits 7 to 5 are reserved: they are always read as 1 and cannot be modified.

Bit 4: IRQ₄ edge select (IEG4)

Bit 4 selects the input sensing of the \overline{IRQ}_4 pin and \overline{ADTRG} pin.

Bit 4 IEG4	Description	
0	Falling edge of IRQ and ADTRG pin input is detected	(initial value)
1	Rising edge of $\overline{\text{IRQ}}_4$ and $\overline{\text{ADTRG}}$ pin input is detected	

Bit 3: IRQ₃ edge select (IEG3)

Bit 3 selects the input sensing of the \overline{IRQ}_3 pin and TMIF pin.

Bit 3 IEG3	Description	
0	Falling edge of $\overline{\text{IRQ}}_{_3}$ and TMIF pin input is detected	(initial value)
1	Rising edge of $\overline{\text{IRQ}}_{_3}$ and TMIF pin input is detected	

Bit 2: IRQ₂ edge select (IEG2)

Bit 2 selects the input sensing of pin \overline{IRQ}_2 .

Bit 2 IEG2	Description	
0	Falling edge of $\overline{\text{IRQ}}_{\scriptscriptstyle 2}$ pin input is detected	(initial value)
1	Rising edge of \overline{IRQ}_2 pin input is detected	

Bit 1: IRQ₁ edge select (IEG1)

Bit 1 selects the input sensing of the \overline{IRQ}_1 pin and TMIC pin.

Bit 1		
IEG1	Description	
0	Falling edge of $\overline{\text{IRQ}}_{_1}$ and TMIC pin input is detected	(initial value)
1	Rising edge of IRQ, and TMIC pin input is detected	

Bit 0: IRQ₀ edge select (IEG0)

Bit 0 selects the input sensing of pin \overline{IRQ}_0 .

Bit 0 IEG0	Description	
0	Falling edge of $\overline{\text{IRQ}}_{\scriptscriptstyle 0}$ pin input is detected	(initial value)
1	Rising edge of $\overline{\text{IRQ}}_{_{0}}$ pin input is detected	

(2) Interrupt Enable Register 1 (IENR1)

Bit	7	6	5	4	3	2	1	0
	IENTA	_	IENWP	IEN4	IEN3	IEN2	IEN1	IEN0
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W

IENR1 is an 8-bit read/write register that enables or disables interrupt requests.

Bit 7: Timer A interrupt enable (IENTA)

Bit 7 enables or disables timer A overflow interrupt requests.

Bit 7		
IENTA	Description	
0	Disables timer A interrupt requests	(initial value)
1	Enables timer A interrupt requests	

Bit 6: Reserved bit

Bit 6 is a readable/writable reserved bit. It is initialized to 0 by a reset.

Bit 5: Wakeup interrupt enable (IENWP)

Bit 5 enables or disables WKP $_{\scriptscriptstyle 7}$ to WKP $_{\scriptscriptstyle 0}$ interrupt requests.

Bit 5 IENWP	Description	
0	Disables $\overline{\text{WKP}}_{_7}$ to $\overline{\text{WKP}}_{_0}$ interrupt requests	(initial value)
1	Enables $\overline{\text{WKP}}_{_{7}}$ to $\overline{\text{WKP}}_{_{0}}$ interrupt requests	

Bits 4 to 0: IRQ₄ to IRQ₀ interrupt enable (IEN4 to IEN0)

Bits 4 to 0 enable or disable IRQ₄ to IRQ₀ interrupt requests.

Bit n IENn	Description	
0	Disables interrupt requests from pin IRQn	(initial value)
1	Enables interrupt requests from pin IRQn	
		(n = 4 to 0)

(3) Interrupt Enable Register 2 (IENR2)

Bit	7	6	5	4	3	2	1	0
	IENDT	IENAD	_	IENTG	IENTFH	IENTFL	IENTC	IENEC
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W

IENR2 is an 8-bit read/write register that enables or disables interrupt requests.

Bit 7: Direct transfer interrupt enable (IENDT)

Bit 7 enables or disables direct transfer interrupt requests.

Bit 7		
IENDT	Description	
0	Disables direct transfer interrupt requests	(initial value)
1	Enables direct transfer interrupt requests	

Bit 6: A/D converter interrupt enable (IENAD)

Bit 6 enables or disables A/D converter interrupt requests.

Bit 6 IENAD	Description	
0	Disables A/D converter interrupt requests	(initial value)
1	Enables A/D converter interrupt requests	

Bit 5: Reserved bit

Bit 5 is a readable/writable reserved bit. It is initialized to 0 by a reset.

Bit 4: Timer G interrupt enable (IENTG)

Bit 4 enables or disables timer G input capture or overflow interrupt requests.

Bit 4 IENTG	Description	
0	Disables timer G interrupt requests	(initial value)
1	Enables timer G interrupt requests	

Bit 3: Timer FH interrupt enable (IENTFH)

Bit 3 enables or disables timer FH compare match and overflow interrupt requests.

Bit 3		
IENTFH	Description	
0	Disables timer FH interrupt requests	(initial value)
1	Enables timer FH interrupt requests	

Bit 2: Timer FL interrupt enable (IENTFL)

Bit 2 enables or disables timer FL compare match and overflow interrupt requests.

Bit 2	Description	
IENTFL	Description	
0	Disables timer FL interrupt requests	(initial value)
1	Enables timer FL interrupt requests	

Bit 1: Timer C interrupt enable (IENTC)

Bit 1 enables or disables timer C overflow and underflow interrupt requests.

Bit 1 IENTC	Description	
0	Disables timer C interrupt requests	(initial value)
1	Enables timer C interrupt requests	

Bit 0: Asynchronous event counter interrupt enable (IENEC)

Bit 0 enables or disables asynchronous event counter interrupt requests.

Bit 0 IENEC	Description	
0	Disables asynchronous event counter interrupt requests	(initial value)
1	Enables asynchronous event counter interrupt requests	

For details of SCI3-1 and SCI3-2 interrupt control, see section 10.2.6, Serial Control Register 3 (SCR3).

(4) Interrupt Request Register 1 (IRR1)

Bit	7	6	5	4	3	2	1	0
	IRRTA	_	_	IRRI4	IRRI3	IRRI2	IRRI1	IRRI0
Initial value	0	0	1	0	0	0	0	0
Read/Write	R/W*	R/W*	_	R/W*	R/W*	R/W*	R/W*	R/W*

Note: * Only a write of 0 for flag clearing is possible

IRR1 is an 8-bit read/write register, in which a corresponding flag is set to 1 when a timer A or IRQ₄ to IRQ₀ interrupt is requested. The flags are not cleared automatically when an interrupt is accepted. It is necessary to write 0 to clear each flag.

Bit 7: Timer A interrupt request flag (IRRTA)

Bit 7 IRRTA	Description	
0	Clearing condition:	(initial value)
	When IRRTA = 1, it is cleared by writing 0	
1	Setting condition:	
	When the timer A counter value overflows from H'FF to H'00	

Bit 6: Reserved bit

Bit 6 is a readable/writable reserved bit. It is initialized to 0 by a reset.



Bit 5: Reserved bit

Bit 5 is reserved; it is always read as 1 and cannot be modified.

Bits 4 to 0: IRQ₄ to IRQ₀ interrupt request flags (IRRI4 to IRRI0)

Bit n	
IRRIn	Description
0	Clearing condition: (initial value)
	When IRRIn = 1, it is cleared by writing 0
1	Setting condition:
	When pin $\overline{\mbox{IRQ}}\mbox{n}$ is designated for interrupt input and the designated signal edge is input
	(n = 4 to 0)

(5) Interrupt Request Register 2 (IRR2)

Bit	7	6	5	4	3	2	1	0
	IRRDT	IRRAD	_	IRRTG	IRRTFH	IRRTFL	IRRTC	IRREC
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W*	R/W*	R/W	R/W*	R/W*	R/W *	R/W *	R/W *

Note: * Only a write of 0 for flag clearing is possible

IRR2 is an 8-bit read/write register, in which a corresponding flag is set to 1 when a direct transfer, A/D converter, Timer G, Timer FH, Timer FC, or Timer C interrupt is requested. The flags are not cleared automatically when an interrupt is accepted. It is necessary to write 0 to clear each flag.

Bit 7: Direct transfer interrupt request flag (IRRDT)

Bit 7 IRRDT	Description
0	Clearing condition: (initial value
	When IRRDT = 1, it is cleared by writing 0
1	Setting condition:
	When a direct transfer is made by executing a SLEEP instruction while $DTON = 1$ i $SYSCR2$

Bit 6: A/D converter interrupt request flag (IRRAD)

Bit 6 IRRAD	Description	
0	Clearing condition:	(initial value)
	When IRRAD = 1, it is cleared by writing 0	
1	Setting condition:	
	When A/D conversion is completed and ADSF is cleared to 0 in ADSR	

Bit 5: Reserved bit

Bit 5 is a readable/writable reserved bit. It is initialized to 0 by a reset.

Bit 4: Timer G interrupt request flag (IRRTG)

Bit 4 IRRTG	Description	
0	Clearing condition:	(initial value)
	When IRRTG = 1, it is cleared by writing 0	
1	Setting condition:	
	When the TMIG pin is designated for TMIG input and the designated input, or when TCG overflows while OVIE is set to 1 in TMG	signal edge is

Bit 3: Timer FH interrupt request flag (IRRTFH)

Bit 3 IRRTFH	Description	
0	Clearing condition:	(initial value)
	When IRRTFH = 1, it is cleared by writing 0	
1	Setting condition:	
	When TCFH and OCRFH match in 8-bit timer mode, or when TCF (T and OCRF (OCRFL, OCRFH) match in 16-bit timer mode	CFL, TCFH)

Bit 2: Timer FL interrupt request flag (IRRTFL)

Bit 2 IRRTFL	Description	
0	Clearing condition:	(initial value)
	When IRRTFL= 1, it is cleared by writing 0	
1	Setting condition:	
	When TCFL and OCRFL match in 8-bit timer mode	

Bit 1: Timer C interrupt request flag (IRRTC)

Bit 1 IRRTC	Description	
0	Clearing condition:	(initial value)
	When IRRTC= 1, it is cleared by writing 0	
1	Setting condition:	
	When the timer C counter value overflows (from H'FF to H'00) or under (from H'00 to H'FF) $$	erflows

Bit 0: Asynchronous event counter interrupt request flag (IRREC)

Bit 0 IRREC	Description	
0	Clearing condition:	(initial value)
	When IRREC = 1, it is cleared by writing 0	
1	Setting condition:	
	When ECH overflows in 16-bit counter mode, or ECH or ECL ove counter mode	erflows in 8-bit

(6) Wakeup Interrupt Request Register (IWPR)

Bit	7	6	5	4	3	2	1	0
	IWPF7	IWPF6	IWPF5	IWPF4	IWPF3	IWPF2	IWPF1	IWPF0
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W*	R/W*	R/W*	R/W*	R/W*	R/W *	R/W *	R/W*

Note: * Only a write of 0 for flag clearing is possible

IWPR is an 8-bit read/write register containing wakeup interrupt request flags. When one of pins \overline{WKP}_7 to \overline{WKP}_0 is designated for wakeup input and a rising or falling edge is input at that pin, the corresponding flag in IWPR is set to 1. A flag is not cleared automatically when the corresponding interrupt is accepted. Flags must be cleared by writing 0.

Bits 7 to 0: Wakeup interrupt request flags (IWPF7 to IWPF0)

Bit n IWPFn	Description
0	Clearing condition: (initial value
	When IWPFn= 1, it is cleared by writing 0
1	Setting condition:
	When pin $\overline{\text{WKP}_{_{\scriptscriptstyle n}}}$ is designated for wakeup input and a rising or falling edge is input at that pin
	(n = 7 to 0

(7) Wakeup Edge Select Register (WEGR)

Bit	7	6	5	4	3	2	1	0
	WKEGS7	WKEGS6	WKEGS5	WKEGS4	WKEGS3	WKEGS2	WKEGS1	WKEGS0
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W							

WEGR is an 8-bit read/write register that specifies rising or falling edge sensing for pins WKPn.

WEGR is initialized to H'00 by a reset.



Bit n: WKPn edge select (WKEGSn)

Bit n selects $\overline{WKP}n$ pin input sensing.

Bit n WKEGSn	Description	
0	WKPn pin falling edge detected	(initial value)
1	WKPn pin rising edge detected	
		(n = 7 to 0)

3.3.3 External Interrupts

There are 13 external interrupts: IRQ₄ to IRQ₀ and WKP₇ to WKP₀.

(1) Interrupts WKP₇ to WKP₀

Interrupts WKP₇ to WKP₀ are requested by either rising or falling edge input to pins \overline{WKP}_7 to \overline{WKP}_0 . When these pins are designated as pins \overline{WKP}_7 to \overline{WKP}_0 in port mode register 5 and a rising or falling edge is input, the corresponding bit in IWPR is set to 1, requesting an interrupt. Recognition of wakeup interrupt requests can be disabled by clearing the IENWP bit to 0 in IENR1. These interrupts can all be masked by setting the I bit to 1 in CCR.

When WKP₇ to WKP₀ interrupt exception handling is initiated, the I bit is set to 1 in CCR. Vector number 9 is assigned to interrupts WKP₇ to WKP₀. All eight interrupt sources have the same vector number, so the interrupt-handling routine must discriminate the interrupt source.

(2) Interrupts IRQ₄ to IRQ₀

Interrupts IRQ₄ to IRQ₀ are requested by input signals to pins $\overline{\text{IRQ}}_4$ to $\overline{\text{IRQ}}_0$. These interrupts are detected by either rising edge sensing or falling edge sensing, depending on the settings of bits IEG4 to IEG0 in IEGR.

When these pins are designated as pins \overline{IRQ}_4 to \overline{IRQ}_0 in port mode register 3 and 1 and the designated edge is input, the corresponding bit in IRR1 is set to 1, requesting an interrupt. Recognition of these interrupt requests can be disabled individually by clearing bits IEN4 to IEN0 to 0 in IENR1. These interrupts can all be masked by setting the I bit to 1 in CCR.

When IRQ_4 to IRQ_0 interrupt exception handling is initiated, the I bit is set to 1 in CCR. Vector numbers 8 to 4 are assigned to interrupts IRQ_4 to IRQ_0 . The order of priority is from IRQ_0 (high) to IRQ_4 (low). Table 3.2 gives details.

3.3.4 Internal Interrupts

There are 23 internal interrupts that can be requested by the on-chip peripheral modules. When a peripheral module requests an interrupt, the corresponding bit in IRR1 or IRR2 is set to 1. Recognition of individual interrupt requests can be disabled by clearing the corresponding bit in IENR1 or IENR2. All these interrupts can be masked by setting the I bit to 1 in CCR. When internal interrupt handling is initiated, the I bit is set to 1 in CCR. Vector numbers from 20 to 11 are assigned to these interrupts. Table 3.2 shows the order of priority of interrupts from on-chip peripheral modules.

3.3.5 Interrupt Operations

Interrupts are controlled by an interrupt controller. Figure 3.2 shows a block diagram of the interrupt controller. Figure 3.3 shows the flow up to interrupt acceptance.

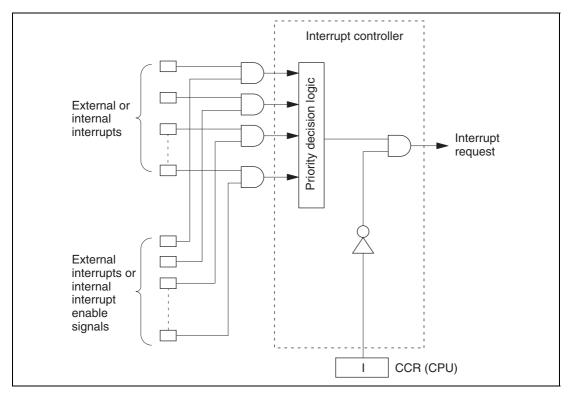


Figure 3.2 Block Diagram of Interrupt Controller

Interrupt operation is described as follows.

- When an interrupt condition is met while the interrupt enable register bit is set to 1, an interrupt request signal is sent to the interrupt controller.
- When the interrupt controller receives an interrupt request, it sets the interrupt request flag.
- From among the interrupts with interrupt request flags set to 1, the interrupt controller selects the interrupt request with the highest priority and holds the others pending. (Refer to table 3.2 for a list of interrupt priorities.)
- The interrupt controller checks the I bit of CCR. If the I bit is 0, the selected interrupt request is accepted; if the I bit is 1, the interrupt request is held pending.
- If the interrupt is accepted, after processing of the current instruction is completed, both PC and CCR are pushed onto the stack. The state of the stack at this time is shown in figure 3.4.
 The PC value pushed onto the stack is the address of the first instruction to be executed upon return from interrupt handling.
- The I bit of CCR is set to 1, masking further interrupts.
- The vector address corresponding to the accepted interrupt is generated, and the interrupt handling routine located at the address indicated by the contents of the vector address is executed.
- Notes: 1. When disabling interrupts by clearing bits in an interrupt enable register, or when clearing bits in an interrupt request register, always do so while interrupts are masked (I = 1).
 - 2. If the above clear operations are performed while I = 0, and as a result a conflict arises between the clear instruction and an interrupt request, exception processing for the interrupt will be executed after the clear instruction has been executed.

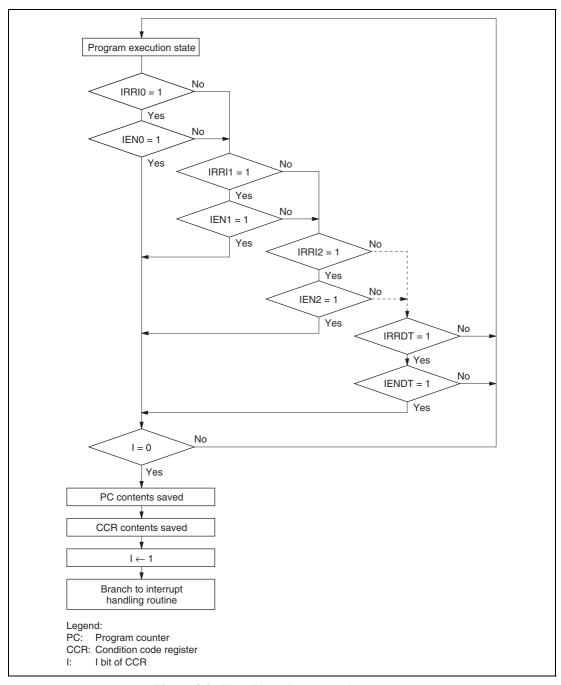


Figure 3.3 Flow Up to Interrupt Acceptance

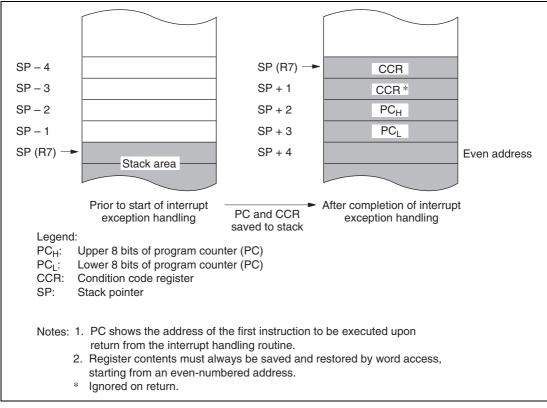


Figure 3.4 Stack State after Completion of Interrupt Exception Handling

Figure 3.5 shows a typical interrupt sequence.

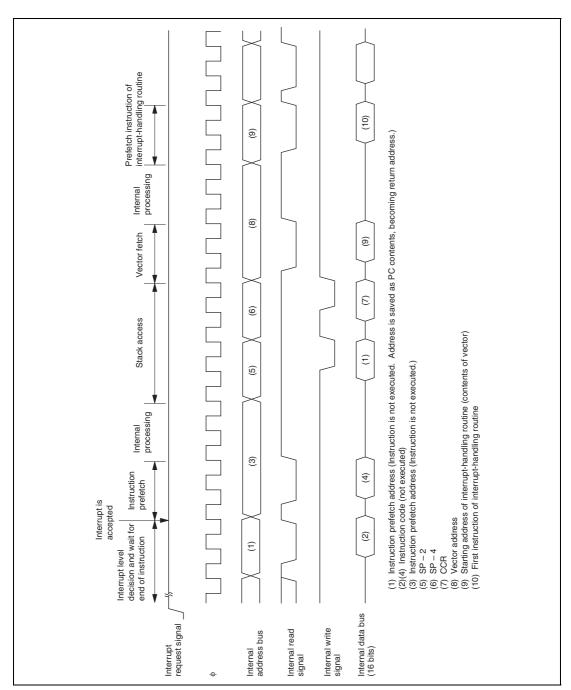


Figure 3.5 Interrupt Sequence

3.3.6 Interrupt Response Time

Table 3.4 shows the number of wait states after an interrupt request flag is set until the first instruction of the interrupt handler is executed.

Table 3.4 Interrupt Wait States

Item	States	Total
Waiting time for completion of executing instruction*	1 to 13	15 to 27
Saving of PC and CCR to stack	4	
Vector fetch	2	
Instruction fetch	4	
Internal processing	4	

Note: * Not including EEPMOV instruction.

3.4 Application Notes

3.4.1 Notes on Stack Area Use

When word data is accessed in the LSI, the least significant bit of the address is regarded as 0. Access to the stack always takes place in word size, so the stack pointer (SP: R7) should never indicate an odd address. Use PUSH Rn (MOV.W Rn, @-SP) or POP Rn (MOV.W @SP+, Rn) to save or restore register values.

Setting an odd address in SP may cause a program to crash. An example is shown in figure 3.6.

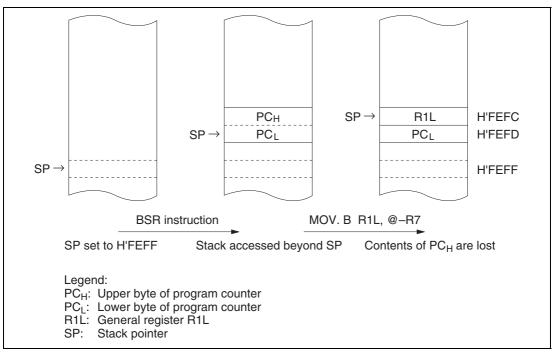


Figure 3.6 Operation when Odd Address is Set in SP

When CCR contents are saved to the stack during interrupt exception handling or restored when RTE is executed, this also takes place in word size. Both the upper and lower bytes of word data are saved to the stack; on return, the even address contents are restored to CCR while the odd address contents are ignored.

3.4.2 Notes on Rewriting Port Mode Registers

When a port mode register is rewritten to switch the functions of external interrupt pins, the following points should be observed.

When an external interrupt pin function is switched by rewriting the port mode register that controls pins \overline{IRQ}_4 to \overline{IRQ}_0 , \overline{WKP}_7 to \overline{WKP}_0 , the interrupt request flag may be set to 1 at the time the pin function is switched, even if no valid interrupt is input at the pin. Be sure to clear the interrupt request flag to 0 after switching pin functions. Table 3.5 shows the conditions under which interrupt request flags are set to 1 in this way.

Table 3.5 Conditions Under which Interrupt Request Flag is Set to 1

Interrup	ot Request	
Flags S	et to 1	Conditions
IRR1	IRRI4	When PMR1 bit IRQ4 is changed from 0 to 1 while pin \overline{IRQ}_4 is low and IEGR bit IEG4 = 0.
		When PMR1 bit IRQ4 is changed from 1 to 0 while pin $\overline{\text{IRQ}}_4$ is low and IEGR bit IEG4 = 1.
	IRRI3	When PMR1 bit IRQ3 is changed from 0 to 1 while pin \overline{IRQ}_3 is low and IEGR bit IEG3 = 0.
		When PMR1 bit IRQ3 is changed from 1 to 0 while pin $\overline{\text{IRQ}}_3$ is low and IEGR bit IEG3 = 1.
	IRRI2	When PMR1 bit IRQ2 is changed from 0 to 1 while pin \overline{IRQ}_2 is low and IEGR bit IEG2 = 0.
		When PMR1 bit IRQ2 is changed from 1 to 0 while pin $\overline{\text{IRQ}}_2$ is low and IEGR bit IEG2 = 1.
	IRRI1	When PMR1 bit IRQ1 is changed from 0 to 1 while pin \overline{IRQ}_1 is low and IEGR bit IEG1 = 0.
		When PMR1 bit IRQ1 is changed from 1 to 0 while pin \overline{IRQ}_1 is low and IEGR bit IEG1 = 1.
	IRRI0	When PMR3 bit IRQ0 is changed from 0 to 1 while pin $\overline{\text{IRQ}}_{\circ}$ is low and IEGR bit IEG0 = 0.
		When PMR3 bit IRQ0 is changed from 1 to 0 while pin $\overline{\text{IRQ}}_0$ is low and IEGR bit IEG0 = 1.
IWPR	IWPF7	When PMR5 bit WKP7 is changed from 0 to 1 while pin $\overline{\text{WKP}}_{7}$ is low.
	IWPF6	When PMR5 bit WKP6 is changed from 0 to 1 while pin $\overline{\text{WKP}}_{\scriptscriptstyle{6}}$ is low.
	IWPF5	When PMR5 bit WKP5 is changed from 0 to 1 while pin $\overline{\text{WKP}}_{\scriptscriptstyle{5}}$ is low.
	IWPF4	When PMR5 bit WKP4 is changed from 0 to 1 while pin $\overline{\text{WKP}}_{_4}$ is low.
	IWPF3	When PMR5 bit WKP3 is changed from 0 to 1 while pin $\overline{\text{WKP}}_3$ is low.
	IWPF2	When PMR5 bit WKP2 is changed from 0 to 1 while pin $\overline{\text{WKP}}_2$ is low.
	IWPF1	When PMR5 bit WKP1 is changed from 0 to 1 while pin WKP, is low.
	IWPF0	When PMR5 bit WKP0 is changed from 0 to 1 while pin $\overline{\text{WKP}}_{0}$ is low.

Figure 3.7 shows the procedure for setting a bit in a port mode register and clearing the interrupt request flag.

When switching a pin function, mask the interrupt before setting the bit in the port mode register. After accessing the port mode register, execute at least one instruction (e.g., NOP), then clear the interrupt request flag from 1 to 0. If the instruction to clear the flag is executed immediately after the port mode register access without executing an intervening instruction, the flag will not be cleared.

An alternative method is to avoid the setting of interrupt request flags when pin functions are switched by keeping the pins at the high level so that the conditions in table 3.5 do not occur.

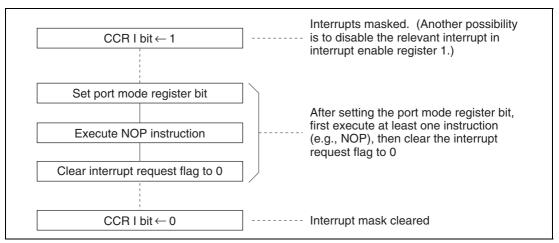


Figure 3.7 Port Mode Register Setting and Interrupt Request Flag Clearing Procedure

3.4.3 Method for Clearing Interrupt Request Flags

Use the recommended method, given below when clearing the flags of interrupt request registers (IRR1, IRR2, IWPR).

Recommended method

Use a single instruction to clear flags. The bit control instruction and byte-size data transfer instruction can be used. Two examples of program code for clearing IRRI1 (bit 1 of IRR1) are given below.

```
BCLR #1, @IRR1:8
MOV.B R1L, @IRR1:8 (set the value of R1L to B'11111101)
```

• Example of a malfunction

When flags are cleared with multiple instructions, other flags might be cleared during execution of the instructions, even though they are currently set, and this will cause a malfunction.

Here is an example in which IRRI0 is cleared and disabled in the process of clearing IRRI1 (bit 1 of IRR1).

```
MOV.B @IRR1:8,R1L ...... IRRIO = 0 at this time
AND.B #B'11111101,R1L ..... Here, IRRIO = 1
MOV.B R1L.@IRR1:8 ...... IRRIO is cleared to 0
```

In the above example, it is assumed that an IRQ0 interrupt is generated while the AND.B instruction is executing.

The IRQ0 interrupt is disabled because, although the original objective is clearing IRRI1, IRRI0 is also cleared.

Section 4 Clock Pulse Generators

4.1 Overview

Clock oscillator circuitry (CPG: clock pulse generator) is provided on-chip, including both a system clock pulse generator and a subclock pulse generator. The system clock pulse generator consists of a system clock oscillator and system clock dividers. The subclock pulse generator consists of a subclock oscillator circuit and a subclock divider.

4.1.1 Block Diagram

Figure 4.1 shows a block diagram of the clock pulse generators.

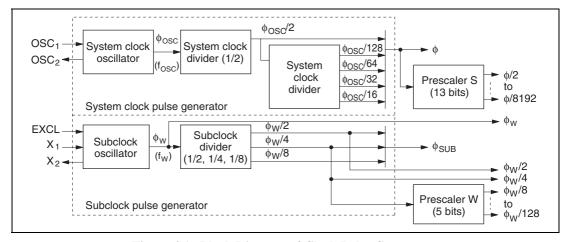


Figure 4.1 Block Diagram of Clock Pulse Generators

4.1.2 System Clock and Subclock

The basic clock signals that drive the CPU and on-chip peripheral modules are ϕ and ϕ_{SUB} . Four of the clock signals have names: ϕ is the system clock, ϕ_{SUB} is the subclock, ϕ_{OSC} is the oscillator clock, and ϕ_{w} is the watch clock.

The clock signals available for use by peripheral modules are $\phi/2$, $\phi/4$, $\phi/8$, $\phi/16$, $\phi/32$, $\phi/64$, $\phi/128$, $\phi/256$, $\phi/512$, $\phi/1024$, $\phi/2048$, $\phi/4096$, $\phi/8192$, ϕ_w , $\phi_w/2$, $\phi_w/4$, $\phi_w/8$, $\phi_w/16$, $\phi_w/32$, $\phi_w/64$, and $\phi_w/128$. The clock requirements differ from one module to another.

4.2 System Clock Generator

Clock pulses can be supplied to the system clock divider either by connecting a crystal or ceramic oscillator, or by providing external clock input.

(1) Connecting a Crystal Oscillator

Figure 4.2 shows a typical method of connecting a crystal oscillator. For information on recommended resonators, see the product AC characteristics listed in section 16, Electrical Characteristics. Please consult with the resonator manufacturer when selecting a resonator model.

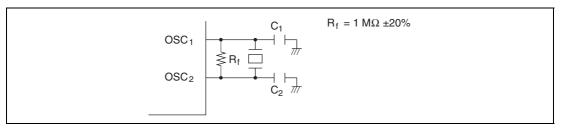


Figure 4.2 Typical Connection to Crystal Oscillator

(2) Connecting a Ceramic Oscillator

Figure 4.3 shows a typical method of connecting a ceramic oscillator. For information on recommended resonators, see the product AC characteristics listed in section 16, Electrical Characteristics. Please consult with the resonator manufacturer when selecting a resonator model.

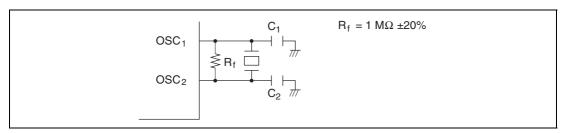


Figure 4.3 Typical Connection to Ceramic Oscillator

(3) Notes on Board Design

When generating clock pulses by connecting a crystal or ceramic oscillator, pay careful attention to the following points.

Avoid running signal lines close to the oscillator circuit, since the oscillator may be adversely affected by induction currents. (See figure 4.4.)

The board should be designed so that the oscillator and load capacitors are located as close as possible to pins OSC₁ and OSC₂.

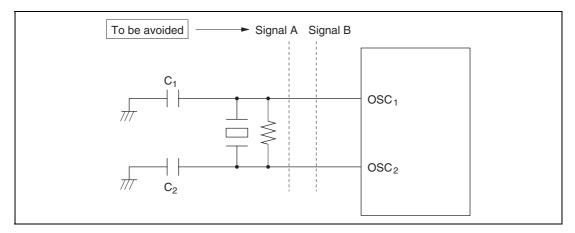


Figure 4.4 Board Design of Oscillator Circuit

(4) External Clock Input Method

Connect an external clock signal to pin OSC₁, and leave pin OSC₂ open. Figure 4.5 shows a typical connection.

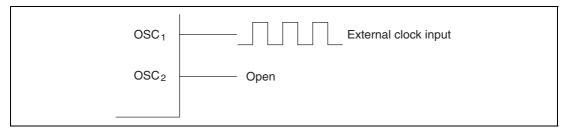


Figure 4.5 External Clock Input (Example)

Frequency	Oscillator Clock (φ _{osc})
Duty cycle	45% to 55%

Note: The circuit parameters above are recommended by the crystal or ceramic oscillator manufacturer.

The circuit parameters are affected by the crystal or ceramic oscillator and floating capacitance when designing the board. When using the oscillator, consult with the crystal or ceramic oscillator manufacturer to determine the circuit parameters.

4.3 Subclock Generator

(1) Connecting a 32.768 kHz/38.4 kHz Crystal Oscillator

Clock pulses can be supplied to the subclock divider by connecting a 32.768 kHz/38.4 kHz crystal oscillator, as shown in figure 4.6. Follow the same precautions as noted under 3. notes on board design for the system clock in 4.2.

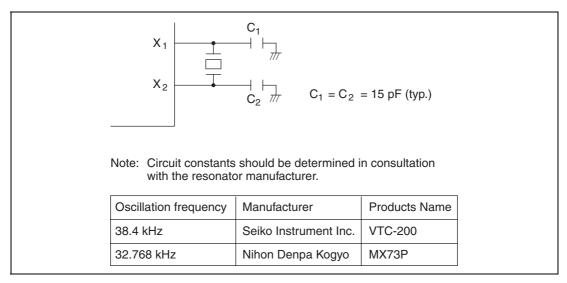


Figure 4.6 Typical Connection to 32.768 kHz/38.4 kHz Crystal Oscillator (Subclock)

Figure 4.7 shows the equivalent circuit of the 32.768 kHz/38.4 kHz crystal oscillator.

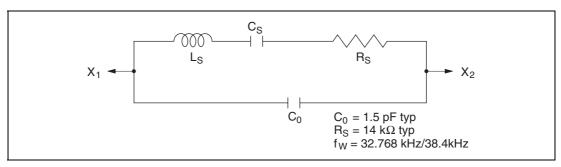


Figure 4.7 Equivalent Circuit of 32.768 kHz/38.4 kHz Crystal Oscillator

(2) Pin Connection when Not Using Subclock

When the subclock is not used, connect pin X_1 to GND and leave pin X_2 open, as shown in figure 4.8.

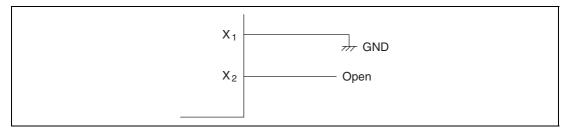


Figure 4.8 Pin Connection when not Using Subclock

(3) External Clock Input

Connect pin X_1 to GND and leave pin X_2 open. Input an external clock to pin EXCL. Set bit EXCL in register PMR2 to 1 to supply the external clock to the internal components of the device. A connection example is shown in figure 4.9.

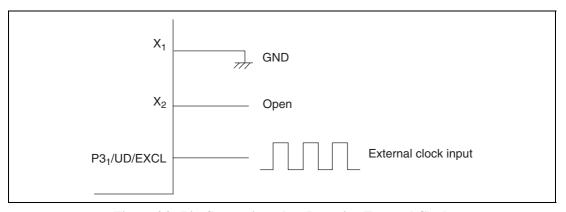


Figure 4.9 Pin Connection when Inputting External Clock

Frequency	Subclock (фw)
Duty	45% to 55%

4.4 Prescalers

This LSI is equipped with two on-chip prescalers having different input clocks (prescaler S and prescaler W). Prescaler S is a 13-bit counter using the system clock (ϕ) as its input clock. Its prescaled outputs provide internal clock signals for on-chip peripheral modules. Prescaler W is a 5-bit counter using a 32.768 kHz or 38.4 kHz signal divided by 4 (ϕ_w /4) as its input clock. Its prescaled outputs are used by timer A as a time base for timekeeping.

(1) Prescaler S (PSS)

Prescaler S is a 13-bit counter using the system clock (ϕ) as its input clock. It is incremented once per clock period.

Prescaler S is initialized to H'0000 by a reset, and starts counting on exit from the reset state.

In standby mode, watch mode, subactive mode, and subsleep mode, the system clock pulse generator stops. Prescaler S also stops and is initialized to H'0000.

The CPU cannot read or write prescaler S.

The output from prescaler S is shared by timer A, timer C, timer F, timer G, SCI3-1, SC3-2, the A/D converter, the LCD controller, the watchdog timer, and the 14-bit PWM. The divider ratio can be set separately for each on-chip peripheral function.

In active (medium-speed) mode the clock input to prescaler S is ϕ osc/16, ϕ osc/32, ϕ osc/64, or ϕ osc/128.

(2) Prescaler W (PSW)

Prescaler W is a 5-bit counter using a 32.768 kHz/38.4 kHz signal divided by 4 (ϕ_w /4) as its input clock.

Prescaler W is initialized to H'00 by a reset, and starts counting on exit from the reset state.

Even in standby mode, watch mode, subactive mode, or subsleep mode, prescaler W continues functioning so long as clock signals are supplied to pins X1 and X2.

Prescaler W can be reset by setting 1s in bits TMA3 and TMA2 of timer mode register A (TMA).

Output from prescaler W can be used to drive timer A, in which case timer A functions as a time base for timekeeping.



4.5 Note on Oscillators

Oscillator characteristics are closely related to board design and should be carefully evaluated by the user, referring to the examples shown in this section. Oscillator circuit constants will differ depending on the oscillator element, stray capacitance in its interconnecting circuit, and other factors. Suitable constants should be determined in consultation with the oscillator element manufacturer. Design the circuit so that the oscillator element never receives voltages exceeding its maximum rating.

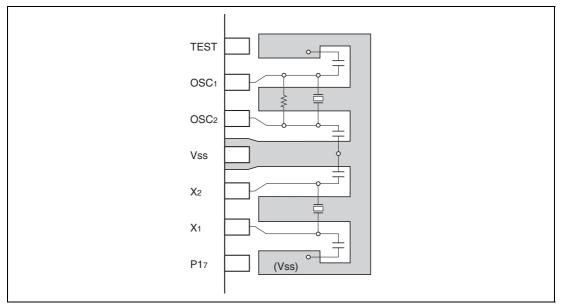


Figure 4.10 Example of Crystal and Ceramic Oscillator Element Arrangement

Figure 4.11 (1) shows an example measuring circuit with the negative resistance suggested by the oscillator manufacturer. Note that if the negative resistance of the circuit is less than that suggested by the oscillator manufacturer, it may be difficult to start the main oscillator.

If it is determined that oscillation is not occurring because the negative resistance is lower than the level suggested by the oscillator manufacturer, the circuit may be modified as shown in figure 4.11 (2) through (4). Which of the modification suggestions to use and the capacitor capacitance should be decided based upon an evaluation of factors such as the negative resistance and the frequency deviation.

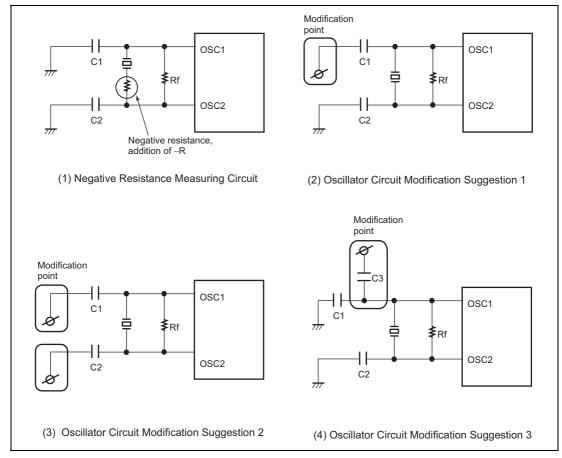


Figure 4.11 Negative Resistance Measurement and Circuit Modification Suggestions

4.5.1 Definition of Oscillation Stabilization Wait Time

Figure 4.12 shows the oscillation waveform (OSC₂), system clock (ϕ), and microcomputer operating mode when a transition is made from standby mode, watch mode, or subactive mode, to active (high-speed/medium-speed) mode, with an oscillator element connected to the system clock oscillator.

As shown in figure 4.12, as the system clock oscillator is halted in standby mode, watch mode, and subactive mode, when a transition is made to active (high-speed/medium-speed) mode, the sum of the following two times (oscillation stabilization time and wait time) is required.

(1) Oscillation Stabilization Time (t_{rc})

The time from the point at which the system clock oscillator oscillation waveform starts to change when an interrupt is generated, until the amplitude of the oscillation waveform increases and the oscillation frequency stabilizes.

(2) Wait Time

The time required for the CPU and peripheral functions to begin operating after the oscillation waveform frequency and system clock have stabilized.

The wait time setting is selected with standby timer select bits 2 to 0 (STS2 to STS0) (bits 6 to 4 in system control register 1 (SYSCR1)).

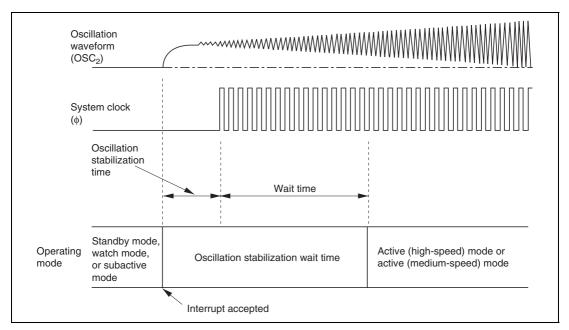


Figure 4.12 Oscillation Stabilization Wait Time

When standby mode, watch mode, or subactive mode is cleared by an interrupt or reset, and a transition is made to active (high-speed/medium-speed) mode, the oscillation waveform begins to change at the point at which the interrupt is accepted. Therefore, when an oscillator element is connected in standby mode, watch mode, or subactive mode, since the system clock oscillator is halted, the time from the point at which this oscillation waveform starts to change until the amplitude of the oscillation waveform increases and the oscillation frequency stabilizes—that is, the oscillation stabilization time—is required.

The oscillation stabilization time in the case of these state transitions is the same as the oscillation stabilization time at power-on (the time from the point at which the power supply voltage reaches the prescribed level until the oscillation stabilizes), specified by "oscillation stabilization time $t_{\Gamma C}$ " in the AC characteristics.

Meanwhile, once the system clock has halted, a wait time of at least 8 states is necessary in order for the CPU and peripheral functions to operate normally.

Thus, the time required from interrupt generation until operation of the CPU and peripheral functions is the sum of the above described oscillation stabilization time and wait time. This total time is called the oscillation stabilization wait time, and is expressed by equation (1) below.

Oscillation stabilization wait time = oscillation stabilization time + wait time =
$$t_{\perp} + (8 \text{ to } 131,072 \text{ states})$$
(1)

Therefore, when a transition is made from standby mode, watch mode, or subactive mode, to active (high-speed/medium-speed) mode, with an oscillator element connected to the system clock oscillator, careful evaluation must be carried out on the installation circuit before deciding on the oscillation stabilization wait time. In particular, since the oscillation stabilization time is affected by installation circuit constants, stray capacitance, and so forth, suitable constants should be determined in consultation with the oscillator element manufacturer.

4.5.2 Notes on Use of Crystal Oscillator Element (Excluding Ceramic Oscillator Element)

When a microcomputer operates, the internal power supply potential fluctuates slightly in synchronization with the system clock. Depending on the individual crystal oscillator element characteristics, the oscillation waveform amplitude may not be sufficiently large immediately after the oscillation stabilization wait time, making the oscillation waveform susceptible to influence by fluctuations in the power supply potential. In this state, the oscillation waveform may be disrupted, leading to an unstable system clock and erroneous operation of the microcomputer.

If erroneous operation occurs, change the setting of standby timer select bits 2 to 0 (STS2 to STS0) (bits 6 to 4 in system control register 1 (SYSCR1)) to give a longer wait time.

For example, if erroneous operation occurs with a wait time setting of 16 states, check the operation with a wait time setting of 8,192 states or more.

If the same kind of erroneous operation occurs after a reset as after a state transition, hold the $\overline{\text{RES}}$ pin low for a longer period.



Section 5 Power-Down Modes

5.1 Overview

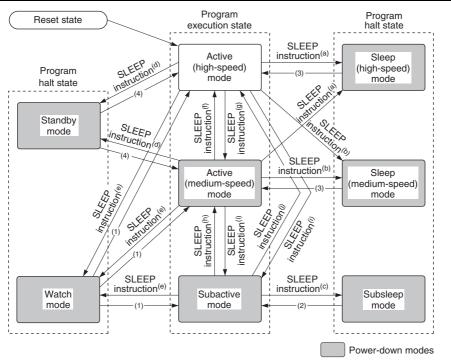
The LSI has nine modes of operation after a reset. These include eight power-down modes, in which power dissipation is significantly reduced. Table 5.1 gives a summary of the nine operating modes.

Table 5.1 Operating Modes

Operating Mode	Description
Active (high-speed) mode	The CPU and all on-chip peripheral functions are operable on the system clock in high-speed operation
Active (medium-speed) mode	The CPU and all on-chip peripheral functions are operable on the system clock in low-speed operation
Subactive mode	The CPU is operable on the subclock in low-speed operation
Sleep (high-speed) mode	The CPU halts. On-chip peripheral functions are operable on the system clock
Sleep (medium-speed) mode	The CPU halts. On-chip peripheral functions operate at a frequency of 1/64, 1/32, 1/16, or 1/8 of the system clock frequency
Subsleep mode	The CPU halts. The time-base function of timer A, timer C, timer G, timer F,WDT, SCl3-1, SCl3-2, AEC, and LCD controller/driver are operable on the subclock
Watch mode	The CPU halts. The time-base function of timer A, timer F, timer G, AEC, and LCD controller/driver are operable on the subclock
Standby mode	The CPU and all on-chip peripheral functions halt
Module standby mode	Individual on-chip peripheral functions specified by software enter standby mode and halt

Of these nine operating modes, all but the active (high-speed) mode are power-down modes. In this section the two active modes (high-speed and medium speed) will be referred to collectively as active mode.

Figure 5.1 shows the transitions among these operation modes. Table 5.2 indicates the internal states in each mode.



Mode Transition Conditions (1)

	LSON	MSON	SSBY	TMA3	DTON
(a)	0	0	0	*	0
(b)	0	1	0	*	0
(c)	1	*	0	1	0
(d)	0	*	1	0	0
(e)	*	*	1	1	0
(f)	0	0	0	*	1
(g)	0	1	0	*	1
(h)	0	1	1	1	1
(i)	1	*	1	1	1
(J)	0	0	1	1	1

Mode Transition Conditions (2)

	Interrupt Sources
(1)	Timer A, Timer F, Timer G interrupt, IRQ ₀ interrupt,
	WKP ₇ to WKP ₀ interrupts
(2)	Timer A, Timer C, Timer F, Timer G, SCI3-1,
	SCI3-2 interrupt, IRQ ₄ to IRQ ₀ interrupts,
	WKP ₇ to WKP ₀ interrupts, AEC
(3)	All interrupts
(4)	IRQ ₁ or IRQ ₀ interrupt, WKP ₇ to WKP ₀ interrupts

* : Don't care

Notes: 1. A transition between different modes cannot be made to occur simply because an interrupt request is generated. Make sure that interrupt handling is performed after the interrupt is accepted.

2. Details on the mode transition conditions are given in the explanations of each mode, in sections 5.2 to 5.9.

Figure 5.1 Mode Transition Diagram

Table 5.2 Internal State in Each Operating Mode

		Activ	e Mode	Sleep	Mode				
Function		High- Speed	Medium- Speed	High- Speed	Medium- Speed	Watch Mode	Subactive Mode	Subsleep Mode	Standby Mode
System cloc	k oscillator	Functions	Functions	Functions	Functions	Halted	Halted	Halted	Halted
Subclock os	scillator	Functions	Functions	Functions	Functions	Functions	Functions	Functions	Functions
CPU	Instructions	Functions	Functions	Halted	Halted	Halted	Functions	Halted	Halted
operations	RAM	_		Retained	Retained	Retained	_	Retained	Retained
	Registers	_							
	I/O ports	_							Retained*1
External	IRQ ₀	Functions	Functions	Functions	Functions	Functions	Functions	Functions	Functions
interrupts	IRQ ₁	_				Retained*6	_		
	IRQ ₂	_							Retained*6
	IRQ ₃	_							
	IRQ ₄	_							
	WKP ₀	Functions	Functions	Functions	Functions	Functions	Functions	Functions	Functions
	WKP ₁	_							
	WKP ₂	_							
	WKP ₃	_							
	WKP ₄	_							
	WKP ₅	_							
	WKP ₆	_							
	WKP ₇	<u> </u>							
Peripheral	Timer A	Functions	Functions	Functions	Functions	Functions*5	Functions*5	Functions*5	Retained
functions	Asynchronous counter	_				Functions*8	Functions	Functions	Functions*8
	Timer C	_				Retained	Functions/ Retained*2	Functions/ Retained*2	Retained
	WDT	_					Functions/ Retained*7	Retained	-
	Timer G, Timer F	_				Functions/ Retained*9	Functions/ Retained*2	Functions/ Retained*2	=
	SCI3-1	_				Reset	Functions/	Functions/	Reset
	SCI3-2	=					Retained*3	Retained*3	-
	PWM	=				Retained	Retained	Retained	Retained
	A/D converter	_				Retained	Retained	Retained	Retained
	LCD	_				Functions/ Retained*4	Functions/ Retained*4	Functions/ Retained*4	Retained

Notes:

- 1. Register contents are retained, but output is high-impedance state.
- 2. Functions if an external clock or the $\phi_W/4$ internal clock is selected; otherwise halted and retained.
- 3. Functions if $\phi_W/2$ is selected as the internal clock; otherwise halted and retained.
- 4. Functions if ϕ_W , $\phi_W/2$ or $\phi_W/4$ is selected as the operating clock; otherwise halted and retained.
- 5. Functions if the timekeeping time-base function is selected.
- 6. External interrupt requests are ignored. Interrupt request register contents are not altered.
- 7. Functions if $\phi_W/32$ is selected as the internal clock; otherwise halted and retained.
- 8. Incrementing is possible, but interrupt generation is not.
- 9. Functions if the $\phi_W/4$ internal clock is selected; otherwise halted and retained.

5.1.1 System Control Registers

The operation mode is selected using the system control registers described in table 5.3.

Table 5.3 System Control Registers

Name	Abbr.	R/W	Initial Value	Address
System control register 1	SYSCR1	R/W	H'07	H'FFF0
System control register 2	SYSCR2	R/W	H'F0	H'FFF1

(1) System Control Register 1 (SYSCR1)

Bit	7	6	5	4	3	2	1	0
	SSBY	STS2	STS1	STS0	LSON	_	MA1	MA0
Initial value	0	0	0	0	0	1	1	1
Read/Write	R/W	R/W	R/W	R/W	R/W	_	R/W	R/W

SYSCR1 is an 8-bit read/write register for control of the power-down modes.

Upon reset, SYSCR1 is initialized to H'07.

Bit 7: Software standby (SSBY)

This bit designates transition to standby mode or watch mode.

SSBY	Description							
0	When a SLEEP instruction is executed in active mode, a transition is made to sleep mode							
	 When a SLEEP instruction is executed in subactive mode, a transition is made to subsleep mode 							
1	When a SLEEP instruction is executed in active mode, a transition is made to standby mode or watch mode							
	 When a SLEEP instruction is executed in subactive mode, a transition is made to watch mode 							

Bits 6 to 4: Standby timer select 2 to 0 (STS2 to STS0)

These bits designate the time the CPU and peripheral modules wait for stable clock operation after exiting from standby mode or watch mode to active mode due to an interrupt. The designation should be made according to the operating frequency so that the waiting time is at least equal to the oscillation settling time.

Bit 6 STS2	Bit 5 STS1	Bit 4 STS0	Description	
0	0	0	Wait time = 8,192 states	(initial value)
0	0	1	Wait time = 16,384 states	
0	1	0	Wait time = 32,768 states	
0	1	1	Wait time = 65,536 states	
1	0	0	Wait time = 131,072 states	
1	0	1	Wait time = 2 states	(External clock mode)
1	1	0	Wait time = 8 states	
1	1	1	Wait time = 16 states	

Note: In the case that external clock is input, set up the "Standby timer select" selection to External clock mode before Mode Transition. Also, do not set up to external clock mode, in the case that it does not use external clock.

Bit 3: Low speed on flag (LSON)

This bit chooses the system clock (ϕ) or subclock (ϕ_{SUB}) as the CPU operating clock when watch mode is cleared. The resulting operation mode depends on the combination of other control bits and interrupt input.

Bit 3 LSON	Description	
0	The CPU operates on the system clock (ϕ)	(initial value)
1	The CPU operates on the subclock (φ _{SUB})	

Bit 2: Reserved bit

Bit 2 is reserved: it is always read as 1 and cannot be modified.

Bits 1 and 0: Active (medium-speed) mode clock select (MA1, MA0)

Bits 1 and 0 choose $\phi_{osc}/128$, $\phi_{osc}/64$, $\phi_{osc}/32$, or $\phi_{osc}/16$ as the operating clock in active (medium-speed) mode and sleep (medium-speed) mode. MA1 and MA0 should be written in active (high-speed) mode or subactive mode.

Bit 1 MA1	Bit 0 MA0	Description	
0	0	φ _{osc} /16	
0	1	φ _{osc} /32	
1	0	φ _{osc} /64	
1	1	φ _{osc} /128	(initial value)

(2) System Control Register 2 (SYSCR2)

Bit	7	6	5	4	3	2	1	0
	_	_	_	NESEL	DTON	MSON	SA1	SA0
Initial value	1	1	1	1	0	0	0	0
Read/Write	_	_	_	R/W	R/W	R/W	R/W	R/W

SYSCR2 is an 8-bit read/write register for power-down mode control.

Bits 7 to 5: Reserved bits

These bits are reserved; they are always read as 1, and cannot be modified.

Bit 4: Noise elimination sampling frequency select (NESEL)

This bit selects the frequency at which the watch clock signal (ϕ_w) generated by the subclock pulse generator is sampled, in relation to the oscillator clock (ϕ_{osc}) generated by the system clock pulse generator. When $\phi_{osc} = 2$ to 16 MHz, clear NESEL to 0.

Bit 4 NESEL	Description	
0	Sampling rate is $\phi_{ m osc}/16$	
1	Sampling rate is $\phi_{ m osc}/4$	(initial value)

Bit 3: Direct transfer on flag (DTON)

This bit designates whether or not to make direct transitions among active (high-speed), active (medium-speed) and subactive mode when a SLEEP instruction is executed. The mode to which the transition is made after the SLEEP instruction is executed depends on a combination of this and other control bits.

Bit 3 DTON	Description					
0	•	When a SLEEP instruction is executed in active mode, (initial value)				
	•	a transition is made to standby mode, watch mode, or sleep mode When a SLEEP instruction is executed in subactive mode, a transition is made				
		to watch mode or subsleep mode				
1	•	When a SLEEP instruction is executed in active (high-speed) mode, a direct transition is made to active (medium-speed) mode if SSBY = 0, MSON = 1, and LSON = 0, or to subactive mode if SSBY = 1, TMA3 = 1, and LSON = 1				
	•	When a SLEEP instruction is executed in active (medium-speed) mode, a direct transition is made to active (high-speed) mode if SSBY = 0, MSON = 0, and LSON = 0, or to subactive mode if SSBY = 1, $TMA3 = 1$, and $TMA3 = 1$				
	•	When a SLEEP instruction is executed in subactive mode, a direct transition is made to active (high-speed) mode if SSBY = 1, TMA3 = 1, LSON = 0, and MSON = 0, or to active (medium-speed) mode if SSBY = 1, TMA3 = 1, LSON = 0, and MSON = 1				

Bit 2: Medium speed on flag (MSON)

After standby, watch, or sleep mode is cleared, this bit selects active (high-speed) or active (medium-speed) mode.

Bit 2 MSON	Description	
0	Operation in active (high-speed) mode	(initial value)
1	Operation in active (medium-speed) mode	

Bits 1 and 0: Subactive mode clock select (SA1 and SA0)

These bits select the CPU clock rate $(\phi_w/2, \phi_w/4, \text{ or } \phi_w/8)$ in subactive mode. SA1 and SA0 cannot be modified in subactive mode.

Bit 1 SA1	Bit 0 SA0	Description	
0	0	φ _w /8	(initial value)
0	1	φ _w /4	
1	*	φ _w /2	

* : Don't care

5.2 Sleep Mode

5.2.1 Transition to Sleep Mode

(1) Transition to Sleep (High-Speed) Mode

The system goes from active mode to sleep (high-speed) mode when a SLEEP instruction is executed while the SSBY and LSON bits in SYSCR1 are cleared to 0, the MSON and DTON bits in SYSCR2 are cleared to 0. In sleep mode CPU operation is halted but the on-chip peripheral functions. CPU register contents are retained.

(2) Transition to Sleep (Medium-Speed) Mode

The system goes from active mode to sleep (medium-speed) mode when a SLEEP instruction is executed while the SSBY and LSON bits in SYSCR1 are cleared to 0, the MSON bit in SYSCR2 is set to 1, and the DTON bit in SYSCR2 is cleared to 0. In sleep (medium-speed) mode, as in sleep (high-speed) mode, CPU operation is halted but the on-chip peripheral functions are operational. The clock frequency in sleep (medium-speed) mode is determined by the MA1 and MA0 bits in SYSCR1. CPU register contents are retained.

Furthermore, it sometimes acts with half state early timing at the time of transition to sleep (medium-speed) mode.



5.2.2 Clearing Sleep Mode

Sleep mode is cleared by any interrupt (timer A, timer C, timer F, timer G, asynchronous counter, IRQ_a to IRQ_o , WKP_a to WKP_o , SCI3-1, SCI3-2, A/D converter, or), or by input at the \overline{RES} pin.

Clearing by interrupt

When an interrupt is requested, sleep mode is cleared and interrupt exception handling starts. A transition is made from sleep (high-speed) mode to active (high-speed) mode, or from sleep (medium-speed) mode to active (medium-speed) mode. Sleep mode is not cleared if the I bit of the condition code register (CCR) is set to 1 or the particular interrupt is disabled in the interrupt enable register.

Interrupt signal and system clock are mutually asynchronous. Synchronization error time in a maximum is $2/\phi$ (s).

Clearing by RES input
 When the RES pin goes low, the CPU goes into the reset state and sleep mode is cleared.

5.2.3 Clock Frequency in Sleep (Medium-Speed) Mode

Operation in sleep (medium-speed) mode is clocked at the frequency designated by the MA1 and MA0 bits in SYSCR1.

5.3 Standby Mode

5.3.1 Transition to Standby Mode

The system goes from active mode to standby mode when a SLEEP instruction is executed while the SSBY bit in SYSCR1 is set to 1, the LSON bit in SYSCR1 is cleared to 0, and bit TMA3 in TMA is cleared to 0. In standby mode the clock pulse generator stops, so the CPU and on-chip peripheral modules stop functioning, but as long as the rated voltage is supplied, the contents of CPU registers, on-chip RAM, and some on-chip peripheral module registers are retained. On-chip RAM contents will be further retained down to a minimum RAM data retention voltage. The I/O ports go to the high-impedance state.

5.3.2 Clearing Standby Mode

Standby mode is cleared by an interrupt (IRQ₁ or IRQ₀), WKP₇ to WKP₀ or by input at the \overline{RES} pin.

• Clearing by interrupt

When an interrupt is requested, the system clock pulse generator starts. After the time set in bits STS2 to STS0 in SYSCR1 has elapsed, a stable system clock signal is supplied to the entire chip, standby mode is cleared, and interrupt exception handling starts. Operation resumes in active (high-speed) mode if MSON = 0 in SYSCR2, or active (medium-speed) mode if MSON = 1. Standby mode is not cleared if the I bit of CCR is set to 1 or the particular interrupt is disabled in the interrupt enable register.

Clearing by RES input

When the \overline{RES} pin goes low, the system clock pulse generator starts. After the pulse generator output has stabilized, if the \overline{RES} pin is driven high, the CPU starts reset exception handling. Since system clock signals are supplied to the entire chip as soon as the system clock pulse generator starts functioning, the \overline{RES} pin should be kept at the low level until the pulse generator output stabilizes.



5.3.3 Oscillator Settling Time after Standby Mode is Cleared

Bits STS2 to STS0 in SYSCR1 should be set as follows.

• When a crystal oscillator is used

The table below gives settings for various operating frequencies. Set bits STS2 to STS0 for a waiting time at least as long as the oscillation settling time.

Table 5.4 Clock Frequency and Settling Time (Times are in ms)

STS2	STS1	STS0	Waiting Time	2 MHz	1 MHz
0	0	0	8,192 states	4.1	8.2
0	0	1	16,384 states	8.2	16.4
0	1	0	32,768 states	16.4	32.8
0	1	1	65,536 states	32.8	65.5
1	0	0	131,072 states	65.5	131.1
1	0	1	2 states (Use prohibited)	0.001	0.002
1	1	0	8 states	0.004	0.008
1	1	1	16 states	0.008	0.016

When an external clock is used
 STS2 = 1, STS1 = 0, and STS0 = 1 should be set. Other values possible use, but CPU sometimes will start operation before waiting time completion.

5.3.4 Standby Mode Transition and Pin States

When a SLEEP instruction is executed in active (high-speed) mode or active (medium-speed) mode while bit SSBY is set to 1 and bit LSON is cleared to 0 in SYSCR1, and bit TMA3 is cleared to 0 in TMA, a transition is made to standby mode. At the same time, pins go to the high-impedance state (except pins for which the pull-up MOS is designated as on). Figure 5.2 shows the timing in this case.

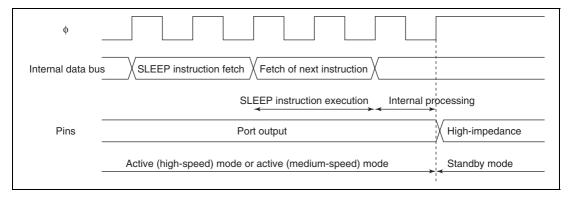


Figure 5.2 Standby Mode Transition and Pin States

5.3.5 Notes on External Input Signal Changes before/after Standby Mode

2. When external input signals cannot be captured because internal clock stops

- 1. When external input signal changes before/after standby mode or watch mode When an external input signal such as IRQ or WKP is input, both the high- and low-level widths of the signal must be at least two cycles of system clock φ or subclock φ_{SUB} (referred to together in this section as the internal clock). As the internal clock stops in standby mode and watch mode, the width of external input signals requires careful attention when a transition is made via these operating modes. Ensure that external input signals conform to the conditions stated in 3, Recommended timing of external input signals, below.
- The case of falling edge capture is illustrated in figure 5.3

 As shown in the case marked "Capture not possible," when an external input signal falls immediately after a transition to active (high-speed or medium-speed) mode or subactive mode, after oscillation is started by an interrupt via a different signal, the external input signal cannot be captured if the high-level width at that point is less than 2 t_{cut} or 2 t_{subcut}.
- 3. Recommended timing of external input signals

To ensure dependable capture of an external input signal, high- and low-level signal widths of at least 2 t_{cyc} or 2 t_{subcyc} are necessary before a transition is made to standby mode or watch mode, as shown in "Capture possible: case 1."

External input signal capture is also possible with the timing shown in "Capture possible: case 2" and "Capture possible: case 3," in which a 2 t_{evc} or 2 t_{subcyc} level width is secured.



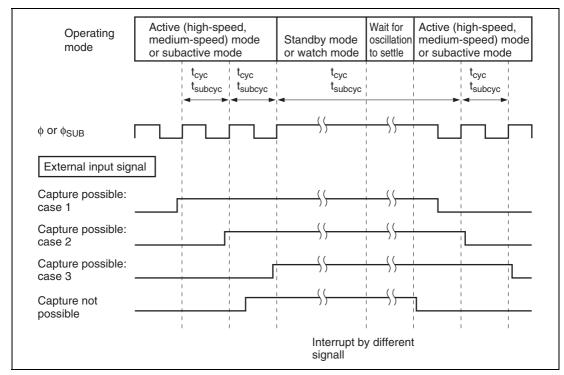


Figure 5.3 External Input Signal Capture when Signal Changes before/after Standby Mode or Watch Mode

4. Input pins to which these notes apply:

 \overline{IRQ}_4 to \overline{IRQ}_0 , \overline{WKP}_7 to \overline{WKP}_0 , \overline{ADTRG} , TMIC, TMIF, TMIG

5.4 Watch Mode

5.4.1 Transition to Watch Mode

The system goes from active or subactive mode to watch mode when a SLEEP instruction is executed while the SSBY bit in SYSCR1 is set to 1 and bit TMA3 in TMA is set to 1.

In watch mode, operation of on-chip peripheral modules is halted except for timer A, timer F, timer G, AEC and the LCD controller/driver (for which operation or halting can be set) is halted. As long as a minimum required voltage is applied, the contents of CPU registers, the on-chip RAM and some registers of the on-chip peripheral modules, are retained. I/O ports keep the same states as before the transition.

5.4.2 Clearing Watch Mode

Watch mode is cleared by an interrupt (timer A, timer F, timer G, IRQ_0 , or WKP_7 to WKP_0) or by input at the \overline{RES} pin.

Clearing by interrupt

When watch mode is cleared by interrupt, the mode to which a transition is made depends on the settings of LSON in SYSCR1 and MSON in SYSCR2. If both LSON and MSON are cleared to 0, transition is to active (high-speed) mode; if LSON = 0 and MSON = 1, transition is to active (medium-speed) mode; if LSON = 1, transition is to subactive mode. When the transition is to active mode, after the time set in SYSCR1 bits STS2 to STS0 has elapsed, a stable clock signal is supplied to the entire chip, watch mode is cleared, and interrupt exception handling starts. Watch mode is not cleared if the I bit of CCR is set to 1 or the particular interrupt is disabled in the interrupt enable register.

Clearing by RES input
 Clearing by RES pin is the same as for standby mode; see Clearing by RES pin in section 5.3.2, Clearing Standby Mode.

5.4.3 Oscillator Settling Time after Watch Mode is Cleared

The waiting time is the same as for standby mode; see section 5.3.3, Oscillator Settling Time after Standby Mode is Cleared.



5.4.4 Notes on External Input Signal Changes before/after Watch Mode

See section 5.3.5, Notes on External Input Signal Changes before/after Standby Mode.

5.5 Subsleep Mode

5.5.1 Transition to Subsleep Mode

The system goes from subactive mode to subsleep mode when a SLEEP instruction is executed while the SSBY bit in SYSCR1 is cleared to 0, LSON bit in SYSCR1 is set to 1, and TMA3 bit in TMA is set to 1. In subsleep mode, operation of on-chip peripheral modules other than the A/D converter WDT and PWM is halted. As long as a minimum required voltage is applied, the contents of CPU registers, the on-chip RAM and some registers of the on-chip peripheral modules are retained. I/O ports keep the same states as before the transition.

5.5.2 Clearing Subsleep Mode

Subsleep mode is cleared by an interrupt (timer A, timer C, timer F, timer G, asynchronous counter, SCI3-2, SCI3-1, IRQ₄ to IRQ₀, WKP₇ to WKP₀) or by a low input at the \overline{RES} pin.

- Clearing by interrupt
 - When an interrupt is requested, subsleep mode is cleared and interrupt exception handling starts. Subsleep mode is not cleared if the I bit of CCR is set to 1 or the particular interrupt is disabled in the interrupt enable register.
 - Interrupt signal and system clock are mutually asynchronous. Synchronization error time in a maximum is $2/\phi_{\text{SUB}}$ (s).
- Clearing by RES input
 Clearing by RES pin is the same as for standby mode; see Clearing by RES pin in section
 5.3.2, Clearing Standby Mode.

5.6 Subactive Mode

5.6.1 Transition to Subactive Mode

Subactive mode is entered from watch mode if a timer A, timer F, timer G, IRQ_0 , or WKP_7 to WKP_0 interrupt is requested while the LSON bit in SYSCR1 is set to 1. From subsleep mode, subactive mode is entered if a timer A, timer C, timer F, timer G, asynchronous counter, SCI3-1, SCI3-2, IRQ_4 to IRQ_0 , or WKP_7 to WKP_0 interrupt is requested. A transition to subactive mode does not take place if the I bit of CCR is set to 1 or the particular interrupt is disabled in the interrupt enable register.

5.6.2 Clearing Subactive Mode

Subactive mode is cleared by a SLEEP instruction or by a low input at the \overline{RES} pin.

- Clearing by SLEEP instruction
 - If a SLEEP instruction is executed while the SSBY bit in SYSCR1 is set to 1 and TMA3 bit in TMA is set to 1, subactive mode is cleared and watch mode is entered. If a SLEEP instruction is executed while SSBY = 0 and LSON = 1 in SYSCR1 and TMA3 = 1 in TMA, subsleep mode is entered. Direct transfer to active mode is also possible; see section 5.8, Direct Transfer, below.
- Clearing by RES pin
 Clearing by RES pin is the same as for standby mode; see Clearing by RES pin in section 5.3.2, Clearing Standby Mode.

5.6.3 Operating Frequency in Subactive Mode

The operating frequency in subactive mode is set in bits SA1 and SA0 in SYSCR2. The choices are $\phi_w/2$, $\phi_w/4$, and $\phi_w/8$.



5.7 Active (Medium-Speed) Mode

5.7.1 Transition to Active (Medium-Speed) Mode

If the \overline{RES} pin is driven low, active (medium-speed) mode is entered. If the LSON bit in SYSCR2 is set to 1 while the LSON bit in SYSCR1 is cleared to 0, a transition to active (medium-speed) mode results from IRQ_0 , IRQ_1 , or WKP_7 to WKP_0 interrupts in standby mode, timer A, timer F, timer G, IRQ_0 , or WKP_7 to WKP_0 interrupts in watch mode, or any interrupt in sleep mode. A transition to active (medium-speed) mode does not take place if the I bit of CCR is set to 1 or the particular interrupt is disabled in the interrupt enable register.

Furthermore, it sometimes acts with half state early timing at the time of transition to active (medium-speed) mode.

5.7.2 Clearing Active (Medium-Speed) Mode

Active (medium-speed) mode is cleared by a SLEEP instruction.

• Clearing by SLEEP instruction

A transition to standby mode takes place if the SLEEP instruction is executed while the SSBY bit in SYSCR1 is set to 1, the LSON bit in SYSCR1 is cleared to 0, and the TMA3 bit in TMA is cleared to 0. The system goes to watch mode if the SSBY bit in SYSCR1 is set to 1 and bit TMA3 in TMA is set to 1 when a SLEEP instruction is executed.

When both SSBY and LSON are cleared to 0 in SYSCR1 and a SLEEP instruction is executed, sleep mode is entered. Direct transfer to active (high-speed) mode or to subactive mode is also possible. See section 5.8, Direct Transfer, below for details.

Clearing by RES pin

When the \overline{RES} pin is driven low, a transition is made to the reset state and active (medium-speed) mode is cleared.

5.7.3 Operating Frequency in Active (Medium-Speed) Mode

Operation in active (medium-speed) mode is clocked at the frequency designated by the MA1 and MA0 bits in SYSCR1.

5.8 Direct Transfer

5.8.1 Overview of Direct Transfer

The CPU can execute programs in three modes: active (high-speed) mode, active (medium-speed) mode, and subactive mode. A direct transfer is a transition among these three modes without the stopping of program execution. A direct transfer can be made by executing a SLEEP instruction while the DTON bit in SYSCR2 is set to 1. After the mode transition, direct transfer interrupt exception handling starts.

If the direct transfer interrupt is disabled in interrupt enable register 2, a transition is made instead to sleep mode or watch mode. Note that if a direct transition is attempted while the I bit in CCR is set to 1, sleep mode or watch mode will be entered, and it will be impossible to clear the resulting mode by means of an interrupt.

- Direct transfer from active (high-speed) mode to active (medium-speed) mode
 When a SLEEP instruction is executed in active (high-speed) mode while the SSBY and
 LSON bits in SYSCR1 are cleared to 0, the MSON bit in SYSCR2 is set to 1, and the DTON
 bit in SYSCR2 is set to 1, a transition is made to active (medium-speed) mode via sleep mode.
- Direct transfer from active (medium-speed) mode to active (high-speed) mode
 When a SLEEP instruction is executed in active (medium-speed) mode while the SSBY and
 LSON bits in SYSCR1 are cleared to 0, the MSON bit in SYSCR2 is cleared to 0, and the
 DTON bit in SYSCR2 is set to 1, a transition is made to active (high-speed) mode via sleep mode.
- Direct transfer from active (high-speed) mode to subactive mode
 When a SLEEP instruction is executed in active (high-speed) mode while the SSBY and
 LSON bits in SYSCR1 are set to 1, the DTON bit in SYSCR2 is set to 1, and the TMA3 bit in
 TMA is set to 1, a transition is made to subactive mode via watch mode.
- Direct transfer from subactive mode to active (high-speed) mode
 When a SLEEP instruction is executed in subactive mode while the SSBY bit in SYSCR1 is set to 1, the LSON bit in SYSCR1 is cleared to 0, the MSON bit in SYSCR2 is cleared to 0, the DTON bit in SYSCR2 is set to 1, and the TMA3 bit in TMA is set to 1, a transition is made directly to active (high-speed) mode via watch mode after the waiting time set in SYSCR1 bits STS2 to STS0 has elapsed.



- Direct transfer from active (medium-speed) mode to subactive mode
 When a SLEEP instruction is executed in active (medium-speed) while the SSBY and LSON bits in SYSCR1 are set to 1, the DTON bit in SYSCR2 is set to 1, and the TMA3 bit in TMA is set to 1, a transition is made to subactive mode via watch mode.
- Direct transfer from subactive mode to active (medium-speed) mode
 When a SLEEP instruction is executed in subactive mode while the SSBY bit in SYSCR1 is
 set to 1, the LSON bit in SYSCR1 is cleared to 0, the MSON bit in SYSCR2 is set to 1, the
 DTON bit in SYSCR2 is set to 1, and the TMA3 bit in TMA is set to 1, a transition is made
 directly to active (medium-speed) mode via watch mode after the waiting time set in SYSCR1
 bits STS2 to STS0 has elapsed.

5.8.2 Direct Transition Times

(1) Time for Direct Transition from Active (High-Speed) Mode to Active (Medium-Speed) Mode

A direct transition from active (high-speed) mode to active (medium-speed) mode is performed by executing a SLEEP instruction in active (high-speed) mode while bits SSBY and LSON are both cleared to 0 in SYSCR1, and bits MSON and DTON are both set to 1 in SYSCR2. The time from execution of the SLEEP instruction to the end of interrupt exception handling (the direct transition time) is given by equation (1) below.

Direct transition time = { (Number of SLEEP instruction execution states) + (number of internal processing states) } × (tcyc before transition) + (number of interrupt exception handling execution states) × (tcyc after transition)

Example: Direct transition time = $(2 + 1) \times 2 tosc + 14 \times 16 tosc = 230 tosc$ (when $\phi/8$ is selected as the CPU operating clock)

Notation:

tosc: OSC clock cycle time tcyc: System clock (φ) cycle time

(2) Time for Direct Transition from Active (Medium-Speed) Mode to Active (High-Speed) Mode

A direct transition from active (medium-speed) mode to active (high-speed) mode is performed by executing a SLEEP instruction in active (medium-speed) mode while bits SSBY and LSON are both cleared to 0 in SYSCR1, and bit MSON is cleared to 0 and bit DTON is set to 1 in SYSCR2.

The time from execution of the SLEEP instruction to the end of interrupt exception handling (the direct transition time) is given by equation (2) below.

Direct transition time = { (Number of SLEEP instruction execution states) + (number of internal processing states) } × (tcyc before transition) + (number of interrupt exception handling execution states) × (tcyc after transition)

Example: Direct transition time = $(2 + 1) \times 16$ tosc + 14×2 tosc = 76tosc (when $\phi/8$ is selected as the CPU operating clock)

Notation:

tosc: OSC clock cycle time tcyc: System clock (φ) cycle time

(3) Time for Direct Transition from Subactive Mode to Active (High-Speed) Mode

A direct transition from subactive mode to active (high-speed) mode is performed by executing a SLEEP instruction in subactive mode while bit SSBY is set to 1 and bit LSON is cleared to 0 in SYSCR1, bit MSON is cleared to 0 and bit DTON is set to 1 in SYSCR2, and bit TMA3 is set to 1 in TMA. The time from execution of the SLEEP instruction to the end of interrupt exception handling (the direct transition time) is given by equation (3) below.

Direct transition time = { (Number of SLEEP instruction execution states) + (number of internal processing states) } \times (tsubcyc before transition) + { (wait time set in STS2 to STS0) + (number of interrupt exception handling execution states) } \times (tcyc after transition)(3)

Example: Direct transition time = $(2 + 1) \times 8tw + (8192 + 14) \times 2tosc = 24tw + 16412tosc$ (when $\phi w/8$ is selected as the CPU operating clock, and wait time = 8192 states)

Notation:

tosc: OSC clock cycle time tw: Watch clock cycle time tcyc: System clock (ϕ) cycle time tsubcyc: Subclock (ϕ_{SUB}) cycle time



(4) Time for Direct Transition from Subactive Mode to Active (Medium-Speed) Mode

A direct transition from subactive mode to active (medium-speed) mode is performed by executing a SLEEP instruction in subactive mode while bit SSBY is set to 1 and bit LSON is cleared to 0 in SYSCR1, bits MSON and DTON are both set to 1 in SYSCR2, and bit TMA3 is set to 1 in TMA. The time from execution of the SLEEP instruction to the end of interrupt exception handling (the direct transition time) is given by equation (4) below.

Direct transition time = { (Number of SLEEP instruction execution states) + (number of internal processing states) } \times (tsubcyc before transition) + { (wait time set in STS2 to STS0) + (number of interrupt exception handling execution states) } \times (tcyc after transition)(4)

Example: Direct transition time = $(2 + 1) \times 8$ tw + $(8192 + 14) \times 16$ tosc = 24tw + 131296tosc (when ϕ w/8 or ϕ 8 is selected as the CPU operating clock, and wait time = 8192 states)

Notation:

tosc: OSC clock cycle time tw: Watch clock cycle time tcyc: System clock (ϕ) cycle time tsubcyc: Subclock (ϕ_{SUB}) cycle time

5.8.3 Notes on External Input Signal Changes before/after Direct Transition

- Direct transition from active (high-speed) mode to subactive mode
 Since the mode transition is performed via watch mode, see section 5.3.5, Notes on External Input Signal Changes before/after Standby Mode.
- Direct transition from active (medium-speed) mode to subactive mode
 Since the mode transition is performed via watch mode, see section 5.3.5, Notes on External Input Signal Changes before/after Standby Mode.
- Direct transition from subactive mode to active (high-speed) mode
 Since the mode transition is performed via watch mode, see section 5.3.5, Notes on External Input Signal Changes before/after Standby Mode.
- 4. Direct transition from subactive mode to active (medium-speed) mode Since the mode transition is performed via watch mode, see section 5.3.5, Notes on External Input Signal Changes before/after Standby Mode.

5.9 Module Standby Mode

5.9.1 Setting Module Standby Mode

Module standby mode is set for individual peripheral functions. All the on-chip peripheral modules can be placed in module standby mode. When a module enters module standby mode, the system clock supply to the module is stopped and operation of the module halts. This state is identical to standby mode.

Module standby mode is set for a particular module by setting the corresponding bit to 0 in clock stop register 1 (CKSTPR1) or clock stop register 2 (CKSTPR2). (See table 5.5.)

5.9.2 Clearing Module Standby Mode

Module standby mode is cleared for a particular module by setting the corresponding bit to 1 in clock stop register 1 (CKSTPR1) or clock stop register 2 (CKSTPR2). (See table 5.5.)

Following a reset, clock stop register 1 (CKSTPR1) and clock stop register 2 (CKSTPR2) are both initialized to H'FF.



Table 5.5 Setting and Clearing Module Standby Mode by Clock Stop Register

Register Name	Bit Name		Operation
CKSTPR1	TACKSTP	1	Timer A module standby mode is cleared
		0	Timer A is set to module standby mode
	TCCKSTP	1	Timer C module standby mode is cleared
		0	Timer C is set to module standby mode
	TFCKSTP	1	Timer F module standby mode is cleared
		0	Timer F is set to module standby mode
	TGCKSTP	1	Timer G module standby mode is cleared
		0	Timer G is set to module standby mode
	ADCKSTP	1	A/D converter module standby mode is cleared
		0	A/D converter is set to module standby mode
	S32CKSTP	1	SCI3-2 module standby mode is cleared
		0	SCI3-2 is set to module standby mode
	S31CKSTP	1	SCI3-1 module standby mode is cleared
		0	SCI3-1 is set to module standby mode
CKSTPR2	LDCKSTP	1	LCD module standby mode is cleared
		0	LCD is set to module standby mode
	PWCKSTP	1	PWM module standby mode is cleared
		0	PWM is set to module standby mode
	WDCKSTP	1	Watchdog timer module standby mode is cleared
		0	Watchdog timer is set to module standby mode
	AECKSTP	1	Asynchronous event counter module standby mode is cleared
		0	Asynchronous event counter is set to module standby mode

Note: For details of module operation, see the sections on the individual modules.

5.9.3 Usage Note

If, due to the timing with which a peripheral module issues interrupt requests, the module in question is set to module standby mode before an interrupt is processed, the module will stop with the interrupt request still pending. In this situation, interrupt processing will be repeated indefinitely unless interrupts are prohibited.

It is therefore necessary to ensure that no interrupts are generated when a module is set to module standby mode. The surest way to do this is to specify the module standby mode setting only when interrupts are prohibited (interrupts prohibited using the interrupt enable register or interrupts masked using bit CCR-I).



Section 6 ROM

6.1 Overview

The H8/38532 has 16 Kbytes of on-chip mask ROM, the H8/38533 has 24 Kbytes, the H8/38534 has 32 Kbytes, the H8/38535 has 40 Kbytes, the H8/38536 has 48 Kbytes, and the H8/38537 has 60 Kbytes. The ROM is connected to the CPU by a 16-bit data bus, allowing high-speed two-state access for both byte data and word data.

The flash memory version of the H8/38537 is equipped with 60 Kbytes of flash memory. The flash memory version of the H8/38534 is equipped with 32 Kbytes of flash memory.

6.2 Flash Memory Overview

6.2.1 Features

The features of the 60 Kbytes or 32 Kbytes of flash memory built into the flash memory versions are summarized below.

- Programming/erase methods
 - The 60-Kbyte flash memory is programmed 128 bytes at a time. Erase is performed in single-block units. The flash memory is configured as follows: 1 Kbyte × 4 blocks, 28 Kbytes × 1 block, 16 Kbytes × 1 block, 8 Kbytes × 1 block, 4 Kbytes × 1 block. The 32-Kbyte flash memory is configured as follows: 1 Kbyte × 4 blocks, 28 Kbytes × 1 block. To erase the entire flash memory, each block must be erased in turn.
- Reprogramming capability
 - The flash memory can be reprogrammed up to 1,000 times.
- On-board programming
 - On-board programming/erasing can be done in boot mode, in which the boot program built into the chip is started to erase or program of the entire flash memory. In normal user program mode, individual blocks can be erased or programmed.
- Programmer mode
 - Flash memory can be programmed/erased in programmer mode using a PROM programmer, as well as in on-board programming mode.
- Automatic bit rate adjustment
 - For data transfer in boot mode, this LSI's bit rate can be automatically adjusted to match the transfer bit rate of the host.

- Programming/erasing protection
 - Sets software protection against flash memory programming/erasing.
- Power-down mode
 - The power supply circuit is partly halted in the subactive mode and can be read in the power-down mode.

6.2.2 Block Diagram

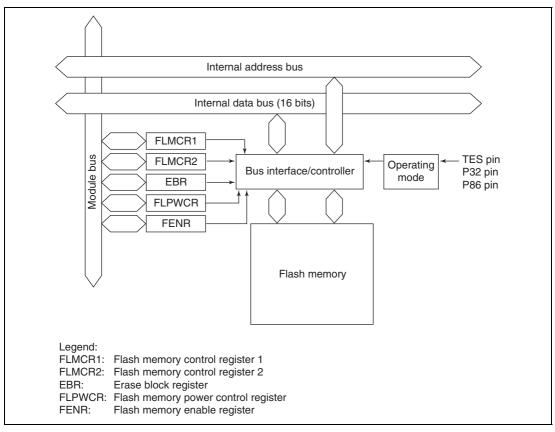


Figure 6.1 Block Diagram of Flash Memory

6.2.3 Block Configuration

Figure 6.2 shows the block configuration of flash memory. The thick lines indicate erasing units, the narrow lines indicate programming units, and the values are addresses. The flash memory is divided into 1 Kbyte \times 4 blocks, 28 Kbytes \times 1 block, 16 Kbytes \times 1 block, 8 Kbytes \times 1 block, and 4 Kbytes \times 1 block. Erasing is performed in these units. Programming is performed in 128-byte units starting from an address with lower eight bits H'00 or H'80.

ſ	H'0000	H'0001	H'0002	← Programming unit: 128 bytes →	H'007F
	H'0080	H'0081	H'0082		H'00FI
Erase unit			 		1
1 Kbyte					
	H'0380	H'0381	H'0382		H'03F
	H'0400	H'0401	H'0402	← Programming unit: 128 bytes →	H'047I
Erase unit	H'0480	H'0481	H'0482		H'04F
1 Kbyte			 		!
İ	H'0780	H'0781	H'0782		H'07F
t	H'0800	H'0801	H'0802	← Programming unit: 128 bytes →	H'087I
ŀ	H'0880	H'0881	H'0882		H'080
Erase unit 1 Kbyte					
Ī	H'0B80	H'0B81	H'0B82		H'0BF
Ī	H'0C00	H'0C01	H'0C02	← Programming unit: 128 bytes →	H'0C7
Ī	H'0C80	H'0C81	H'0C82		H'0CF
Erase unit 1 Kbyte					
	H'0F80	H'0F81	H'0F82		H'0FF
[H'1000	H'1001	H'1002	← Programming unit: 128 bytes →	H'107
	H'1080	H'1081	H'1082		H'10F
Erase unit 28 Kbytes					- - - - - -
ŀ	H'7F80	H'7F81	H'7F82		H'7FF
Ī	H'8000	H'8001	H'8002	← Programming unit: 128 bytes →	H'807
Ī	H'8080	H'8081	H'8082		H'8CF
Erase unit 16 Kbytes			 		1
Ī	H'BF80	H'BF81	H'BF82		H'BFF
Ī	H'C000	H'C001	H'C002	← Programming unit: 128 bytes →	H'C07
Ţ	H'C080	H'C081	H'C082		H'CCF
Erase unit 8 Kbytes					
	H'DF80	H'DF81	H'DF82		H'DFF
[H'E000	H'E001	H'E002	← Programming unit: 128 bytes →	H'E07
	H'E080	H'E081	H'E082		H'ECF
Erase unit 4 Kbytes			 		1
[H'EF80	H'EF81	H'EF82		H'EFF

Figure 6.2 Flash Memory Block Configuration



6.2.4 Register Configuration

Table 6.1 lists the register configuration to control the flash memory when the built in flash memory is effective.

Table 6.1 Register Configuration

Register Name	Abbreviation	R/W	Initial Value	Address
Flash memory control register 1	FLMCR1	R/W	H'00	H'F020
Flash memory control register 2	FLMCR2	R	H'00	H'F021
Flash memory power control register	FLPWCR	R/W	H'00	H'F022
Erase block register	EBR	R/W	H'00	H'F023
Flash memory enable register	FENR	R/W	H'00	H'F02B

Note: FLMCR1, FLMCR2, FLPWCR, EBR, and FENR are 8 bit registers. Only byte access is enabled which are two-state access. These registers are dedicated to the product in which flash memory is included. The product in which mask ROM is included does not have these registers. When the corresponding address is read in these products, the value is undefined. A write is disabled.

6.3 Descriptions of Registers of the Flash Memory

6.3.1 Flash Memory Control Register 1 (FLMCR1)

Bit	7	6	5	4	3	2	1	0
	_	SWE	ESU	PSU	EV	PV	E	Р
Initial value	0	0	0	0	0	0	0	0
Read/Write	_	R/W						

FLMCR1 is a register that makes the flash memory change to program mode, program-verify mode, erase mode, or erase-verify mode. For details on register setting, refer to section 6.5, Flash Memory Programming/Erasing. By setting this register, the flash memory enters program mode, erase mode, program-verify mode, or erase-verify mode. Read the data in the state that bits 6 to 0 of this register are cleared when using flash memory as normal built-in ROM.

Bit 7—Reserved

This bit is always read as 0 and cannot be modified.

Bit 6—Software Write Enable (SWE)

This bit is to set enabling/disabling of programming/enabling of flash memory (set when bits 5 to 0 and the EBR register are to be set).

Bit 6 SWE	Description
0	Programming/erasing is disabled. Other FLMCR1 register bits and all EBR bits cannot be set. (initial value)
1	Flash memory programming/erasing is enabled.

Bit 5—Erase Setup (ESU)

This bit is to prepare for changing to erase mode. Set this bit to 1 before setting the E bit to 1 in FLMCR1 (do not set SWE, PSU, EV, PV, E, and P bits at the same time).

Bit 5 ESU	Description	
0	The erase setup state is cancelled	(initial value)
1	The flash memory changes to the erase setup state. Set this bit to 1 b the E bit to 1 in FLMCR1.	efore setting

Bit 4—Program Setup (PSU)

This bit is to prepare for changing to program mode. Set this bit to 1 before setting the P bit to 1 in FLMCR1 (do not set SWE, ESU, EV, PV, E, and P bits at the same time).

Bit 4 PSU	Description	
0	The program setup state is cancelled	(initial value)
1	The flash memory changes to the program setup state. Set this bit to setting the P bit to 1 in FLMCR1.	1 before

Bit 3—Erase-Verify (EV)

This bit is to set changing to or cancelling erase-verify mode (do not set SWE, ESU, PSU, PV, E, and P bits at the same time).

Bit 3 EV	Description	
0	Erase-verify mode is cancelled	(initial value)
1	The flash memory changes to erase-verify mode	

Bit 2—Program-Verify (PV)

This bit is to set changing to or cancelling program-verify mode (do not set SWE, ESU, PSU, EV, E, and P bits at the same time).

Bit 2 PV	Description	
0	Program-verify mode is cancelled	(initial value)
1	The flash memory changes to program-verify mode	

Bit 1—Erase (E)

This bit is to set changing to or cancelling erase mode (do not set SWE, ESU, PSU, EV, PV, and P bits at the same time).

Bit 1 E	Description	
0	Erase mode is cancelled	(initial value)
1	When this bit is set to 1, while the SWE = 1 and ESU = 1, the flash meaning to erase mode.	nemory

Bit 0—Program (P)

This bit is to set changing to or cancelling program mode (do not set SWE, ESU, PSU, EV, PV, and E bits at the same time).

Bit 0 P	Description	
0	Program mode is cancelled	(initial value)
1	When this bit is set to 1, while the SWE = 1 and PSU = 1, the flash n changes to program mode.	nemory

6.3.2 Flash Memory Control Register 2 (FLMCR2)

Bit	7	6	5	4	3	2	1	0	
	FLER	_	_	_	_	_	_	_	
Initial value	0	0	0	0	0	0	0	0	_
Read/Write	R/W	_	_	_	_	_	_	_	

FLMCR2 is a register that displays the state of flash memory programming/erasing. FLMCR2 is a read-only register, and should not be written to.

Bit 7—Flash Memory Error (FLER)

This bit is set when the flash memory detects an error and goes to the error-protection state during programming or erasing to the flash memory. See section 6.6.3, Error Protection, for details.

Bit 7 FLER	Description	
0	The flash memory operates normally.	(initial value)
1	Indicates that an error has occurred during an operation (programming or erasing).	n on flash memory

Bits 6 to 0—Reserved

These bits are always read as 0 and cannot be modified.

6.3.3 Erase Block Register (EBR)

Bit	7	6	5	4	3	2	1	0
	EB7	EB6	EB5	EB4	EB3	EB2	EB1	EB0
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W							

EBR specifies the flash memory erase area block. EBR is initialized to H'00 when the SWE bit in FLMCR1 is 0. Do not set more than one bit at a time, as this will cause all the bits in EBR to be automatically cleared to 0. When each bit is set to 1 in EBR, the corresponding block can be erased. Other blocks change to the erase-protection state. See table 6.2 for the method of dividing blocks of the flash memory. When the whole bits are to be erased, erase them in turn in unit of a block.

Table 6.2 Division of Blocks to Be Erased

EBR	Bit Name	Block (Size)	Address
0	EB0	EB0 (1 Kbyte)	H'0000 to H'03FF
1	EB1	EB1 (1 Kbyte)	H'0400 to H'07FF
2	EB2	EB2 (1 Kbyte)	H'0800 to H'0BFF
3	EB3	EB3 (1 Kbyte)	H'0C00 to H'0FFF
4	EB4	EB4 (28 Kbytes)	H'1000 to H'7FFF
5	EB5	EB5 (16 Kbytes)	H'8000 to H'BFFF
6	EB6	EB6 (8 Kbytes)	H'C000 to H'DFFF
7	EB7	EB7 (4 Kbytes)	H'E000 to H'EFFF

6.3.4 Flash Memory Power Control Register (FLPWCR)

Bit	7	6	5	4	3	2	1	0	
	PDWND	_	_	_	_	_	_	_	
Initial value	0	0	0	0	0	0	0	0	_
Read/Write	R/W	_	_	_	_	_	_	_	

FLPWCR enables or disables a transition to the flash memory power-down mode when the LSI switches to subactive mode. The power supply circuit can be read in the subactive mode, although it is partly halted in the power-down mode.

Bit 7—Power-down Disable (PDWND)

This bit selects the power-down mode of the flash memory when a transition to the subactive mode is made.

Bit 7 PDWND	Description
0	When this bit is 0 and a transition is made to the subactive mode, the flash memory enters the power-down mode. (initial value)
1	When this bit is 1, the flash memory remains in the normal mode even after a transition is made to the subactive mode.

Bits 6 to 0—Reserved

These bits are always read as 0 and cannot be modified.

6.3.5 Flash Memory Enable Register (FENR)

Bit	7	6	5	4	3	2	1	0	_
	FLSHE	_	_	_	_	_	_	_	
Initial value	0	0	0	0	0	0	0	0	_
Read/Write	R/W	_	_	_	_	_	_	_	

FENR controls CPU access to the flash memory control registers, FLMCR1, FLMCR2, EBR, and FLPWCR.

Bit 7—Flash Memory Control Register Enable (FLSHE)

This bit controls access to the flash memory control registers.

Bit 7 FLSHE	Description	
0	Flash memory control registers cannot be accessed	(initial value)
1	Flash memory control registers can be accessed	

Bits 6 to 0—Reserved

These bits are always read as 0 and cannot be modified.



6.4 On-Board Programming Modes

There are two modes for programming/erasing of the flash memory; boot mode, which enables on-board programming/erasing, and programmer mode, in which programming/erasing is performed with a PROM programmer. On-board programming/erasing can also be performed in user program mode. At reset-start in reset mode, this LSI changes to a mode depending on the TEST pin settings, P32 pin settings, and input level of each port, as shown in table 6.3. The input level of each pin must be defined four states before the reset ends.

When changing to boot mode, the boot program built into this LSI is initiated. The boot program transfers the programming control program from the externally-connected host to on-chip RAM via SCI32. After erasing the entire flash memory, the programming control program is executed. This can be used for programming initial values in the on-board state or for a forcible return when programming/erasing can no longer be done in user program mode. In user program mode, individual blocks can be erased and programmed by branching to the user program/erase control program prepared by the user.

Table 6.3 Setting Programming Modes

TEST	P32	P86	PB0	PB1	PB2	LSI State after Reset End
0	1	Х	Х	Х	Х	User Mode
0	0	1	Χ	Х	Х	Boot Mode
1	Х	Х	0	0	0	Programmer Mode

X: Don't care

6.4.1 Boot Mode

Table 6.4 shows the boot mode operations between reset end and branching to the programming control program. The device uses SCI32 in the boot mode.

- 1. When boot mode is used, the flash memory programming control program must be prepared in the host beforehand. Prepare a programming control program in accordance with the description in section 6.5, Flash Memory Programming/Erasing.
- 2. SCI3 should be set to asynchronous mode, and the transfer format as follows: 8-bit data, 1 stop bit, and no parity. The inversion function of TXD and RXD pins by the SPCR register is set to "Not to be inverted," so do not put the circuit for inverting a value between the host and this LSI.

- 3. When the boot program is initiated, the chip measures the low-level period of asynchronous SCI communication data (H'00) transmitted continuously from the host. The chip then calculates the bit rate of transmission from the host, and adjusts the SCI3 bit rate to match that of the host. The reset should end with the RXD pin high. The RXD and TXD pins should be pulled up on the board if necessary. After the reset is complete, it takes approximately 100 states before the chip is ready to measure the low-level period.
- 4. After matching the bit rates, the chip transmits one H'00 byte to the host to indicate the completion of bit rate adjustment. The host should confirm that this adjustment end indication (H'00) has been received normally, and transmit one H'55 byte to the chip. If reception could not be performed normally, initiate boot mode again by a reset. Depending on the host's transfer bit rate and system clock frequency of this LSI, there will be a discrepancy between the bit rates of the host and the chip. To operate the SCI properly, set the host's transfer bit rate and system clock frequency of this LSI within the ranges listed in table 6.5.
- 5. In boot mode, a part of the on-chip RAM area is used by the boot program. The area H'F780 to H'FEEF is the area to which the programming control program is transferred from the host. The boot program area cannot be used until the execution state in boot mode switches to the programming control program.
- 6. Before branching to the programming control program, the chip terminates transfer operations by SCI3 (by clearing the RE and TE bits in SCR to 0), however the adjusted bit rate value remains set in BRR. Therefore, the programming control program can still use it for transfer of write data or verify data with the host. The TXD pin is high (PCR42 = 1, P42 = 1). The contents of the CPU general registers are undefined immediately after branching to the programming control program. These registers must be initialized at the beginning of the programming control program, as the stack pointer (SP), in particular, is used implicitly in subroutine calls, etc.
- 7. Boot mode can be cleared by a reset. End the reset after driving the reset pin low, waiting at least 20 states, and then setting the TEST pin and P32 pin. Boot mode is also cleared when a WDT overflow occurs.
- 8. Do not change the TEST pin and P32 pin input levels in boot mode.



Table 6.4 Boot Mode Operation

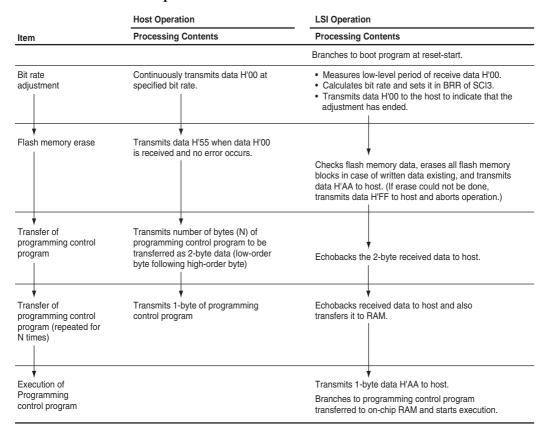


Table 6.5 Oscillating Frequencies (f_{osc}) for which Automatic Adjustment of LSI Bit Rate Is Possible

Product Group	Host Bit Rate	Oscillating Frequencies (f_{osc}) Range of LSI
H8/38537	19,200 bps	16 MHz
H8/38534 Flash memory version	9,600 bps	8 to 16 MHz
riadirinemory version	4,800 bps	6 to 16 MHz
	2,400 bps	2 to 16 MHz
	1,200 bps	2 to 16 MHz

6.4.2 Programming/Erasing in User Program Mode

The term user mode refers to the status when a user program is being executed. On-board programming/erasing of an individual flash memory block can also be performed in user program mode by branching to a user program/erase control program. The user must set branching conditions and provide on-board means of supplying programming data. The flash memory must contain the user program/erase control program or a program that provides the user program/erase control program from external memory. As the flash memory itself cannot be read during programming/erasing, transfer the user program/erase control program to on-chip RAM, as in boot mode. Figure 6.3 shows a sample procedure for programming/erasing in user program mode. Prepare a user program/erase control program in accordance with the description in section 6.5, Flash Memory Programming/Erasing.

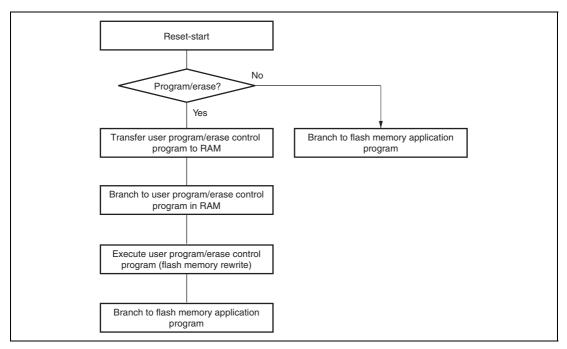


Figure 6.3 Programming/Erasing Flowchart Example in User Program Mode

6.5 Flash Memory Programming/Erasing

A software method using the CPU is employed to program and erase flash memory in the on-board programming modes. Depending on the FLMCR1 setting, the flash memory operates in one of the following four modes: Program mode, program-verify mode, erase mode, and erase-verify mode. The programming control program in boot mode and the user program/erase control program in user program mode use these operating modes in combination to perform programming/erasing. Flash memory programming and erasing should be performed in accordance with the descriptions in section 6.5.1, Program/Program-Verify and section 6.5.2, Erase/Erase-Verify, respectively.

6.5.1 Program/Program-Verify

When writing data or programs to the flash memory, the program/program-verify flowchart shown in figure 6.4 should be followed. Performing programming operations according to this flowchart will enable data or programs to be written to the flash memory without subjecting the chip to voltage stress or sacrificing program data reliability.

- 1. Programming must be done to an empty address. Do not reprogram an address to which programming has already been performed.
- 2. Programming should be carried out 128 bytes at a time. A 128-byte data transfer must be performed even if writing fewer than 128 bytes. In this case, H'FF data must be written to the extra addresses.
- 3. Prepare the following data storage areas in RAM: A 128-byte programming data area, a 128-byte reprogramming data area, and a 128-byte additional-programming data area. Perform reprogramming data computation according to table 6.6, and additional programming data computation according to table 6.7.
- 4. Consecutively transfer 128 bytes of data in byte units from the reprogramming data area or additional-programming data area to the flash memory. The program address and 128-byte data are latched in the flash memory. The lower 8 bits of the start address in the flash memory destination area must be H'00 or H'80.
- 5. The time during which the P bit is set to 1 is the programming time. Figure 6.8 shows the allowable programming times.
- 6. The watchdog timer (WDT) is set to prevent overprogramming due to program runaway, etc. An overflow cycle of approximately 6.6 ms is allowed.
- 7. For a dummy write to a verify address, write 1-byte data H'FF to an address whose lower 1 bit is b'0. Verify data can be read in word size from the address to which a dummy write was performed.
- 8. The maximum number of repetitions of the program/program-verify sequence of the same bit is 1,000.



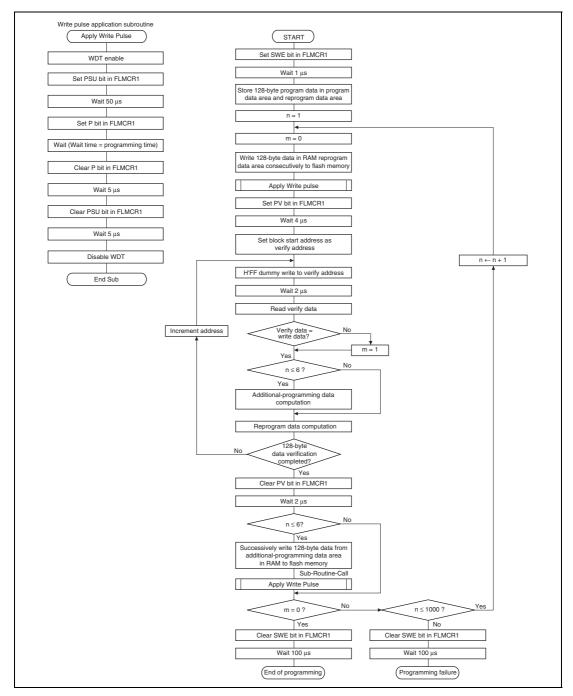


Figure 6.4 Program/Program-Verify Flowchart

Table 6.6 Reprogram Data Computation Table

Program Data	Verify Data	Reprogram Data	Comments
0	0	1	Programming completed
0	1	0	Reprogram bit
1	0	1	_
1	1	1	Remains in erased state

 Table 6.7
 Additional-Program Data Computation Table

Reprogram Data	Verify Data	Additional-Program Data	Comments
0	0	0	Additional-program bit
0	1	1	No additional programming
1	0	1	No additional programming
1	1	1	No additional programming

Table 6.8 Programming Time

n (Number of Writes)	Programming Time	In Additional Programming	Comments
1 to 6	30	10	
7 to 1,000	200	_	

Note: Time shown in μs.

6.5.2 Erase/Erase-Verify

When erasing flash memory, the erase/erase-verify flowchart shown in figure 6.5 should be followed.

- 1. Prewriting (setting erase block data to all 0s) is not necessary.
- 2. Erasing is performed in block units. Make only a single-bit specification in the erase block register (EBR). To erase multiple blocks, each block must be erased in turn.
- 3. The time during which the E bit is set to 1 is the flash memory erase time.
- 4. The watchdog timer (WDT) is set to prevent overerasing due to program runaway, etc. An overflow cycle of approximately 19.8 ms is allowed.
- 5. For a dummy write to a verify address, write 1-byte data H'FF to an address whose lower 1 bit is b'0. Verify data can be read in word size from the address to which a dummy write was performed.
- 6. If the read data is not erased successfully, set erase mode again, and repeat the erase/erase-verify sequence as before. The maximum number of repetitions of the erase/erase-verify sequence is 100.

6.5.3 Interrupt Handling when Programming/Erasing Flash Memory

All interrupts, are disabled while flash memory is being programmed or erased, or while the boot program is executing, for the following three reasons:

- 1. Interrupt during programming/erasing may cause a violation of the programming or erasing algorithm, with the result that normal operation cannot be assured.
- 2. If interrupt exception handling starts before the vector address is written or during programming/erasing, a correct vector cannot be fetched and the CPU malfunctions.
- 3. If an interrupt occurs during boot program execution, normal boot mode sequence cannot be carried out.



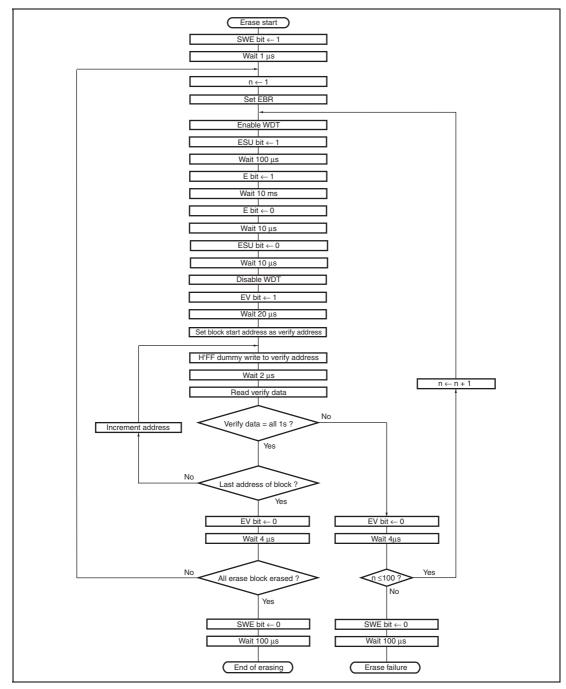


Figure 6.5 Erase/Erase-Verify Flowchart

6.6 Program/Erase Protection

There are three kinds of flash memory program/erase protection; hardware protection, software protection, and error protection.

6.6.1 Hardware Protection

Hardware protection refers to a state in which programming/erasing of flash memory is forcibly disabled or aborted because of a transition to reset, subactive mode, subsleep mode, watch mode, or standby mode. Flash memory control register 1 (FLMCR1), flash memory control register 2 (FLMCR2), and erase block register (EBR) are initialized. In a reset via the \overline{RES} pin, the reset state is not entered unless the \overline{RES} pin is held low until oscillation stabilizes after powering on. In the case of a reset during operation, hold the \overline{RES} pin low for the \overline{RES} pulse width specified in the AC Characteristics section.

6.6.2 Software Protection

Software protection can be implemented against programming/erasing of all flash memory blocks by clearing the SWE bit in FLMCR1. When software protection is in effect, setting the P or E bit in FLMCR1 does not cause a transition to program mode or erase mode. By setting the erase block register (EBR), erase protection can be set for individual blocks. When EBR is set to H'00, erase protection is set for all blocks.

6.6.3 Error Protection

In error protection, an error is detected when CPU runaway occurs during flash memory programming/erasing, or operation is not performed in accordance with the program/erase algorithm, and the program/erase operation is aborted. Aborting the program/erase operation prevents damage to the flash memory due to overprogramming or overerasing.

When the following errors are detected during programming/erasing of flash memory, the FLER bit in FLMCR2 is set to 1, and the error protection state is entered.

- When the flash memory of the relevant address area is read during programming/erasing (including vector read and instruction fetch)
- Immediately after exception handling excluding a reset during programming/erasing
- When a SLEEP instruction is executed during programming/erasing



The FLMCR1, FLMCR2, and EBR settings are retained, however program mode or erase mode is aborted at the point at which the error occurred. Program mode or erase mode cannot be re-entered by re-setting the P or E bit. However, PV and EV bit setting is enabled, and a transition can be made to verify mode. Error protection can be cleared only by a power-on reset.

6.7 Programmer Mode

In programmer mode, a PROM programmer can be used to perform programming/erasing via a socket adapter, just as a discrete flash memory. Use a PROM programmer that supports the MCU device type with the on-chip Renesas Technology (former Hitachi Ltd.) 64-Kbyte flash memory (F-ZTAT64V3). A 10-MHz input clock is required. For the conditions for transition to programmer mode, see table 6.3.

6.7.1 Socket Adapter

The socket adapter converts the pin allocation of the flash memory device to that of the discrete flash memory HN28F101. The address of the on-chip flash memory is H'0000 to H'EFFF. Figure 6.6 shows a socket-adapter-pin correspondence diagram.

6.7.2 Programmer Mode Commands

The following commands are supported in programmer mode.

- Memory Read Mode
- Auto-Program Mode
- Auto-Erase Mode
- Status Read Mode

Status polling is used for auto-programming, auto-erasing, and status read modes. In status read mode, detailed internal information is output after the execution of auto-programming or auto-erasing. Table 6.9 shows the sequence of each command. In auto-programming mode, 129 cycles are required since 128 bytes are written at the same time. In memory read mode, the number of cycles depends on the number of address write cycles (n).



 Table 6.9
 Command Sequence in Programmer Mode

	Number	1st Cycle			2nd Cycle		
Command Name	of Cycles	Mode	Address	Data	Mode	Address	Data
Memory read	1 + n	Write	Х	H'00	Read	RA	Dout
Auto-program	129	Write	Х	H'40	Write	WA	Din
Auto-erase	2	Write	Х	H'20	Write	Х	H'20
Status read	2	Write	Х	H'71	Write	Х	H'71

n: the number of address write cycles



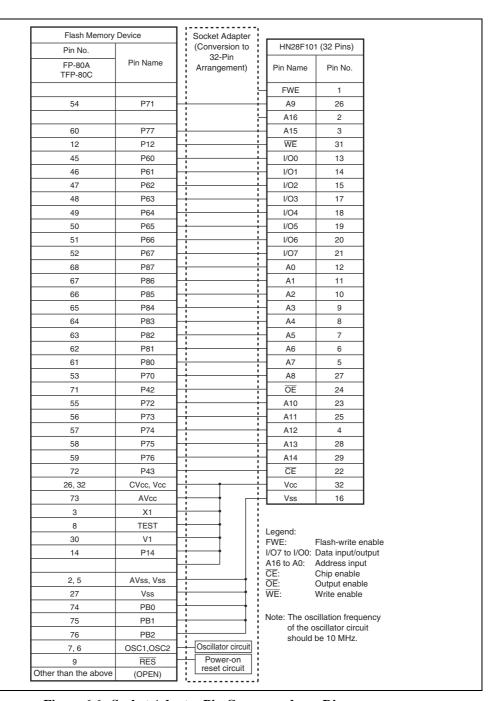


Figure 6.6 Socket Adapter Pin Correspondence Diagram

6.7.3 Memory Read Mode

- After completion of auto-program/auto-erase/status read operations, a transition is made to the command wait state. When reading memory contents, a transition to memory read mode must first be made with a command write, after which the memory contents are read. Once memory read mode has been entered, consecutive reads can be performed.
- 2. In memory read mode, command writes can be performed in the same way as in the command wait state.
- 3. After powering on, memory read mode is entered.
- 4. Tables 6.10 to 6.12 show the AC characteristics.

Table 6.10 AC Characteristics in Transition to Memory Read Mode

Conditions: $V_{cc} = 3.3 \text{ V} \pm 0.3 \text{ V}, V_{ss} = 0 \text{ V}, T_a = 25^{\circ}\text{C} \pm 5^{\circ}\text{C}$

Item	Symbol	Min	Max	Unit	Notes
Command write cycle	t _{nxtc}	20	_	μs	Figure 6.7
CE hold time	t _{ceh}	0	_	ns	
CE setup time	t _{ces}	0	_	ns	_
Data hold time	t _{dh}	50	_	ns	_
Data setup time	t _{ds}	50	_	ns	_
Write pulse width	t _{wep}	70	_	ns	_
WE rise time	t,	_	30	ns	
WE fall time	t,	_	30	ns	_

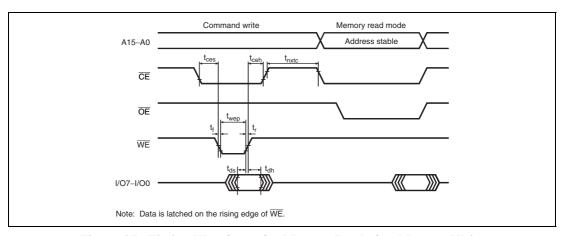


Figure 6.7 Timing Waveforms for Memory Read after Memory Write

Table 6.11 AC Characteristics in Transition from Memory Read Mode to Another Mode

Conditions: $V_{CC} = 3.3 \text{ V} \pm 0.3 \text{ V}$, $V_{SS} = 0 \text{ V}$, $T_a = 25^{\circ}\text{C} \pm 5^{\circ}\text{C}$

Item	Symbol	Min	Max	Unit	Notes
Command write cycle	t _{nxte}	20	_	μs	Figure 6.8
CE hold time	t _{ceh}	0	_	ns	
CE setup time	t _{ces}	0	_	ns	
Data hold time	t _{dh}	50	_	ns	
Data setup time	t _{ds}	50	_	ns	
Write pulse width	t _{wep}	70	_	ns	
WE rise time	t _r	_	30	ns	
WE fall time	t,	_	30	ns	

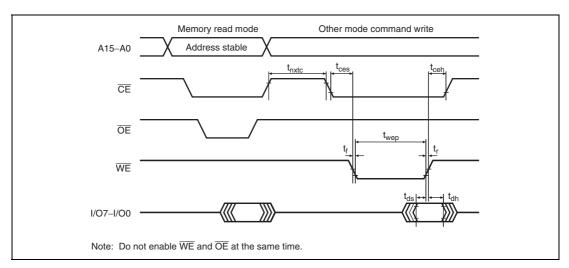


Figure 6.8 Timing Waveforms in Transition from Memory Read Mode to Another Mode

Table 6.12 AC Characteristics in Memory Read Mode

Conditions: $V_{cc} = 3.3 \text{ V} \pm 0.3 \text{ V}, V_{ss} = 0 \text{ V}, T_a = 25^{\circ}\text{C} \pm 5^{\circ}\text{C}$

Item	Symbol	Min	Max	Unit	Notes
Access time	t _{acc}	_	20	μs	Figure 6.9
CE output delay time	t _{ce}	_	150	ns	Figure 6.10
OE output delay time	t _{oe}	_	150	ns	
Output disable delay time	t _{df}	_	100	ns	
Data output hold time	t _{oh}	5	_	ns	

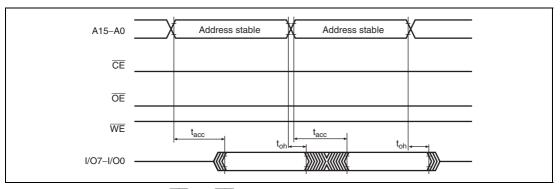


Figure 6.9 $\overline{\text{CE}}$ and $\overline{\text{OE}}$ Enable State Read Timing Waveforms

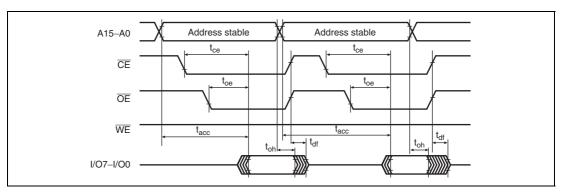


Figure 6.10 $\overline{\text{CE}}$ and $\overline{\text{OE}}$ Clock System Read Timing Waveforms

6.7.4 Auto-Program Mode

- 1. When reprogramming previously programmed addresses, perform auto-erasing before autoprogramming.
- 2. Perform auto-programming once only on the same address block. It is not possible to program an address block that has already been programmed.
- 3. In auto-program mode, 128 bytes are programmed simultaneously. This should be carried out by executing 128 consecutive byte transfers. A 128-byte data transfer is necessary even when programming fewer than 128 bytes. In this case, H'FF data must be written to the extra addresses.
- 4. The lower 7 bits of the transfer address must be low. If a value other than an effective address is input, processing will switch to a memory write operation but a write error will be flagged.
- 5. Memory address transfer is performed in the second cycle (figure 6.11). Do not perform transfer after the third cycle.
- 6. Do not perform a command write during a programming operation.
- 7. Perform one auto-program operation for a 128-byte block for each address. Two or more additional programming operations cannot be performed on a previously programmed address block.
- 8. Confirm normal end of auto-programming by checking I/O6. Alternatively, status read mode can also be used for this purpose (I/O7 status polling uses the auto-program operation end decision pin).
- Status polling I/O6 and I/O7 pin information is retained until the next command write. As long
 as the next command write has not been performed, reading is possible by enabling \(\overline{CE}\) and
 \(\overline{OE}\).
- 10. Table 6.13 shows the AC characteristics.

Table 6.13 AC Characteristics in Auto-Program Mode

Conditions: $V_{cc} = 3.3 \text{ V} \pm 0.3 \text{ V}, V_{ss} = 0 \text{ V}, T_a = 25^{\circ}\text{C} \pm 5^{\circ}\text{C}$

Item	Symbol	Min	Max	Unit	Notes
Command write cycle	t _{nxtc}	20	_	μs	Figure 6.11
CE hold time	t _{ceh}	0	_	ns	_
CE setup time	t _{ces}	0	_	ns	_
Data hold time	t _{dh}	50	_	ns	_
Data setup time	t _{ds}	50	_	ns	_
Write pulse width	t _{wep}	70	_	ns	_
Status polling start time	t _{wsts}	1	_	ms	_
Status polling access time	t _{spa}	_	150	ns	_
Address setup time	t _{as}	0	_	ns	_
Address hold time	t _{ah}	60	_	ns	_
Memory write time	t _{write}	1	3000	ms	_
WE rise time	t _r	_	30	ns	_
WE fall time	t,	_	30	ns	_

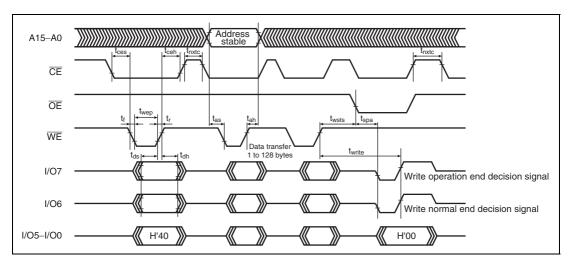


Figure 6.11 Auto-Program Mode Timing Waveforms

6.7.5 Auto-Erase Mode

- 1. Auto-erase mode supports only entire memory erasing.
- 2. Do not perform a command write during auto-erasing.
- 3. Confirm normal end of auto-erasing by checking I/O6. Alternatively, status read mode can also be used for this purpose (I/O7 status polling uses the auto-erase operation end decision pin).
- Status polling I/O6 and I/O7 pin information is retained until the next command write. As long
 as the next command write has not been performed, reading is possible by enabling \(\overline{CE}\) and
 \(\overline{OE}\).
- 5. Table 6.14 shows the AC characteristics.

Table 6.14 AC Characteristics in Auto-Erase Mode

Conditions: $V_{cc} = 3.3 \text{ V} \pm 0.3 \text{ V}, V_{ss} = 0 \text{ V}, T_a = 25^{\circ}\text{C} \pm 5^{\circ}\text{C}$

Item	Symbol	Min	Max	Unit	Notes
Command write cycle	t _{nxtc}	20	_	μs	Figure 6.12
CE hold time	t _{ceh}	0	_	ns	_
CE setup time	t _{ces}	0	_	ns	_
Data hold time	t _{dh}	50	_	ns	_
Data setup time	t _{ds}	50	_	ns	_
Write pulse width	t _{wep}	70	_	ns	_
Status polling start time	t _{ests}	1	_	ms	_
Status polling access time	t _{spa}	_	150	ns	_
Memory erase time	t _{erase}	100	40000	ms	_
WE rise time	t,	_	30	ns	<u> </u>
WE fall time	t,	_	30	ns	_

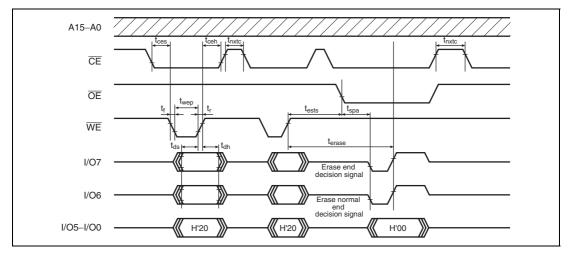


Figure 6.12 Auto-Erase Mode Timing Waveforms

6.7.6 Status Read Mode

- 1. Status read mode is provided to identify the kind of abnormal end. Use this mode when an abnormal end occurs in auto-program mode or auto-erase mode.
- 2. The return code is retained until a command write other than a status read mode command write is executed.
- 3. Table 6.15 shows the AC characteristics and 6.16 shows the return codes.

Table 6.15 AC Characteristics in Status Read Mode

Conditions: $V_{cc} = 3.3 \text{ V} \pm 0.3 \text{ V}, V_{ss} = 0 \text{ V}, T_a = 25^{\circ}\text{C} \pm 5^{\circ}\text{C}$

Item	Symbol	Min	Max	Unit	Notes
Read time after command write	t _{nxtc}	20	_	μs	Figure 6.13
CE hold time	t _{ceh}	0	_	ns	
CE setup time	t _{ces}	0	_	ns	
Data hold time	t _{dh}	50	_	ns	
Data setup time	t _{ds}	50	_	ns	
Write pulse width	t _{wep}	70	_	ns	
OE output delay time	t _{oe}	_	150	ns	
Disable delay time	t _{df}	_	100	ns	
CE output delay time	t _{ce}	_	150	ns	
WE rise time	t,	_	30	ns	
WE fall time	t,	_	30	ns	

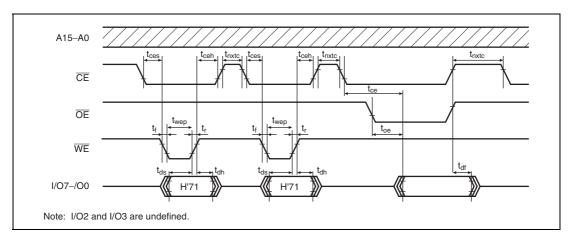


Figure 6.13 Status Read Mode Timing Waveforms

Table 6.16 Status Read Mode Return Codes

Pin Name	Initial Value	Indications
I/O7	0	1: Abnormal end
		0: Normal end
I/O6	0	1: Command error
		0: Otherwise
I/O5	0	1: Programming error
		0: Otherwise
I/O4	0	1: Erasing error
		0: Otherwise
I/O3	0	_
I/O2	0	_
I/O1	0	1: Over counting of writing or erasing
		0: Otherwise
I/O0	0	1: Effective address error
		0: Otherwise

6.7.7 Status Polling

- 1. The I/O7 status polling flag indicates the operating status in auto-program/auto-erase mode.
- 2. The I/O6 status polling flag indicates a normal or abnormal end in auto-program/auto-erase mode.

Table 6.17 Status Polling Output Truth Table

I/O7	I/O6	I/O0 to 5	Status
0	0	0	During internal operation
1	0	0	Abnormal end
1	1	0	Normal end
0	1	0	_

6.7.8 Programmer Mode Transition Time

Commands cannot be accepted during the oscillation stabilization period or the programmer mode setup period. After the programmer mode setup time, a transition is made to memory read mode.

Table 6.18 Stipulated Transition Times to Command Wait State

Item	Symbol	Min	Max	Unit	Notes
Oscillation stabilization time (crystal oscillator)	T _{osc1}	10	_	ms	Figure 6.14
Oscillation stabilization time (ceramic oscillator)	T _{osc1}	5	_	ms	
Programmer mode setup time	T _{bmv}	10	_	ms	
Vcc hold time	T _{dwn}	0	_	ms	

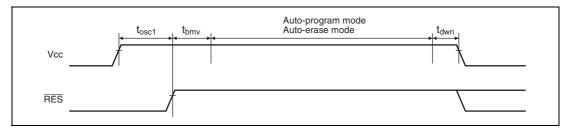


Figure 6.14 Oscillation Stabilization Time, Boot Program Transfer Time, and Power-Down Sequence

6.7.9 Notes on Memory Programming

- 1. When performing programming using programmer mode on a chip that has been programmed/erased in an on-board programming mode, auto-erasing is recommended before carrying out auto-programming.
- 2. The flash memory is initially in the erased state when the device is shipped by Renesas Technology. For other chips for which the erasure history is unknown, it is recommended that auto-erasing be executed to check and supplement the initialization (erase) level.

6.8 Power-Down States for Flash Memory

In user mode, the flash memory will operate in either of the following states:

- Normal operating mode
 The flash memory can be read and written to at high speed.
- Power-down operating mode
 The power supply circuit of the flash memory is partly halted and can be read under low power consumption.
- Standby mode
 All flash memory circuits are halted.

Table 6.19 shows the correspondence between the operating modes of this LSI and the flash memory. In subactive mode, the flash memory can be set to operate in power-down mode with the PDWND bit in FLPWCR. When the flash memory returns to its normal operating state from power-down mode or standby mode, a period to stabilize the power supply circuits that were stopped is needed. When the flash memory returns to its normal operating state, bits STS2 to STS0 in SYSCR1 must be set to provide a wait time of at least 20 μ s, even when the external clock is being used.

Table 6.19 Flash Memory Operating States

Flash	Memory	Operating	State
гіазіі	MEHIOIA	Operannu	State

LSI Operating State	PDWND = 0 (Initial value)	PDWND = 1
Active mode	Normal operating mode	Normal operating mode
Subactive mode	Power-down mode	Normal operating mode
Sleep mode	Normal operating mode	Normal operating mode
Subsleep mode	Standby mode	Standby mode
Standby mode	Standby mode	Standby mode
Watch mode	Standby mode	Standby mode

Section 7 RAM

7.1 Overview

The H8/38532 and H8/38533 have 1 Kbyte of high-speed static RAM on-chip, and the H8/38534, H8/38535, H8/38536, and H8/38537 have 2 Kbytes. The RAM is connected to the CPU by a 16-bit data bus, allowing high-speed 2-state access for both byte data and word data.

7.1.1 Block Diagram

Figure 7.1 shows a block diagram of the on-chip RAM.

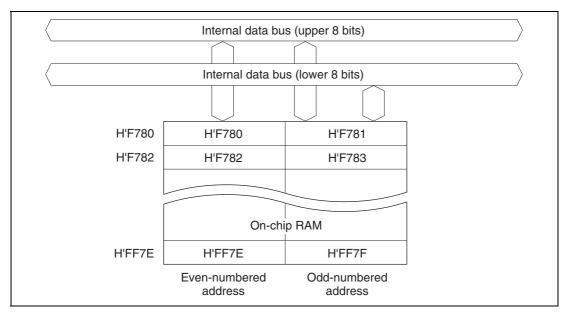


Figure 7.1 RAM Block Diagram (H8/38534)

Section 8 I/O Ports

8.1 Overview

The LSI is provided with six 8-bit I/O ports, one 4-bit I/O port, one 3-bit I/O port, one 8-bit input-only port, and one 1-bit input-only port. Table 8.1 indicates the functions of each port.

Each port has of a port control register (PCR) that controls input and output, and a port data register (PDR) for storing output data. Input or output can be assigned to individual bits. See section 2.8.3, Bit-Manipulation Instruction, for information on executing bit-manipulation instructions to write data in PCR or PDR.

Ports 5, 6, 7, 8, and A are also used as liquid crystal display segment and common pins, selectable in 8-bit units.

Block diagrams of each port are given in Appendix B, I/O Port Block Diagrams.

Table 8.1 Port Functions

Port	Description	Pins	Other Functions	Function Switching Registers
Port 1	8-bit I/O port MOS input pull-up option	P1, to P1,/ IRQ, to IRQ,/ TMIF, TMIC	External interrupts 3 to 1 Timer event interrupts TMIF, TMIC	PMR1 TCRF, TMC
		P1 ₄ /IRQ ₄ /ADTRG	External interrupt 4 and A/D converter external trigger	PMR1, AMR
		P1 ₃ /TMIG	Timer G input capture input	PMR1
		P1 ₂ , P1 ₁ / TMOFH, TMOFL	Timer F output compare output	PMR1
		P1 ₀ /TMOW	Timer A clock output	PMR1
Port 3	8-bit I/O portMOS input pull-up optionLarge-current port	P3,/AEVL P3,/AEVH P3,/TXD ₃₁ P3,/RXD ₃₁ P3,/SCK ₃₁	SCI31 data output (TXD ₃₁), data input (RXD ₃₁), clock input/output (SCK ₃₁), and asynchronous counter event inputs AEVL, AEVH	PMR3 SCR31 SMR31
		P3 ₂ P3 ₁ /UD/EXCL P3 ₀ /PWM	Timer C count-up/down select input, 14-bit PWM output, and external subclock input	

Port	Description	Pins	Other Functions	Function Switching Registers
Port 4	1-bit input port	P4₃/IRQ₀	External interrupt 0	PMR3
	• 3-bit I/O port	P4 ₂ /TXD ₃₂ P4 ₁ /RXD ₃₂ P4 ₀ /SCK ₃₂	SCI32 data output (TXD ₃₂), data input (RXD ₃₂), clock input/output (SCK ₃₂)	SCR32 SMR32
Port 5	8-bit I/O port	P5, to P5,/	Wakeup input (WKP, to	PMR5
	 MOS input pull-up option 	WKP, to WKP,/ SEG, to SEG,	\overline{WKP}_0), segment output (SEG ₈ to SEG ₁)	LPCR
Port 6			Segment output (SEG ₁₆ to	LPCR
	 MOS input pull-up option 	SEG ₁₆ to SEG ₉	SEG ₉)	
Port 7	• 8-bit I/O port	P7, to P7,/ SEG ₂₄ to SEG ₁₇	Segment output (SEG $_{24}$ to SEG $_{17}$)	LPCR
Port 8	• 8-bit I/O port	P8/SEG ₃₂ P8/SEG ₃₁ P8/SEG ₃₀ P8/SEG ₂₉ P8, to P8/ SEG ₂₈ to SEG ₂₅	Segment output (SEG ₃₂ to SEG ₂₅)	LPCR
Port A	4-bit I/O port	PA ₃ to PA ₀ / COM ₄ to COM ₁	Common output (COM ₄ to COM ₁)	LPCR
Port B	8-bit input port	PB ₇ to PB ₀ / AN ₇ to AN ₀	A/D converter analog input	AMR



8.2 Port 1

8.2.1 Overview

Port 1 is a 8-bit I/O port. Figure 8.1 shows its pin configuration.

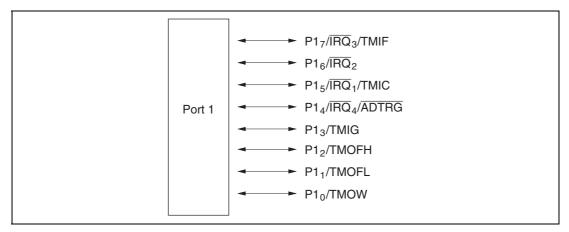


Figure 8.1 Port 1 Pin Configuration

8.2.2 Register Configuration and Description

Table 8.2 shows the port 1 register configuration.

Table 8.2 Port 1 Registers

Name	Abbr.	R/W	Initial Value	Address
Port data register 1	PDR1	R/W	H'00	H'FFD4
Port control register 1	PCR1	W	H'00	H'FFE4
Port pull-up control register 1	PUCR1	R/W	H'00	H'FFE0
Port mode register 1	PMR1	R/W	H'00	H'FFC8

(1) Port Data Register 1 (PDR1)

Bit	7	6	5	4	3	2	1	0
	P1 ₇	P1 ₆	P1 ₅	P1 ₄	P1 ₃	P1 ₂	P1 ₁	P1 ₀
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W							

PDR1 is an 8-bit register that stores data for port 1 pins P1, to P1₀. If port 1 is read while PCR1 bits are set to 1, the values stored in PDR1 are read, regardless of the actual pin states. If port 1 is read while PCR1 bits are cleared to 0, the pin states are read.

Upon reset, PDR1 is initialized to H'00.

(2) Port Control Register 1 (PCR1)

Bit	7	6	5	4	3	2	1	0
	PCR1 ₇	PCR1 ₆	PCR1 ₅	PCR1 ₄	PCR1 ₃	PCR1 ₂	PCR1 ₁	PCR1 ₀
Initial value	0	0	0	0	0	0	0	0
Read/Write	W	W	W	W	W	W	W	W

PCR1 is an 8-bit register for controlling whether each of the port 1 pins $P1_7$ to $P1_0$ functions as an input pin or output pin. Setting a PCR1 bit to 1 makes the corresponding pin an output pin, while clearing the bit to 0 makes the pin an input pin. The settings in PCR1 and in PDR1 are valid only when the corresponding pin is designated in PMR1 as a general I/O pin.

Upon reset, PCR1 is initialized to H'00.

PCR1 is a write-only register, which is always read as all 1s.

(3) Port Pull-Up Control Register 1 (PUCR1)

Bit	7	6	5	4	3	2	1	0
	PUCR1 ₇	PUCR1 ₆	PUCR1 ₅	PUCR1 ₄	PUCR1 ₃	PUCR1 ₂	PUCR1 ₁	PUCR1 ₀
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W							



PUCR1 controls whether the MOS pull-up of each of the port 1 pins P1₇ to P1₀ is on or off. When a PCR1 bit is cleared to 0, setting the corresponding PUCR1 bit to 1 turns on the MOS pull-up for the corresponding pin, while clearing the bit to 0 turns off the MOS pull-up.

Upon reset, PUCR1 is initialized to H'00.

(4) Port Mode Register 1 (PMR1)

Bit	7	6	5	4	3	2	1	0
	IRQ3	IRQ2	IRQ1	IRQ4	TMIG	TMOFH	TMOFL	TMOW
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W

PMR1 is an 8-bit read/write register, controlling the selection of pin functions for port 1 pins.

Upon reset, PMR1 is initialized to H'00.

Bit 7: P1,/IRQ,/TMIF pin function switch (IRQ3)

This bit selects whether pin P1/ \overline{IRQ}_3 /TMIF is used as P1, or as \overline{IRQ}_3 /TMIF.

Bit 7 IRQ3	Description	
0	Functions as P1, I/O pin	(initial value)
1	Functions as IRQ ₃ /TMIF input pin	

Note: Rising or falling edge sensing can be designated for IRQ₃, TMIF. For details on TMIF settings, see (3) Timer Control Register F (TCRF) in section 9.4.2, Register Descriptions.

Bit 6: $P1_6/\overline{IRQ}_2$ pin function switch (IRQ2)

This bit selects whether pin $P1_6/\overline{IRQ}_2$ is used as $P1_6$ or as \overline{IRQ}_2 .

Bit 6 IRQ2	Description	
0	Functions as P1 ₆ I/O pin	(initial value)
1	Functions as $\overline{IRQ}_{_2}$ input pin	

Note: Rising or falling edge sensing can be designated for \overline{IRQ}_2 .

Bit 5: P1,/IRQ,/TMIC pin function switch (IRQ1)

This bit selects whether pin P1₅/ \overline{IRQ}_1 /TMIC is used as P1₅ or as \overline{IRQ}_1 /TMIC.

Bit 5		
IRQ1	Description	
0	Functions as P1 ₅ I/O pin	(initial value)
1	Functions as IRQ,/TMIC input pin	

Note: Rising or falling edge sensing can be designated for $\overline{IRQ}_1/TMIC$.

For details of TMIC pin setting, see (1) Timer Mode Register C (TMC) in section 9.3.2, Register Descriptions.

Bit 4: P1₄/IRQ₄/ADTRG pin function switch (IRQ4)

This bit selects whether pin P1₄/IRQ₄/ADTRG is used as P1₄ or as IRQ₄/ADTRG.

Bit 4		
IRQ4	Description	
0	Functions as P1, I/O pin	(initial value)
1	Functions as IRQ₄/ADTRG input pin	

Note: For details of ADTRG pin setting, see section 12.3.2, Start of A/D Conversion by External Trigger Input.

Bit 3: P1₃/TMIG pin function switch (TMIG)

This bit selects whether pin P1₃/TMIG is used as P1₃ or as TMIG.

Bit 3 TMIG	Description	
0	Functions as P1 ₃ I/O pin	(initial value)
1	Functions as TMIG input pin	

Bit 2: P1₂/TMOFH pin function switch (TMOFH)

This bit selects whether pin P1₂/TMOFH is used as P1₂ or as TMOFH.

Bit 1: P1₁/TMOFL pin function switch (TMOFL)

This bit selects whether pin P1₁/TMOFL is used as P1₁ or as TMOFL.

Bit 1 TMOFL	Description	
0	Functions as P1, I/O pin	(initial value)
1	Functions as TMOFL output pin	

Bit 0: P1₀/TMOW pin function switch (TMOW)

This bit selects whether pin P1₀/TMOW is used as P1₀ or as TMOW.

Bit 0 TMOW	Description	
0	Functions as P1 ₀ I/O pin	(initial value)
1	Functions as TMOW output pin	

8.2.3 **Pin Functions**

Table 8.3 shows the port 1 pin functions.

Table 8.3 Port 1 Pin Functions

Pin **Pin Functions and Selection Method**

P1/IRQ./TMIF The pin function depends on bit IRQ3 in PMR1, bits CKSL2 to CKSL0 in TCRF, and bit PCR1, in PCR1.

IRQ ₃	0		1	
PCR1 ₇	0	1	*	*
CKSL2 to CKSL0	,	*	Not 0**	0**
Pin function	P1, input pin	P1, output pin	ĪRQ₃ input pin	ĪRQ₃/TMIF input pin

Note: When this pin is used as the TMIF input pin, clear bit IEN3 to 0 in IENR1 to disable the IRQ₃ interrupt.

P1 /IRQ,

The pin function depends on bits IRQ2 in PMR1 and bit PCR1, in PCR1.

IRQ2	()	1
PCR1 ₆	0	1	*
Pin function	P1 ₆ input pin	P1 ₆ output pin	$\overline{IRQ}_{\scriptscriptstyle 2}$ input pin

P1_e/IRQ_e **TMIC**

The pin function depends on bit IRQ1 in PMR1, bits TMC2 to TMC0 in TMC, and bit PCR1, in PCR1.

IRQ1)	1		
PCR1₅	0	1	*		
TMC2 to TMC0	:	*	Not 111	111	
Pin function	P1 _s input pin P1 _s output pin		ĪRQ₁ input pin	ĪRQ₁/TMIC input pin	

Note: When this pin is used as the TMIC input pin, clear bit IEN1 to 0 in IENR1 to disable the IRQ, interrupt.

Pin	Pin Functions and Selection Method					
P1 ₄ /IRQ ₄ ADTRG	The pin function do in PCR1.	epends on bit IR	Q4 in PMR1, bit	TRGE in AMR,	and bit PCR1 ₄	
	IRQ4	(0	-	1	
	PCR1₄	0	1	;	*	
	TRGE	:	*	0	1	
	Pin function	P1₄ input pin	P1 ₄ output pin	ĪRQ₄ input pin	IRQ₄/ADTRG input pin	
	Note: When this p to disable th	in is used as the e IRQ₄ interrupt		oin, clear bit IEN	4 to 0 in IENR1	
P1 ₃ /TMIG	The pin function de	epends on bit TN	MIG in PMR1 an	d bit PCR1 ₃ in P	CR1.	
	TMIG	(0	-	1	
	PCR1 ₃	0	1	>	*	
	Pin function	P1 ₃ input pin	P1 ₃ output pin	TMIG in	nput pin	
P1 ₂ /TMOFH	The pin function de	epends on bit TN	MOFH in PMR1	and bit PCR1 ₂ in	PCR1.	
	TMOFH	(0	1		
	PCR1 ₂	0	1	*		
	Pin function	P1 ₂ input pin	P1 ₂ output pin	TMOFH	output pin	
P1₁/TMOFL	The pin function de	epends on bit TN	MOFL in PMR1 a	and bit PCR1, in	PCR1.	
	TMOFL	(0	-	1	
	PCR1,	0	1	>	*	
	Pin function	P1₁ input pin	P1, output pin	TMOFL o	output pin	
P1 ₀ /TMOW	The pin function de	epends on bit TN	MOW in PMR1 a	nd bit PCR1 ₀ in	PCR1.	
	TMOW	(0		1	
	PCR1₀	0	1	;	*	
	Pin function	P1 _o input pin	P1 _o output pin	TMOW o	output pin	

*: Don't care

8.2.4 Pin States

Table 8.4 shows the port 1 pin states in each operating mode.

Table 8.4 Port 1 Pin States

Pins	Reset	Sleep	Subsleep	Standby	Watch	Subactive	Active
P1/IRQ_/TMIF P1_/IRQ_ P1_/IRQ_/TMIC P1_/IRQ_/ADTRG P1_/TMIG P1_/TMOFH P1_/TMOFL P1_/TMOW	High- impedance	Retains previous state	Retains previous state	High- impedance*	Retains previous state	Functional	Functional

Note: * A high-level signal is output when the MOS pull-up is in the on state.

8.2.5 MOS Input Pull-Up

Port 1 has a built-in MOS input pull-up function that can be controlled by software. When a PCR1 bit is cleared to 0, setting the corresponding PUCR1 bit to 1 turns on the MOS input pull-up for that pin. The MOS input pull-up function is in the off state after a reset.

PCR1 _n	0	0	1	
PUCR1 _n	0	1	*	
MOS input pull-up	Off	On	Off	

(n = 7 to 0)

*: Don't care

8.3 Port 3

8.3.1 Overview

Port 3 is a 8-bit I/O port, configured as shown in figure 8.2.

In the flash memory version, the on-chip pull-up MOS for pin P3₂ is on during the reset period. It turns off and normal operation resumes after the reset is cleared. This should be considered when making connections to external circuitry. Note that in the mask ROM version P3₂ continues to operate normally.

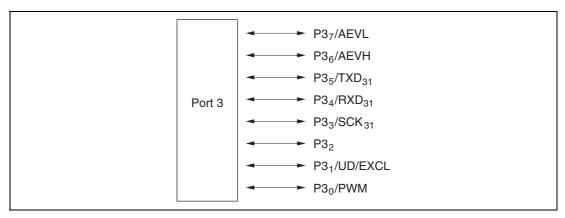


Figure 8.2 Port 3 Pin Configuration

8.3.2 Register Configuration and Description

Table 8.5 shows the port 3 register configuration.

Table 8.5 Port 3 Registers

Name	Abbr.	R/W	Initial Value	Address
Port data register 3	PDR3	R/W	H'00	H'FFD6
Port control register 3	PCR3	W	H'00	H'FFE6
Port pull-up control register 3	PUCR3	R/W	H'00	H'FFE1
Port mode register 2	PMR2	R/W	H'58	H'FFC9
Port mode register 3	PMR3	R/W	H'04	H'FFCA

(1) Port Data Register 3 (PDR3)

Bit	7	6	5	4	3	2	1	0
	P3 ₇	P3 ₆	P3 ₅	P3 ₄	P3 ₃	P3 ₂	P3 ₁	P3 ₀
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W							

PDR3 is an 8-bit register that stores data for port 3 pins P3₇ to P3₀. If port 3 is read while PCR3 bits are set to 1, the values stored in PDR3 are read, regardless of the actual pin states. If port 3 is read while PCR3 bits are cleared to 0, the pin states are read.

Upon reset, PDR3 is initialized to H'00.

(2) Port Control Register 3 (PCR3)

Bit	7	6	5	4	3	2	1	0
	PCR3 ₇	PCR3 ₆	PCR3 ₅	PCR3 ₄	PCR3 ₃	PCR3 ₂	PCR3 ₁	PCR3 ₀
Initial value	0	0	0	0	0	0	0	0
Read/Write	W	W	W	W	W	W	W	W

PCR3 is an 8-bit register for controlling whether each of the port 3 pins $P3_7$ to $P3_0$ functions as an input pin or output pin. Setting a PCR3 bit to 1 makes the corresponding pin an output pin, while clearing the bit to 0 makes the pin an input pin. The settings in PCR3 and in PDR3 are valid only when the corresponding pin is designated in PMR3 as a general I/O pin.

Upon reset, PCR3 is initialized to H'00.

PCR3 is a write-only register, which is always read as all 1s.



(3) Port Pull-Up Control Register 3 (PUCR3)

Bit	7	6	5	4	3	2	1	0
	PUCR3 ₇	PUCR3 ₆	PUCR3 ₅	PUCR3 ₄	PUCR3 ₃	PUCR3 ₂	PUCR3 ₁	PUCR3 ₀
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W							

PUCR3 controls whether the MOS pull-up of each of the port 3 pins $P3_7$ to $P3_0$ is on or off. When a PCR3 bit is cleared to 0, setting the corresponding PUCR3 bit to 1 turns on the MOS pull-up for the corresponding pin, while clearing the bit to 0 turns off the MOS pull-up.

Upon reset, PUCR3 is initialized to H'00.

(4) Port Mode Register 3 (PMR3)

Bit	7	6	5	4	3	2	1	0
	AEVL	AEVH	WDCKS	NCS	IRQ0	_	UD	PWM
Initial value	0	0	0	0	0	1	0	0
Read/Write	R/W	R/W	R/W	R/W	R/W	_	R/W	R/W

PMR3 is an 8-bit read/write register, controlling the selection of pin functions for port 3 pins.

Upon reset, PMR3 is initialized to H'04.

Bit 7: P3₇/AEVL pin function switch (AEVL)

This bit selects whether pin P3₇/AEVL is used as P3₇ or as AEVL.

Bit 7 AEVL	Description	
0	Functions as P3, I/O pin	(initial value)
1	Functions as AEVL input pin	

Bit 6: P3₆/AEVH pin function switch (AEVH)

This bit selects whether pin P3₆/AEVH is used as P3₆ or as AEVH.

Bit 6 AEVH	Description	
0	Functions as P3 ₆ I/O pin	(initial value)
1	Functions as AEVH input pin	

Bit 5: Watchdog timer source clock select (WDCKS)

This bit selects the watchdog timer source clock.

Bit 5 WDCKS	Description	
0	φ/8192 selected	(initial value)
1	φw/32 selected	

Bit 4: TMIG noise canceler select (NCS)

This bit controls the noise canceler for the input capture input signal (TMIG).

Bit 4 NCS	Description	
0	Noise cancellation function not used	(initial value)
1	Noise cancellation function used	

Bit 3: $P4_3/\overline{IRQ}_0$ pin function switch (IRQ₀)

This bit selects whether pin $P4_3/\overline{IRQ}_0$ is used as $P4_3$ or as \overline{IRQ}_0 .

Bit 3 IRQ ₀	Description	
0	Functions as P4 ₃ input pin	(initial value)
1	Functions as $\overline{\text{IRQ}}_{_0}$ input pin	

Bit 2: Reserved bit

This bit cannot be written to.

Bit 1: P3₁/UD pin function switch (UD)

This bit selects whether pin P3₁/UD is used as P3₁ or as UD.

Bit 1		
UD	Description	
0	Functions as P3, I/O pin	(initial value)
1	Functions as UD input pin	

Bit 0: P3₀/PWM pin function switch (PWM)

This bit selects whether pin P3₀/PWM is used as P3₀ or as PWM.

Bit 0		
PWM	Description	
0	Functions as P3 _o I/O pin	(initial value)
1	Functions as PWM output pin	

(5) Port Mode Register 2 (PMR2)

Bit	7	6	5	4	3	2	1	0	
	EXCL	_	_	_	_	_	_	_	
Initial value	0	1	0	1	1	0	0	0	_
Read/Write	R/W	R	R/W	R	R	R/W	R/W	R/W	

PMR2 is an 8-bit read/write register that controls external clock input to pin P3₁.

Upon reset, PMR2 is initialized to H'58.

Bit 7: P3₁/UD/EXCL pin function switch (EXCL)

This bit selects whether pin P3₁/UD/EXCL is used as P3₁/UD or as EXCL. When the pin is used as EXCL an external clock should be input to it. See section 4, Clock Pulse Generators, for a connection example.

Bit 7		
EXCL	Description	
0	Functions as P3,/UD I/O pin	(initial value)
1	Functions as EXCL input pin	

Bit 6: Reserved bit

Bit 6 is a reserved bit. It is always read as 1 and cannot be modified.

Bit 5: Reserved bit

Bit 5 is a readable/writable reserved bit.

Bits 4 and 3: Reserved bits

Bits 4 and 3 are reserved bits. They are always read as 1 and cannot be modified.

Bits 2 to 0: Reserved bits

Bits 2 to 0 are readable/writable reserved bits.



8.3.3 Pin Functions

Table 8.6 shows the port 3 pin functions.

Table 8.6 Port 3 Pin Functions

Pin	Pin Functions and	Selection Method		
P3 ₇ /AEVL	The pin function depends on bit SO1 in PMR3 and bit PCR3, in PCR3.			
	AEVL	()	1
	PCR3 ₇	0	1	*
	Pin function	P3, input pin	P3, output pin	AEVL input pin
P3 ₆ /AEVH	The pin function dep	pends on bit AEVH in	PMR3 and bit PCR	3 ₆ in PCR3.
	AEVH	()	1
	PCR3 ₆	0	1	*
	Pin function	P3 ₆ input pin	P3 ₆ output pin	AEVH input pin
P3 ₅ /TXD ₃₁	The pin function dep PCR3 ₅ in PCR3.	pends on bit TE ₃₁ in S	SCR31, bit SPC31 in	SPCR, and bit
	SPC31	()	1
	TE ₃₁	()	1
	PCR3₅	0	1	*
	Pin function	P3₅ input pin	P3 ₅ output pin	TXD ₃₁ output pin
P3 ₄ /RXD ₃₁	The pin function depends on bit RE ₃₁ in SCR31 and bit PCR3 ₄ in PCR3.			
	RE ₃₁	()	1
	PCR3₄	0	1	*
	Pin function	P3₄ input pin	P3₄ output pin	RXD ₃₁ input pin
			_	_

Pin	Pin Functions and Selection	Method
T III I	FIII FUIICIIOIIS AIIU SEIECIIOI	ı welliou

P3₃/SCK₃₁

The pin function depends on bits CKE311 and CKE310 in SCR31, bit COM31 in SMR31, and bit PCR3 $_{\rm s}$ in PCR3.

CKE311	0			1	
CKE310	0			1	*
COM3 ₁	()	1	*	*
PCR3 ₃	0	1	*	ķ	*
Pin function	P3 ₃ input pin	P3 ₃ output pin	SCK ₃₁		SCK ₃₁ input pin

P3,

The pin function depends on bit $PCR3_2$ in PCR3.

PCR3 ₂	0	1
Pin function	P3 ₂ input pin	P3 ₂ output pin

P3,/UD/EXCL

The pin function depends on bit EXCL in PMR2, bit UD in PMR3, and bit PCR3, in PCR3.

EXCL		1		
UD	()	1	*
PCR3 ₁	0	1	*	*
Pin function	P3, input pin	P3, output pin	UD input pin	EXCL input pin

P3₀/PWM

The pin function depends on bit PWM in PMR3 and bit PCR3_o in PCR3.

PWM	(0		
PCR3₀	0	1	*	
Pin function	P3 _o input pin	P3₀ output pin	PWM output pin	

*: Don't care



8.3.4 Pin States

Table 8.7 shows the port 3 pin states in each operating mode.

Table 8.7 Port 3 Pin States

Pins	Reset	Sleep	Subsleep	Standby	Watch	Subactive	Active
P3/AEVL P3/AEVH P3/TXD ₃₁ P3/RXD ₃₁ P3/SCK ₃₁	High- impedance	Retains previous state	Retains previous state	High- impedance*1	Retains previous state	Functional	Functional
P3 ₂ *3	Pull-up MOS on	_					
P3 ₂ *2 P3 ₁ /UD/EXCL P3 ₀ /PWM	High- impedance	_					

Notes: 1. A high-level signal is output when the MOS pull-up is in the on state.

- 2. Applies to the mask ROM version.
- 3. Applies to the flash memory version.

8.3.5 MOS Input Pull-Up

Port 3 has a built-in MOS input pull-up function that can be controlled by software. When a PCR3 bit is cleared to 0, setting the corresponding PUCR3 bit to 1 turns on the MOS pull-up for that pin. The MOS pull-up function is in the off state after a reset.

PCR3 _n	0	0	1
PUCR3 _n	0	1	*
MOS input pull-up	Off	On	Off

(n = 7 to 0)

*: Don't care

8.4 Port 4

8.4.1 Overview

Port 4 is a 3-bit I/O port and 1-bit input port, configured as shown in figure 8.3.

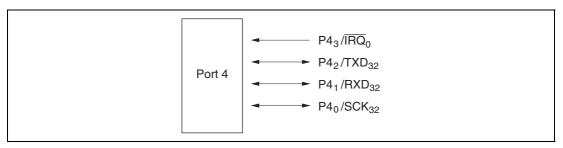


Figure 8.3 Port 4 Pin Configuration

8.4.2 Register Configuration and Description

Table 8.8 shows the port 4 register configuration.

Table 8.8 Port 4 Registers

Name	Abbr.	R/W	Initial Value	Address
Port data register 4	PDR4	R/W	H'F8	H'FFD7
Port control register 4	PCR4	W	H'F8	H'FFE7

(1) Port Data Register 4 (PDR4)

Bit	7	6	5	4	3	2	1	0
	_	_	_	_	P4 ₃	P4 ₂	P4 ₁	P4 ₀
Initial value	1	1	1	1	1	0	0	0
Read/Write	_	_	_	_	R	R/W	R/W	R/W

PDR4 is an 8-bit register that stores data for port 4 pins P4₂ to P4₀. If port 4 is read while PCR4 bits are set to 1, the values stored in PDR4 are read, regardless of the actual pin states. If port 4 is read while PCR4 bits are cleared to 0, the pin states are read.

Upon reset, PDR4 is initialized to H'F8.

(2) Port Control Register 4 (PCR4)

Bit	7	6	5	4	3	2	1	0
	_	_	_	_	_	PCR4 ₂	PCR4 ₁	PCR4 ₀
Initial value	1	1	1	1	1	0	0	0
Read/Write	_	_	_	_	_	W	W	W

PCR4 is an 8-bit register for controlling whether each of port 4 pins P4₂ to P4₀ functions as an input pin or output pin. Setting a PCR4 bit to 1 makes the corresponding pin an output pin, while clearing the bit to 0 makes the pin an input pin. PCR4 and PDR4 settings are valid when the corresponding pins are designated for general-purpose input/output by SCR32.

Upon reset, PCR4 is initialized to H'F8.

PCR4 is a write-only register, which always reads all 1s.

8.4.3 Pin Functions

Table 8.9 shows the port 4 pin functions.

Table 8.9 Port 4 Pin Functions

Pin	Pin Functions and Selection Method						
P4 ₃ /IRQ ₀	The pin function depends on bit IRQ0 in PMR3.						
	IRQ0	0		1			
	Pin function	P4 ₃ input pi	n	ĪR	IQ₀ input pin		
P4 ₂ /TXD ₃₂	The pin function depends on bit TE ₃₂ in SCR32, bit SPC32 in SPCR, and bit PCR4 ₂ in PCR4.						
	TE ₃₂	()		1		
	PCR4 ₂	0 1			*		
	Pin function P4 ₂ input pin P4 ₂ output pin						

Pin	Pin Functions and Selection Method							
P4 ₁ /RXD ₃₂	The pin function depends on bit RE ₃₂ in SCR32 and bit PCR4, in PCR4.							
	RE ₃₂		0)			1	
	PCR4 ₁	0		1			*	
	Pin function	P4₁ input pi	n	P4₁ out	put pin	RX	(D ₃₂ input pin	
P4 ₀ /SCK ₃₂	The pin function depends on bits CKE321 and CKE320 in SCR32, bit COM32 SMR32, and bit PCR4 ₀ in PCR4.						1 bit COM32 in	
	CKE321			0			1	
	CKE320		0			1	*	
	COM32	0			1	*	*	
	PCR4₀	0 1		1	*	•	*	
	Pin function	P4 ₀ input pin	P4 ₀ 0	output pin	ut pin SCK ₃₂ out		SCK ₃₂ input pin	
					Ρ'		F	

*: Don't care

8.4.4 Pin States

Table 8.10 shows the port 4 pin states in each operating mode.

Table 8.10 Port 4 Pin States

Pins	Reset	Sleep	Subsleep	Standby	Watch	Subactive	Active
P4 ₃ /IRQ ₀ P4 ₂ /TXD ₃₂ P4 ₁ /RXD ₃₂ P4 ₂ /SCK ₃₂	High- impedance	Retains previous state		High- impedance	Retains previous state	Functional	Functional

8.5 Port 5

8.5.1 Overview

Port 5 is an 8-bit I/O port, configured as shown in figure 8.4.

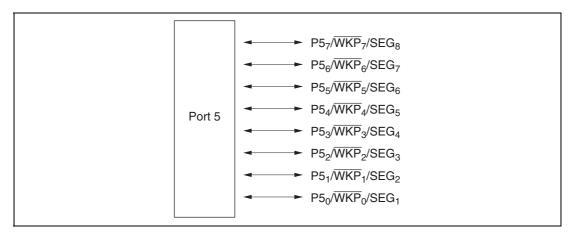


Figure 8.4 Port 5 Pin Configuration

8.5.2 Register Configuration and Description

Table 8.11 shows the port 5 register configuration.

Table 8.11 Port 5 Registers

Name	Abbr.	R/W	Initial Value	Address
Port data register 5	PDR5	R/W	H'00	H'FFD8
Port control register 5	PCR5	W	H'00	H'FFE8
Port pull-up control register 5	PUCR5	R/W	H'00	H'FFE2
Port mode register 5	PMR5	R/W	H'00	H'FFCC

(1) Port Data Register 5 (PDR5)

Bit	7	6	5	4	3	2	1	0
	P5 ₇	P5 ₆	P5 ₅	P5 ₄	P5 ₃	P5 ₂	P5 ₁	P5 ₀
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W							

PDR5 is an 8-bit register that stores data for port 5 pins P5, to P5₀. If port 5 is read while PCR5 bits are set to 1, the values stored in PDR5 are read, regardless of the actual pin states. If port 5 is read while PCR5 bits are cleared to 0, the pin states are read.

Upon reset, PDR5 is initialized to H'00.

(2) Port Control Register 5 (PCR5)

Bit	7	6	5	4	3	2	1	0
	PCR5 ₇	PCR5 ₆	PCR5 ₅	PCR5 ₄	PCR5 ₃	PCR5 ₂	PCR5 ₁	PCR5 ₀
Initial value	0	0	0	0	0	0	0	0
Read/Write	W	W	W	W	W	W	W	W

PCR5 is an 8-bit register for controlling whether each of the port 5 pins P5₇ to P5₀ functions as an input pin or output pin. Setting a PCR5 bit to 1 makes the corresponding pin an output pin, while clearing the bit to 0 makes the pin an input pin. PCR5 and PDR5 settings are valid when the corresponding pins are designated for general-purpose input/output by PMR5 and bits SGS3 to SGS0 in LPCR.

Upon reset, PCR5 is initialized to H'00.

PCR5 is a write-only register, which is always read as all 1s.

(3) Port Pull-Up Control Register 5 (PUCR5)

Bit	7	6	5	4	3	2	1	0
	PUCR5 ₇	PUCR5 ₆	PUCR5 ₅	PUCR5 ₄	PUCR5 ₃	PUCR5 ₂	PUCR5 ₁	PUCR5 ₀
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W							

PUCR5 controls whether the MOS pull-up of each of port 5 pins $P5_7$ to $P5_0$ is on or off. When a PCR5 bit is cleared to 0, setting the corresponding PUCR5 bit to 1 turns on the MOS pull-up for the corresponding pin, while clearing the bit to 0 turns off the MOS pull-up.

Upon reset, PUCR5 is initialized to H'00.

(4) Port Mode Register 5 (PMR5)

Bit	7	6	5	4	3	2	1	0
	WKP ₇	WKP ₆	WKP_5	WKP ₄	WKP ₃	WKP ₂	WKP ₁	WKP ₀
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W

PMR5 is an 8-bit read/write register, controlling the selection of pin functions for port 5 pins.

Upon reset, PMR5 is initialized to H'00.

Bit n: $P5_n/\overline{WKP}_n/SEG_{n+1}$ pin function switch (WKPn)

When pin P5n/WKPn/SEGn+1 is not used as SEG_{n+1}, these bits select whether the pin is used as P5n or \overline{WKP}_n .

Bit n WKPn	Description	
0	Functions as P5n I/O pin	(initial value)
1	Functions as WKP input pin	
		(n - 7 to 0)

Note: For use as SEG_{n+1} , see section 13.2.1, LCD Port Control Register (LPCR).

8.5.3 Pin Functions

Table 8.12 shows the port 5 pin functions.

Table 8.12 Port 5 Pin Functions

Pin	Pin Functions and S	Selection Method
1 111	i iii i aiictions ana t	Sciconon Michiga

 $P5_{/}\overline{WKP}_{/}SEG_{_8}$ The pin function depends on bit WKP $_{_n}$ in PMR5, bit PCR5 $_{_n}$ in PCR5, and bits SGS3 to SGS0 in LPCR.

P5₀/WKP₀/SEG₁

(n = 7 to 0)

SGS3 to SGS0		1***		
WKP _n	()	1	*
PCR5 _n	0	1	*	*
Pin function	P5 _n input pin	P5 _n output pin	WKPn input pin	SEGn+1 output pin

^{*:} Don't care

8.5.4 Pin States

Table 8.13 shows the port 5 pin states in each operating mode.

Table 8.13 Port 5 Pin States

Pins	Reset	Sleep	Subsleep	Standby	Watch	Subactive	Active
P5 ₇ /WKP ₇ / SEG ₈ to P5 ₀ / WKP ₀ /SEG ₁	High- impedance	Retains previous state		High- impedance*		Functional	Functional

Note: * A high-level signal is output when the MOS pull-up is in the on state.



8.5.5 MOS Input Pull-Up

Port 5 has a built-in MOS input pull-up function that can be controlled by software. When a PCR5 bit is cleared to 0, setting the corresponding PUCR5 bit to 1 turns on the MOS pull-up for that pin. The MOS pull-up function is in the off state after a reset.

PCR5 _n	0	0	1
PUCR5 _n	0	1	*
MOS input pull-up	Off	On	Off

(n = 7 to 0)

*: Don't care

8.6 Port 6

8.6.1 Overview

Port 6 is an 8-bit I/O port. The port 6 pin configuration is shown in figure 8.5.

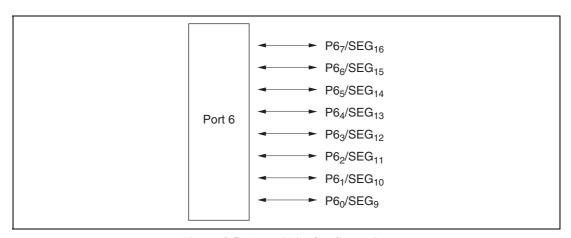


Figure 8.5 Port 6 Pin Configuration

8.6.2 Register Configuration and Description

Table 8.14 shows the port 6 register configuration.

Table 8.14 Port 6 Registers

Name	Abbr.	R/W	Initial Value	Address
Port data register 6	PDR6	R/W	H'00	H'FFD9
Port control register 6	PCR6	W	H'00	H'FFE9
Port pull-up control register 6	PUCR6	R/W	H'00	H'FFE3

(1) Port Data Register 6 (PDR6)

Bit	7	6	5	4	3	2	1	0
	P6 ₇	P6 ₆	P6 ₅	P6 ₄	P6 ₃	P6 ₂	P6 ₁	P6 ₀
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W							

PDR6 is an 8-bit register that stores data for port 6 pins P6, to P6.

If port 6 is read while PCR6 bits are set to 1, the values stored in PDR6 are read, regardless of the actual pin states. If port 6 is read while PCR6 bits are cleared to 0, the pin states are read.

Upon reset, PDR6 is initialized to H'00.

(2) Port Control Register 6 (PCR6)

Bit	7	6	5	4	3	2	1	0
	PCR6 ₇	PCR6 ₆	PCR6 ₅	PCR6 ₄	PCR6 ₃	PCR6 ₂	PCR6 ₁	PCR6 ₀
Initial value	0	0	0	0	0	0	0	0
Read/Write	W	W	W	W	W	W	W	W

PCR6 is an 8-bit register for controlling whether each of the port 6 pins $P6_7$ to $P6_0$ functions as an input pin or output pin.

Setting a PCR6 bit to 1 makes the corresponding pin (P6₇ to P6₀) an output pin, while clearing the bit to 0 makes the pin an input pin. PCR6 and PDR6 settings are valid when the corresponding pins are designated for general-purpose input/output by bits SGS3 to SGS0 in LPCR.



Upon reset, PCR6 is initialized to H'00.

PCR6 is a write-only register, which always reads all 1s.

(3) Port Pull-Up Control Register 6 (PUCR6)

Bit	7	6	5	4	3	2	1	0
	PUCR6 ₇	PUCR6 ₆	PUCR6 ₅	PUCR6 ₄	PUCR6 ₃	PUCR6 ₂	PUCR6 ₁	PUCR6 ₀
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W							

PUCR6 controls whether the MOS pull-up of each of the port 6 pins $P6_7$ to $P6_9$ is on or off. When a PCR6 bit is cleared to 0, setting the corresponding PUCR6 bit to 1 turns on the MOS pull-up for the corresponding pin, while clearing the bit to 0 turns off the MOS pull-up.

Upon reset, PUCR6 is initialized to H'00.

8.6.3 Pin Functions

Table 8.15 shows the port 6 pin functions.

Table 8.15 Port 6 Pin Functions

Pin Functions and Selection Method								
The pin function depends on bit PCR6n in PCR6 and bits SGS3 to SGS0 in LPCR. $(n=7\ to\ 0)$								
SGS3 to SGS0	00**,	011**, 1***						
PCR6 _n	0	1	*					
Pin function P6, input pin P6, output pin S								
	The pin function dep LPCR. SGS3 to SGS0 PCR6,	The pin function depends on bit PCR6n LPCR. SGS3 to SGS0 00**, PCR6, 0	The pin function depends on bit PCR6n in PCR6 and bits SG LPCR. SGS3 to SGS0 00**, 010* PCR6 _n 0 1					

*: Don't care

8.6.4 Pin States

Table 8.16 shows the port 6 pin states in each operating mode.

Table 8.16 Port 6 Pin States

Pin	Reset	Sleep	Subsleep	Standby	Watch	Subactive	Active
P6 ₇ /SEG ₁₆ to P6 ₀ /SEG ₉	High- impedance	Retains previous state		High- impedance*		Functional	Functional

Note: * A high-level signal is output when the MOS pull-up is in the on state.

8.6.5 MOS Input Pull-Up

Port 6 has a built-in MOS pull-up function that can be controlled by software. When a PCR6 bit is cleared to 0, setting the corresponding PUCR6 bit to 1 turns on the MOS pull-up for that pin. The MOS pull-up function is in the off state after a reset.

PCR6 _n	0	0	1
PUCR6 _n	0	1	*
MOS input pull-up	Off	On	Off

(n = 7 to 0)

*: Don't care

8.7 Port 7

8.7.1 Overview

Port 7 is a 8-bit I/O port, configured as shown in figure 8.6.

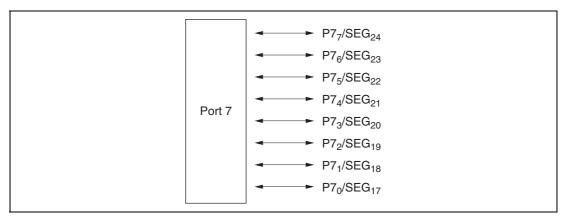


Figure 8.6 Port 7 Pin Configuration

8.7.2 Register Configuration and Description

Table 8.17 shows the port 7 register configuration.

Table 8.17 Port 7 Registers

Name	Abbr.	R/W	Initial Value	Address
Port data register 7	PDR7	R/W	H'00	H'FFDA
Port control register 7	PCR7	W	H'00	H'FFEA

(1) Port Data Register 7 (PDR7)

Bit	7	6	5	4	3	2	1	0
	P7 ₇	P7 ₆	P7 ₅	P7 ₄	P7 ₃	P7 ₂	P7 ₁	P7 ₀
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W							

PDR7 is an 8-bit register that stores data for port 7 pins $P7_7$ to $P7_0$. If port 7 is read while PCR7 bits are set to 1, the values stored in PDR7 are read, regardless of the actual pin states. If port 7 is read while PCR7 bits are cleared to 0, the pin states are read.

Upon reset, PDR7 is initialized to H'00.

(2) Port Control Register 7 (PCR7)

Bit	7	6	5	4	3	2	1	0
	PCR7 ₇	PCR7 ₆	PCR7 ₅	PCR7 ₄	PCR7 ₃	PCR7 ₂	PCR7 ₁	PCR7 ₀
Initial value	0	0	0	0	0	0	0	0
Read/Write	W	W	W	W	W	W	W	W

PCR7 is an 8-bit register for controlling whether each of the port 7 pins $P7_7$ to $P7_0$ functions as an input pin or output pin. Setting a PCR7 bit to 1 makes the corresponding pin an output pin, while clearing the bit to 0 makes the pin an input pin. PCR7 and PDR7 settings are valid when the corresponding pins are designated for general-purpose input/output by bits SGS3 to SGS0 in LPCR.

Upon reset, PCR7 is initialized to H'00.

PCR7 is a write-only register, which always reads as all 1s.



8.7.3 Pin Functions

Table 8.18 shows the port 7 pin functions.

Table 8.18 Port 7 Pin Functions

Pin	Pin Functions and Selection Method								
P7 ₇ /SEG ₂₄ to P7 ₀ /SEG ₁₇	The pin function dep LPCR.	pin function depends on bit PCR7 $_{_{n}}$ in PCR7 and bits SGS3 to SGS0 in R. $(n=7\ to\ 0)$							
	SGS3 to SGS0	00** 01**, 1***							
	PCR7 _n	0	*						
	Pin function	P7 _n input pin P7 _n output pin SEG _{n+17} output pin							

*: Don't care

8.7.4 Pin States

Table 8.19 shows the port 7 pin states in each operating mode.

Table 8.19 Port 7 Pin States

Pins	Reset	Sleep	Subsleep	Standby	Watch	Subactive	Active
P7 ₇ /SEG ₂₄ to P7 ₀ /SEG ₁₇	High- impedance	Retains previous state		High- impedance		Functional	Functional

8.8 Port 8

8.8.1 Overview

Port 8 is an 8-bit I/O port configured as shown in figure 8.7.

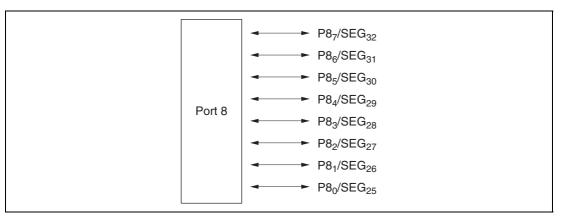


Figure 8.7 Port 8 Pin Configuration

8.8.2 Register Configuration and Description

Table 8.20 shows the port 8 register configuration.

Table 8.20 Port 8 Registers

Name	Abbr.	R/W	Initial Value	Address
Port data register 8	PDR8	R/W	H'00	H'FFDB
Port control register 8	PCR8	W	H'00	H'FFEB

(1) Port Data Register 8 (PDR8)

Bit	7	6	5	4	3	2	1	0
	P8 ₇	P8 ₆	P8 ₅	P8 ₄	P8 ₃	P8 ₂	P8 ₁	P8 ₀
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W							

PDR8 is an 8-bit register that stores data for port 8 pins P8, to P8₀. If port 8 is read while PCR8 bits are set to 1, the values stored in PDR8 are read, regardless of the actual pin states. If port 8 is read while PCR8 bits are cleared to 0, the pin states are read.

Upon reset, PDR8 is initialized to H'00.

(2) Port Control Register 8 (PCR8)

Bit	7	6	5	4	3	2	1	0
	PCR8 ₇	PCR8 ₆	PCR8 ₅	PCR8 ₄	PCR8 ₃	PCR8 ₂	PCR8 ₁	PCR8 ₀
Initial value	0	0	0	0	0	0	0	0
Read/Write	W	W	W	W	W	W	W	W

PCR8 is an 8-bit register for controlling whether each of the port 8 pins $P8_7$ to $P8_0$ functions as an input or output pin. Setting a PCR8 bit to 1 makes the corresponding pin an output pin, while clearing the bit to 0 makes the pin an input pin. PCR8 and PDR8 settings are valid when the corresponding pins are designated for general-purpose input/output by bits SGS3 to SGS0 in LPCR.

Upon reset, PCR8 is initialized to H'00.

PCR8 is a write-only register, which is always read as all 1s.

8.8.3 Pin Functions

Table 8.21 shows the port 8 pin functions.

Table 8.21 Port 8 Pin Functions

1 abic 0.21	tort of mir unctions								
Pin	Pin Functions and	d Selection Method							
P8 ₇ /SEG ₃₂	The pin function de	epends on bit PCR8, ir	PCR8 and bits SGS	3 to SGS0 in LPCR.					
	SGS3 to SGS0	00	0*	001*, 01**, 1***					
	PCR8 ₇	0	1	*					
	Pin function	P8, input pin	P8, output pin	SEG ₃₂ output pin					
P8 ₆ /SEG ₃₁	The pin function de	epends on bit PCR8 ₆ ir	PCR8 and bits SGS	3 to SGS0 in LPCR.					
	SGS3 to SGS0	00	0*	001*, 01**, 1***					
	PCR8 ₆	0	1	*					
	Pin function	P8 ₆ input pin	P8 ₆ output pin	SEG ₃₁ output pin					
P8₅/SEG₃₀	The pin function de	pends on bit PCR8 ₅ ir	n PCR8 and bits SGS	3 to SGS0 in LPCR.					
	SGS3 to SGS0	00	0*	001*, 01**, 1***					
	PCR8₅	0	1	*					
	Pin function	P8₅ input pin	P8₅ output pin	SEG ₃₀ output pin					
P8 ₄ /SEG ₂₉	The pin function de	epends on bit PCR8₄ ir	PCR8 and bits SGS	3 to SGS0 in LPCR.					
	SGS3 to SGS0	00	0*	001*, 01**, 1***					
	PCR8 ₄	0	1	*					
	Pin function	P8₄ input pin	P8₄ output pin	SEG ₂₉ output pin					
P8 ₃ /SEG ₂₈ to	The pin function de	epends on bit PCR8, ir	PCR8 and bits SGS	3 to SGS0 in LPCR. (n = 3 to 0)					
ÿ <u>2</u> 5	SGS3 to SGS0	00	0*	001*, 01**, 1***					
	PCR8 _n	0	1	*					
	Pin function	P8 input pin	P8 output pin	SEG output pin					

*: Don't care



8.8.4 Pin States

Table 8.22 shows the port 8 pin states in each operating mode.

Table 8.22 Port 8 Pin States

Pins	Reset	Sleep	Subsleep	Standby	Watch	Subactive	Active
P8/SEG ₃₂ P8/SEG ₃₁ P8/SEG ₃₀ P8 ₄ /SEG ₂₉ P8 ₃ /SEG ₂₈ to P8 ₃ /SEG ₂₅	High- impedance	Retains previous state		High- impedance	Retains previous state	Functional	Functional

8.9 Port A

8.9.1 Overview

Port A is a 4-bit I/O port, configured as shown in figure 8.8.

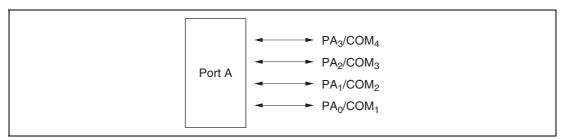


Figure 8.8 Port A Pin Configuration

8.9.2 Register Configuration and Description

Table 8.23 shows the port A register configuration.

Table 8.23 Port A Registers

Name	Abbr.	R/W	Initial Value	Address
Port data register A	PDRA	R/W	H'F0	H'FFDD
Port control register A	PCRA	W	H'F0	H'FFED

(1) Port Data Register A (PDRA)

Bit	7	6	5	4	3	2	1	0
	_	_	_	_	PA ₃	PA ₂	PA ₁	PA ₀
Initial value	1	1	1	1	0	0	0	0
Read/Write	_	_	_	_	R/W	R/W	R/W	R/W

PDRA is an 8-bit register that stores data for port A pins PA₃ to PA₀. If port A is read while PCRA bits are set to 1, the values stored in PDRA are read, regardless of the actual pin states. If port A is read while PCRA bits are cleared to 0, the pin states are read.

Upon reset, PDRA is initialized to H'F0.

(2) Port Control Register A (PCRA)

Bit	7	6	5	4	3	2	1	0
	_	_	_	_	PCRA ₃	PCRA ₂	PCRA ₁	PCRA ₀
Initial value	1	1	1	1	0	0	0	0
Read/Write	_	_	_	_	R/W	R/W	R/W	R/W

PCRA controls whether each of port A pins PA₃ to PA₀ functions as an input pin or output pin. Setting a PCRA bit to 1 makes the corresponding pin an output pin, while clearing the bit to 0 makes the pin an input pin. PCRA and PDRA settings are valid when the corresponding pins are designated for general-purpose input/output by LPCR.

Upon reset, PCRA is initialized to H'F0.

PCRA is a write-only register, which always reads all 1s.



8.9.3 Pin Functions

Table 8.24 shows the port A pin functions.

Table 8.24 Port A Pin Functions

Pin	Pin Functions and	Selection Method		
PA ₃ /COM ₄	The pin function dep	oends on bit PCRA ₃ i	n PCRA and bits SG	iS3 to SGS0.
	SGS3 to SGS0	0000	0000	Not 0000
	PCRA₃	0	1	*
	Pin function	PA ₃ input pin	PA ₃ output pin	COM ₄ output pin
PA ₂ /COM ₃	The pin function dep	pends on bit PCRA ₂ i	n PCRA and bits SG	S3 to SGS0.
	SGS3 to SGS0	0000	0000	Not 0000
	PCRA ₂	0	1	*
	Pin function	PA ₂ input pin	PA ₂ output pin	COM ₃ output pin
PA ₁ /COM ₂	The pin function dep	pends on bit PCRA, i	n PCRA and bits SG	iS3 to SGS0.
	SGS3 to SGS0	0000	0000	Not 0000
	PCRA,	0	1	*
	Pin function	PA, input pin	PA, output pin	COM ₂ output pin
PA ₀ /COM ₁	The pin function dep	pends on bit PCRA ₀ i	n PCRA and bits SG	iS3 to SGS0.
	SGS3 to SGS0	00	00	Not 0000
	PCRA₀	0	1	*
	Pin function	PA₀ input pin	PA₀ output pin	COM, output pin

*: Don't care

8.9.4 Pin States

Table 8.25 shows the port A pin states in each operating mode.

Table 8.25 Port A Pin States

Pins	Reset	Sleep	Subsleep	Standby	Watch	Subactive	Active
PA ₃ /COM ₄ PA ₂ /COM ₃ PA ₁ /COM ₂ PA ₃ /COM ₁	High- impedance	Retains previous state		High- impedance	Retains previous state	Functional	Functional

8.10 Port B

8.10.1 Overview

Port B is an 8-bit input-only port, configured as shown in figure 8.9.

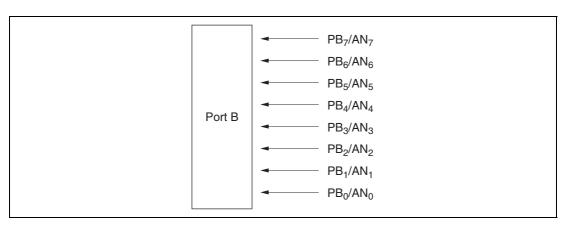


Figure 8.9 Port B Pin Configuration

8.10.2 Register Configuration and Description

Table 8.26 shows the port B register configuration.

Table 8.26 Port B Register

Name	Abbr.	R/W	Initial Value	Address
Port data register B	PDRB	R	Undefined	H'FFDE

Port Data Register B (PDRB)

Bit	7	6	5	4	3	2	1	0
	PB ₇	PB ₆	PB_5	PB ₄	PB ₃	PB ₂	PB ₁	PB ₀
Read/Write	R	R	R	R	R	R	R	R

Reading PDRB always gives the pin states. However, if a port B pin is selected as an analog input channel for the A/D converter by AMR bits CH3 to CH0, that pin reads 0 regardless of the input voltage.

8.11 Input/Output Data Inversion Function

8.11.1 Overview

With input pins \overline{WKP}_0 to \overline{WKP}_7 , RXD₃₁, and RXD₃₂, and output pins TXD₃₁ and TXD₃₂, the data can be handled in inverted form.

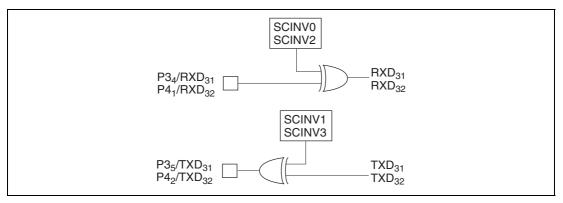


Figure 8.10 Input/Output Data Inversion Function

8.11.2 Register Configuration and Descriptions

Table 8.27 shows the registers used by the input/output data inversion function.

Table 8.27 Register Configuration

Name	Abbr.	R/W	Initial Value	Address
Serial port control register	SPCR	R/W	H'C0	H'FF91

Serial Port Control Register (SPCR)

Bit	7	6	5	4	3	2	1	0
	_	_	SPC32	SPC31	SCINV3	SCINV2	SCINV1	SCINV0
Initial value	1	1	0	0	0	0	0	0
Read/Write	_	_	R/W	R/W	R/W	R/W	R/W	R/W

SPCR is an 8-bit readable/writable register that performs RXD₃₁, RXD₃₂, TXD₃₁, and TXD₃₂ pin input/output data inversion switching. SPCR is initialized to H'C0 by a reset.

Bits 7 and 6: Reserved bits

Bits 7 and 6 are reserved; they are always read as 1 and cannot be modified.

Bit 5: P4₂/TXD₃₂ pin function switch (SPC32)

This bit selects whether pin P4₂/TXD₃₂ is used as P4₂ or as TXD₃₂.

Bit 5 SPC32	Description	
0	Functions as P4 ₂ I/O pin	(initial value)
1	Functions as TXD ₃₂ output pin*	

Note: * Set the TE bit in SCR3 after setting this bit to 1.

Bit 4: P3₅/TXD₃₁ pin function switch (SPC31)

This bit selects whether pin P35/TXD31 is used as P35 or as TXD31.

Bit 4 SPC31	Description	
0	Functions as P3 _s I/O pin	(initial value)
1	Functions as TXD ₃₁ output pin*	

Note: * Set the TE bit in SCR3 after setting this bit to 1.

Bit 3: TXD₃₂ pin output data inversion switch

Bit 3 specifies whether or not TXD₃₂ pin output data is to be inverted.

Bit 3 SCINV3	Description	
0	TXD ₃₂ output data is not inverted	(initial value)
1	TXD ₃₂ output data is inverted	

Bit 2: RXD₃₂ pin input data inversion switch

Bit 2 specifies whether or not RXD₃₂ pin input data is to be inverted.

Bit 2 SCINV2	Description	
0	RXD ₃₂ input data is not inverted	(initial value)
1	RXD ₃₂ input data is inverted	

Bit 1: TXD₃₁ pin output data inversion switch

Bit 1 specifies whether or not TXD₃₁ pin output data is to be inverted.

Bit 1 SCINV1	Description	
0	TXD ₃₁ output data is not inverted	(initial value)
1	TXD ₃₁ output data is inverted	

Bit 0: RXD₃₁ pin input data inversion switch

Bit 0 specifies whether or not RXD₃₁ pin input data is to be inverted.

Bit 0 SCINV0	Description	
0	RXD ₃₁ input data is not inverted	(initial value)
1	RXD ₃₁ input data is inverted	



8.11.3 Note on Modification of Serial Port Control Register

When a serial port control register is modified, the data being input or output up to that point is inverted immediately after the modification, and an invalid data change is input or output. When modifying a serial port control register, do so in a state in which data changes are invalidated.

8.12 Application Note

8.12.1 The Management of the Un-Use Terminal

If an I/O pin not used by the user system is floating, pull it up or down.

- If an unused pin is an input pin, handle it in one of the following ways:
 - Pull it up to V_{cc} with an on-chip pull-up MOS.
 - Pull it up to V_{cc} with an external resistor of approximately 100 k Ω .
 - Pull it down to V_{ss} with an external resistor of approximately 100 k Ω .
 - For a pin also used by the A/D converter, pull it up to AV_{cc}.
- If an unused pin is an output pin, handle it in one of the following ways:
 - Set the output of the unused pin to high and pull it up to V_{cc} with an external resistor of approximately 100 k Ω .
 - Set the output of the unused pin to low and pull it down to V_{ss} with an external resistor of approximately 100 k Ω .

Section 9 Timers

9.1 Overview

This LSI provides six timers: timers A, C, F, G, and a watchdog timer, and an asynchronous event counter. The functions of these timers are outlined in table 9.1.

Table 9.1 Timer Functions

Name	Functions	Internal Clock	Event Input Pin	Waveform Output Pin	Remarks
Timer A	• 8-bit interval timer	φ/8 to φ/8192	_	_	
	Interval function	(8 choices)			
	Time base	ϕ_w /128 (choice of 4 overflow periods)	_		
	Clock output	$\phi/4$ to $\phi/32$ ϕ_{w} , $\phi_{w}/4$ to $\phi_{w}/32$ (9 choices)	_	TMOW	-
Timer C	8-bit timer	φ/4 to φ/8192, φ _w /4	TMIC	_	Up-count/down-
	 Interval function 	(7 choices)			count controllable by
	 Event counting function 				software or hardware
	 Up-count/down- count selectable 				
Timer F	16-bit timer	φ/4 to φ/32, φ _w /4	TMIF	TMOFL	
	Event counting function Also usable as two independent 8-bit timers Output compare output function	(4 choices)		TMOFH	
Timer G	8-bit timer	φ/2 to φ/64, φ _w /4	TMIG	_	Counter
	 Input capture function 	(4 choices)			clearing option
	Interval function				 Built-in capture input signal noise canceler

Name	Functions	Internal Clock	Event Input Pin	Waveform Output Pin Remarks
Watchdog timer	Reset signal generated when8-bit counter overflows	φ/8192 φ _w /32	_	_
Asynchro-	16-bit counter	_	AEVL	_
nous event counter	 Also usable as two independent 8-bit counters 		AEVH	
	 Counts events asynchronous to φ and φw 			



9.2 Timer A

9.2.1 Overview

Timer A is an 8-bit timer with interval timing and real-time clock time-base functions. The clock time-base function is available when a 32.768-kHz crystal oscillator is connected. A clock signal divided from 32.768 kHz, from 38.4 kHz (if a 38.4 kHz crystal oscillator is connected), or from the system clock, can be output at the TMOW pin.

(1) Features

Features of timer A are given below.

- Choice of eight internal clock sources (\$\phi/8192\$, \$\phi/4096\$, \$\phi/2048\$, \$\phi/512\$, \$\phi/256\$, \$\phi/128\$, \$\phi/32\$, \$\phi/8\$).
- Choice of four overflow periods (1 s, 0.5 s, 0.25 s, 31.25 ms) when timer A is used as a clock time base (using a 32.768 kHz crystal oscillator).
- An interrupt is requested when the counter overflows.
- Any of nine clock signals can be output at the TMOW pin: 32.768 kHz divided by 32, 16, 8, or 4 (1 kHz, 2 kHz, 4 kHz, 8 kHz, 32,768 kHz) or 38.4 kHz divided by 32, 16, 8, or 4 (1.2 kHz, 2.4 kHz, 4.8 kHz, 9.6 kHz, 38.4 kHz), and the system clock divided by 32, 16, 8, or 4.
- Use of module standby mode enables this module to be placed in standby mode independently when not used.

(2) Block Diagram

Figure 9.1 shows a block diagram of timer A.

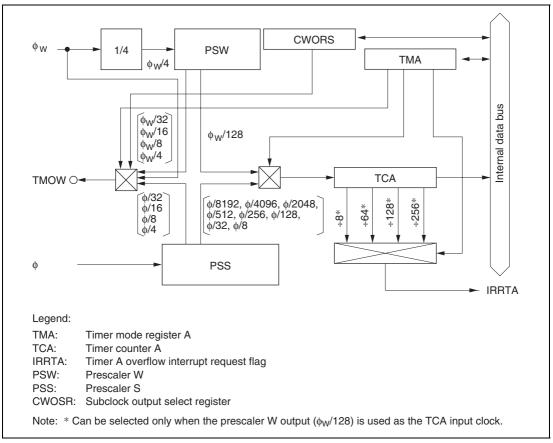


Figure 9.1 Block Diagram of Timer A

(3) Pin Configuration

Table 9.2 shows the timer A pin configuration.

Table 9.2 Pin Configuration

Name	Abbr.	I/O	Function
Clock output	TMOW	Output	Output of waveform generated by timer A output circuit



(4) Register Configuration

Table 9.3 shows the register configuration of timer A.

Table 9.3 Timer A Registers

Name	Abbr.	R/W	Initial Value	Address
Timer mode register A	TMA	R/W	H'10	H'FFB0
Timer counter A	TCA	R	H'00	H'FFB1
Clock stop register 1	CKSTPR1	R/W	H'FF	H'FFFA
Subclock output select register	CWOSR	R/W	H'FE	H'FF92

9.2.2 Register Descriptions

(1) Timer Mode Register A (TMA)

Bit	7	6	5	4	3	2	1	0
	TMA7	TMA6	TMA5	_	TMA3	TMA2	TMA1	TMA0
Initial value	0	0	0	1	0	0	0	0
Read/Write	R/W	R/W	R/W	_	R/W	R/W	R/W	R/W

TMA is an 8-bit read/write register for selecting the prescaler, input clock, and output clock.

Upon reset, TMA is initialized to H'10.

Bits 7 to 5: Clock output select (TMA7 to TMA5)

Bits 7 to 5 choose which of eight clock signals is output at the TMOW pin. The system clock divided by 32, 16, 8, or 4 can be output in active mode and sleep mode. A 32.768 kHz or 38.4 kHz signal divided by 32, 16, 8, or 4 can be output in active mode, sleep mode, and subactive mode. ϕ_w is output in all modes except the reset state.

CWOSR	TMA				
cwos	Bit 7 TMA7	Bit 6 TMA6	Bit 5 TMA5	Clock Output	
0	0	0	0	φ/32	(initial value)
			1	φ/16	
		1	0	φ/8	
			1	φ/4	
	1	0	0	φ _w /32	
			1	φ _w /16	
		1	0	φ _w /8	
			1	φ_/4	
1	*	*	*	ϕ_{w}	

*: Don't care

Bit 4: Reserved bit

Bit 4 is reserved; it is always read as 1, and cannot be modified.

Bits 3 to 0: Internal clock select (TMA3 to TMA0)

Bits 3 to 0 select the clock input to TCA. The selection is made as follows.

				Description	
Bit 3 TMA3	Bit 2 TMA2	Bit 1 TMA1	Bit 0 TMA0	Prescaler and Divider Ratio or Overflow Period	Function
0	0	0	0	PSS, $\phi/8192$ (initial value)	Interval timer
			1	PSS,	_
		1	0	PSS,	_
			1	PSS, ø/512	_
	1	0	0	PSS, ø/256	_
			1	PSS, ø/128	_
		1	0	PSS, ¢/32	_
			1	PSS, ø/8	_
1	0	0	0	PSW, 1 s	Clock time
			1	PSW, 0.5 s	base (when using 32.768
		1	0	PSW, 0.25 s	kHz)
			1	PSW, 0.03125 s	_
	1	0	0	PSW and TCA are reset	_
			1		
		1	0	_	
			1	_	

(2) Timer Counter A (TCA)

Bit	7	6	5	4	3	2	1	0
	TCA7	TCA6	TCA5	TCA4	TCA3	TCA2	TCA1	TCA0
Initial value	0	0	0	0	0	0	0	0
Read/Write	R	R	R	R	R	R	R	R

TCA is an 8-bit read-only up-counter, which is incremented by internal clock input. The clock source for input to this counter is selected by bits TMA3 to TMA0 in timer mode register A (TMA). TCA values can be read by the CPU in active mode, but cannot be read in subactive mode. When TCA overflows, the IRRTA bit in interrupt request register 1 (IRR1) is set to 1.

TCA is cleared by setting bits TMA3 and TMA2 of TMA to 11.

Upon reset, TCA is initialized to H'00.

(3) Clock Stop Register 1 (CKSTPR1)

Bit	7	6	5	4	3	2	1	0
		S31CKSTP	S32CKSTP	ADCKSTP	TGCKSTP	TFCKSTP	TCCKSTP	TACKSTP
Initial value	1	1	1	1	1	1	1	1
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W

CKSTPR1 is an 8-bit read/write register that performs module standby mode control for peripheral modules. Only the bit relating to timer A is described here. For details of the other bits, see the sections on the relevant modules.

Bit 0: Timer A module standby mode control (TACKSTP)

Bit 0 controls setting and clearing of module standby mode for timer A.

TACKSTP	Description	
0	Timer A is set to module standby mode	_
1	Timer A module standby mode is cleared	(initial value)

(4) Subclock Output Select Register (CWOSR)

Bit	7	6	5	4	3	2	1	0
	_		_	_	_	_	_	cwos
Initial value	1	1	1	1	1	1	1	0
Read/Write	R	R	R	R	R	R	R	R/W

CWOSR is an 8-bit read/write register that selects the clock to be output from the TMOW pin.

CWOSR is initialized to H'FE by a reset.

Bits 7 to 1: Reserved bits

Bits 7 to 1 are reserved; they are always read as 1 and cannot be modified.



Bit 0: TMOW pin clock select (CWOS)

Bit 0 selects the clock to be output from the TMOW pin.

Bit 0		
cwos	Description	
0	Clock output from timer A is output (see TMA)	(initial value)
1	ϕ_{w} is output	

9.2.3 Timer Operation

(1) Interval Timer Operation

When bit TMA3 in timer mode register A (TMA) is cleared to 0, timer A functions as an 8-bit interval timer.

Upon reset, TCA is cleared to H'00 and bit TMA3 is cleared to 0, so up-counting and interval timing resume immediately. The clock input to timer A is selected by bits TMA2 to TMA0 in TMA; any of eight internal clock signals output by prescaler S can be selected.

After the count value in TCA reaches H'FF, the next clock signal input causes timer A to overflow, setting bit IRRTA to 1 in interrupt request register 1 (IRR1). If IENTA = 1 in interrupt enable register 1 (IENR1), a CPU interrupt is requested.*

At overflow, TCA returns to H'00 and starts counting up again. In this mode timer A functions as an interval timer that generates an overflow output at intervals of 256 input clock pulses.

Note: * For details on interrupts, see section 3.3, Interrupts.

(2) Real-Time Clock Time Base Operation

When bit TMA3 in TMA is set to 1, timer A functions as a real-time clock time base by counting clock signals output by prescaler W. The overflow period of timer A is set by bits TMA1 and TMA0 in TMA. A choice of four periods is available. In time base operation (TMA3 = 1), setting bit TMA2 to 1 clears both TCA and prescaler W to their initial values of H'00.

(3) Clock Output

Setting bit TMOW in port mode register 1 (PMR1) to 1 causes a clock signal to be output at pin TMOW. Nine different clock output signals can be selected by means of bits TMA7 to TMA5 in TMA and bit CWOS in CWOSR. The system clock divided by 32, 16, 8, or 4 can be output in active mode and sleep mode. A 32.768 kHz or 38.4 kHz signal divided by 32, 16, 8, or 4 can be output in active mode, sleep mode, watch mode, subactive mode, and subsleep mode. The 32.768 kHz or 38.4 kHz clock is output in all modes except the reset state.

9.2.4 Timer A Operation States

Table 9.4 summarizes the timer A operation states.

Table 9.4 Timer A Operation States

Oper	ation Mode	Reset	Active	Sleep	Watch	Sub- active	Sub- sleep	Standby	Module Standby
TCA	Interval	Reset	Functions	Functions	Halted	Halted	Halted	Halted	Halted
	Clock time base	Reset	Functions	Functions	Functions	Functions	Functions	Halted	Halted
TMA	CWOSR	Reset	Functions	Retained	Retained	Functions	Retained	Retained	Retained

Note: When the real-time clock time base function is selected as the internal clock of TCA in active mode or sleep mode, the internal clock is not synchronous with the system clock, so it is synchronized by a synchronizing circuit. This may result in a maximum error of 1/φ (s) in the count cycle.

9.2.5 Application Note

When bit 0 (TACKSTP) of the clock stop register 1 (CKSTPR1) is cleared to 0, bit 3 (TMA3) of the timer mode register A (TMA) cannot be rewritten.

Set bit 0 (TACKSTP) of the clock stop register 1 (CKSTPR1) to 1 before rewriting bit 3 (TMA3) of the timer mode register A (TMA).



9.3 Timer C

9.3.1 Overview

Timer C is an 8-bit timer that increments each time a clock pulse is input. This timer has two operation modes, interval and auto reload.

(1) Features

Features of timer C are given below.

- Choice of seven internal clock sources (φ/8192, φ/2048, φ/512, φ/64, φ/16, φ/4, φw/4) or an external clock (can be used to count external events).
- An interrupt is requested when the counter overflows.
- Up/down-counter switching is possible by hardware or software.
- Subactive mode and subsleep mode operation is possible when ϕ w/4 is selected as the internal clock, or when an external clock is selected.
- Use of module standby mode enables this module to be placed in standby mode independently when not used.

(2) Block Diagram

Figure 9.2 shows a block diagram of timer C.

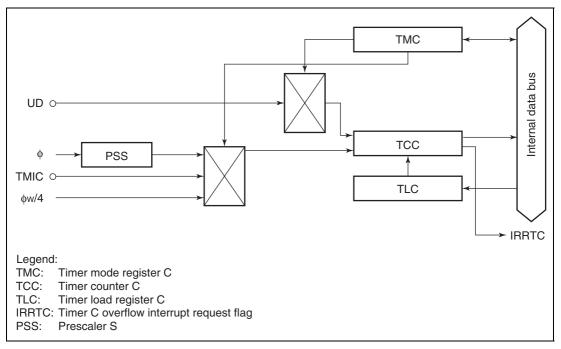


Figure 9.2 Block Diagram of Timer C

(3) Pin Configuration

Table 9.5 shows the timer C pin configuration.

Table 9.5 Pin Configuration

Name	Abbr.	I/O	Function
Timer C event input	TMIC	Input	Input pin for event input to TCC
Timer C up/down-count selection	UD	Input	Timer C up/down select

(4) Register Configuration

Table 9.6 shows the register configuration of timer C.

Table 9.6 Timer C Registers

Name	Abbr.	R/W	Initial Value	Address
Timer mode register C	TMC	R/W	H'18	H'FFB4
Timer counter C	TCC	R	H'00	H'FFB5
Timer load register C	TLC	W	H'00	H'FFB5
Clock stop register 1	CKSTPR1	R/W	H'FF	H'FFFA

9.3.2 Register Descriptions

(1) Timer Mode Register C (TMC)

Bit	7	6	5	4	3	2	1	0
	TMC7	TMC6	TMC5	_	_	TMC2	TMC1	TMC0
Initial value	0	0	0	1	1	0	0	0
Read/Write	R/W	R/W	R/W	_	_	R/W	R/W	R/W

TMC is an 8-bit read/write register for selecting the auto-reload function and input clock, and performing up/down-counter control.

Upon reset, TMC is initialized to H'18.

Bit 7: Auto-reload function select (TMC7)

Bit 7 selects whether timer C is used as an interval timer or auto-reload timer.

Bit 7 TMC7	Description	
0	Interval timer function selected	(initial value)
1	Auto-reload function selected	

Bits 6 and 5: Counter up/down control (TMC6, TMC5)

Selects whether TCC up/down control is performed by hardware using UD pin input, or whether TCC functions as an up-counter or a down-counter.

Bit 6 TMC6	Bit 5 TMC5	Description	
0	0	TCC is an up-counter	(initial value)
0	1	TCC is a down-counter	
1	*	Hardware control by UD pin input UD pin input high: Down-counter UD pin input low: Up-counter	

*: Don't care

Bits 4 and 3: Reserved bits

Bits 4 and 3 are reserved; they are always read as 1 and cannot be modified.

Bits 2 to 0: Clock select (TMC2 to TMC0)

Bits 2 to 0 select the clock input to TCC. For external event counting, either the rising or falling edge can be selected.

Bit 2 TMC2	Bit 1 TMC1	Bit 0 TMC0	Description
0	0	0	Internal clock: $\phi/8192$ (initial value)
0	0	1	Internal clock: $\phi/2048$
0	1	0	Internal clock:
0	1	1	Internal clock: $\phi/64$
1	0	0	Internal clock: $\phi/16$
1	0	1	Internal clock: $\phi/4$
1	1	0	Internal clock: ϕ w/4
1	1	1	External event (TMIC): rising or falling edge*

Note: * The edge of the external event signal is selected by bit IEG1 in the IRQ edge select register (IEGR). See (1) IRQ Edge Select Register (IEGR) in 3.3.2, Interrupt Control Registers, for details. IRQ1 must be set to 1 in port mode register 1 (PMR1) before setting 111 in bits TMC2 to TMC0.



(2) Timer Counter C (TCC)

Bit	7	6	5	4	3	2	1	0
	TCC7	TCC6	TCC5	TCC4	TCC3	TCC2	TCC1	TCC0
Initial value	0	0	0	0	0	0	0	0
Read/Write	R	R	R	R	R	R	R	R

TCC is an 8-bit read-only up-counter, which is incremented by internal clock or external event input. The clock source for input to this counter is selected by bits TMC2 to TMC0 in timer mode register C (TMC). TCC values can be read by the CPU at any time.

When TCC overflows from H'FF to H'00 or to the value set in TLC, or underflows from H'00 to H'FF or to the value set in TLC, the IRRTC bit in IRR2 is set to 1.

TCC is allocated to the same address as TLC.

Upon reset, TCC is initialized to H'00.

(3) Timer Load Register C (TLC)

Bit	7	6	5	4	3	2	1	0
	TLC7	TLC6	TLC5	TLC4	TLC3	TLC2	TLC1	TLC0
Initial value	0	0	0	0	0	0	0	0
Read/Write	W	W	W	W	W	W	W	W

TLC is an 8-bit write-only register for setting the reload value of timer counter C (TCC).

When a reload value is set in TLC, the same value is loaded into timer counter C as well, and TCC starts counting up from that value. When TCC overflows or underflows during operation in autoreload mode, the TLC value is loaded into TCC. Accordingly, overflow/underflow periods can be set within the range of 1 to 256 input clocks.

The same address is allocated to TLC as to TCC.

Upon reset, TLC is initialized to H'00.

(4) Clock Stop Register 1 (CKSTPR1)

Bit	7	6	5	4	3	2	1	0
	_	S31CKSTP	S32CKSTP	ADCKSTP	TGCKSTP	TFCKSTP	TCCKSTP	TACKSTP
Initial value	1	1	1	1	1	1	1	1
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W

CKSTPR1 is an 8-bit read/write register that performs module standby mode control for peripheral modules. Only the bit relating to timer C is described here. For details of the other bits, see the sections on the relevant modules.

Bit 1: Timer C module standby mode control (TCCKSTP)

Bit 1 controls setting and clearing of module standby mode for timer C.

TCCKSTP	Description	
0	Timer C is set to module standby mode	
1	Timer C module standby mode is cleared	(initial value)

9.3.3 Timer Operation

(1) Interval Timer Operation

When bit TMC7 in timer mode register C (TMC) is cleared to 0, timer C functions as an 8-bit interval timer.

Upon reset, TCC is initialized to H'00 and TMC to H'18, so TCC continues up-counting as an interval up-counter without halting immediately after a reset. The timer C operating clock is selected from seven internal clock signals output by prescalers S and W, or an external clock input at pin TMIC. The selection is made by bits TMC2 to TMC0 in TMC.

TCC up/down-count control can be performed either by software or hardware. The selection is made by bits TMC6 and TMC5 in TMC.

After the count value in TCC reaches H'FF (H'00), the next clock input causes timer C to overflow (underflow), setting bit IRRTC to 1 in IRR2. If IENTC = 1 in interrupt enable register 2 (IENR2), a CPU interrupt is requested.

At overflow (underflow), TCC returns to H'00 (H'FF) and starts counting up (down) again.



During interval timer operation (TMC7 = 0), when a value is set in timer load register C (TLC), the same value is set in TCC.

Note: For details on interrupts, see section 3.3, Interrupts.

(2) Auto-Reload Timer Operation

Setting bit TMC7 in TMC to 1 causes timer C to function as an 8-bit auto-reload timer. When a reload value is set in TLC, the same value is loaded into TCC, becoming the value from which TCC starts its count.

After the count value in TCC reaches H'FF (H'00), the next clock signal input causes timer C to overflow/underflow. The TLC value is then loaded into TCC, and the count continues from that value. The overflow/underflow period can be set within a range from 1 to 256 input clocks, depending on the TLC value.

The clock sources, up/down control, and interrupts in auto-reload mode are the same as in interval mode.

In auto-reload mode (TMC7 = 1), when a new value is set in TLC, the TLC value is also set in TCC.

(3) Event Counter Operation

Timer C can operate as an event counter, counting rising or falling edges of an external event signal input at pin TMIC. External event counting is selected by setting bits TMC2 to TMC0 in timer mode register C to all 1s (111).

When timer C is used to count external event input, bit IRQ1 in PMR1 should be set to 1 and bit IEN1 in IENR1 cleared to 0 to disable interrupt IRQ₁ requests.

(4) TCC Up/Down Control by Hardware

With timer C, TCC up/down control can be performed by UD pin input. When bit TMC6 is set to 1 in TMC, TCC functions as an up-counter when UD pin input is high, and as a down-counter when low.

When using UD pin input, set bit UD to 1 in PMR3.



9.3.4 Timer C Operation States

Table 9.7 summarizes the timer C operation states.

Table 9.7 Timer C Operation States

						Sub-	Sub-sleep		Module
Operation Mode		Reset	Active	Sleep	Watch	active		Standby	Standby
TCC	Interval	Reset	Functions	Functions	Halted	Functions/ Halted*	Functions/ Halted*	Halted	Halted
	Auto reload	Reset	Functions	Functions	Halted	Functions/ Halted*	Functions/ Halted*	Halted	Halted
TMC		Reset	Functions	Retained	Retained	Functions	Retained	Retained	Retained

Note: * When ϕ w/4 is selected as the TCC internal clock in active mode or sleep mode, since the system clock and internal clock are mutually asynchronous, synchronization is maintained by a synchronization circuit. This results in a maximum count cycle error of $1/\phi$ (s). When the counter is operated in subactive mode or subsleep mode, either select ϕ w/4 as the internal clock or select an external clock. The counter will not operate on any other internal clock. If ϕ w/4 is selected as the internal clock for the counter when ϕ w/8 has been selected as subclock ϕ _{SuB}, the lower 2 bits of the counter operate on the same cycle, and the operation of the least significant bit is unrelated to the operation of the counter.

9.3.5 Usage Note

Note the following regarding the operation of timer C.

(1) Counting errors caused by external event input

Timer counter errors may occur under the following conditions.

Conditions

• An external event (TMIC) is used in subsleep mode.

Symptom

The counter increments or decrements twice for a single external event input.

Approximate rate of occurrence

The approximate rate of occurrence in cases where the external event input is not synchronized with internal operation is defined by the following equation.

Approximate rate of occurrence P = 30 ns / tsubcyc

For example, if tsubcyc = $61.06 \,\mu s$ (subclock $\phi w/2$), $P = 0.0005 \,(0.05\%)$. If 2,000 external event inputs occur, there is a likelihood that one of them will cause the counter to increment or decrement twice (+2 or -2).

The symptom described is caused by the internal circuit configuration of the device and therefore difficult to avoid. Therefore, it is not advisable to use the clock counter for applications requiring a high degree of accuracy.

9.4 Timer F

9.4.1 Overview

Timer F is a 16-bit timer with a built-in output compare function. As well as counting external events, timer F also provides for counter resetting, interrupt request generation, toggle output, etc., using compare match signals. Timer F can also be used as two independent 8-bit timers (timer FH and timer FL).

(1) Features

Features of timer F are given below.

- Choice of four internal clock sources ($\phi/32$, $\phi/16$, $\phi/4$, $\phi w/4$) or an external clock (can be used as an external event counter)
- TMOFH pin (TMOFL pin) toggle output provided using a single compare match signal (toggle output initial value can be set)
- Counter resetting by a compare match signal
- Two interrupt sources: one compare match, one overflow
- Can operate as two independent 8-bit timers (timer FH and timer FL) (in 8-bit mode).

		Timer FL			
	Timer FH 8-Bit Timer*	8-Bit Timer/Event Counter			
Internal clock	Choice of 4 (\$\phi/32, \$\phi/16, \$\phi/4, \$\phi w/4)				
Event input	_	TMIF pin			
Toggle output	One compare match signal, output to TMOFH pin(initial value settable)	One compare match signal, output to TMOFL pin (initial value settable)			
Counter reset	Counter can be reset by compare match signal				
Interrupt sources	One compare match One overflow				

Note: * When timer F operates as a 16-bit timer, it operates on the timer FL overflow signal.

- Operation in watch mode, subactive mode, and subsleep mode
 When \(\phi \widetilde{4} \) is selected as the internal clock, timer F can operate in watch mode, subactive mode, and subsleep mode.
- Use of module standby mode enables this module to be placed in standby mode independently when not used.



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Block Diagram (2)

Figure 9.3 shows a block diagram of timer F.

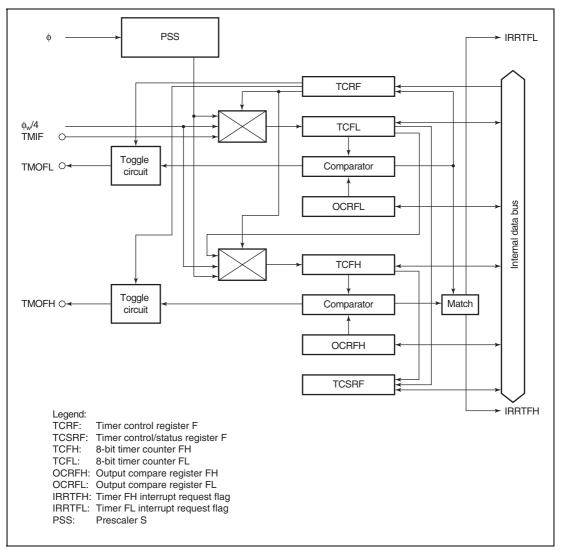


Figure 9.3 Block Diagram of Timer F

(3) Pin Configuration

Table 9.8 shows the timer F pin configuration.

Table 9.8 Pin Configuration

Name	Abbr.	I/O	Function
Timer F event input	TMIF	Input	Event input pin for input to TCFL
Timer FH output	TMOFH	Output	Timer FH toggle output pin
Timer FL output	TMOFL	Output	Timer FL toggle output pin

(4) Register Configuration

Table 9.9 shows the register configuration of timer F.

Table 9.9 Timer F Registers

Name	Abbr.	R/W	Initial Value	Address
Timer control register F	TCRF	W	H'00	H'FFB6
Timer control/status register F	TCSRF	R/W	H'00	H'FFB7
8-bit timer counter FH	TCFH	R/W	H'00	H'FFB8
8-bit timer counter FL	TCFL	R/W	H'00	H'FFB9
Output compare register FH	OCRFH	R/W	H'FF	H'FFBA
Output compare register FL	OCRFL	R/W	H'FF	H'FFBB
Clock stop register 1	CKSTPR1	R/W	H'FF	H'FFFA

9.4.2 **Register Descriptions**

(1) 16-bit Timer Counter (TCF) 8-bit Timer Counter (TCFH)

8-bit Timer Counter (TCFL)

								TO	CF							
Bit	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
Initial value	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
Read/Write	R/W															
				TC	FH							TC	FL			

TCF is a 16-bit read/write up-counter configured by cascaded connection of 8-bit timer counters TCFH and TCFL. In addition to the use of TCF as a 16-bit counter with TCFH as the upper 8 bits and TCFL as the lower 8 bits, TCFH and TCFL can also be used as independent 8-bit counters.

TCFH and TCFL can be read and written by the CPU, but when they are used in 16-bit mode, data transfer to and from the CPU is performed via a temporary register (TEMP). For details of TEMP, see section 9.4.3, CPU Interface.

TCFH and TCFL are each initialized to H'00 upon reset.

a. 16-bit mode (TCF)

When CKSH2 is cleared to 0 in TCRF, TCF operates as a 16-bit counter. The TCF input clock is selected by bits CKSL2 to CKSL0 in TCRF.

TCF can be cleared in the event of a compare match by means of CCLRH in TCSRF.

When TCF overflows from H'FFFF to H'0000, OVFH is set to 1 in TCSRF. If OVIEH in TCSRF is 1 at this time, IRRTFH is set to 1 in IRR2, and if IENTFH in IENR2 is 1, an interrupt request is sent to the CPU.

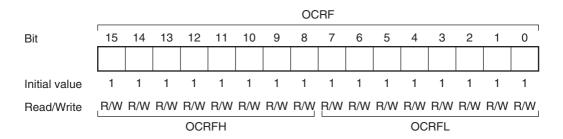
b. 8-bit mode (TCFL/TCFH)

When CKSH2 is set to 1 in TCRF, TCFH and TCFL operate as two independent 8-bit counters. The TCFH (TCFL) input clock is selected by bits CKSH2 to CKSH0 (CKSL2 to CKSL0) in TCRF.

TCFH (TCFL) can be cleared in the event of a compare match by means of CCLRH (CCLRL) in TCSRF.

When TCFH (TCFL) overflows from H'FF to H'00, OVFH (OVFL) is set to 1 in TCSRF. If OVIEH (OVIEL) in TCSRF is 1 at this time, IRRTFH (IRRTFL) is set to 1 in IRR2, and if IENTFH (IENTFL) in IENR2 is 1, an interrupt request is sent to the CPU.

(2) 16-bit Output Compare Register (OCRF) 8-bit Output Compare Register (OCRFH) 8-bit Output Compare Register (OCRFL)



OCRF is a 16-bit read/write register composed of the two registers OCRFH and OCRFL. In addition to the use of OCRF as a 16-bit register with OCRFH as the upper 8 bits and OCRFL as the lower 8 bits, OCRFH and OCRFL can also be used as independent 8-bit registers.

OCRFH and OCRFL can be read and written by the CPU, but when they are used in 16-bit mode, data transfer to and from the CPU is performed via a temporary register (TEMP). For details of TEMP, see section 9.4.3, CPU Interface.

OCRFH and OCRFL are each initialized to H'FF upon reset.

a. 16-bit mode (OCRF)

When CKSH2 is cleared to 0 in TCRF, OCRF operates as a 16-bit register. OCRF contents are constantly compared with TCF, and when both values match, CMFH is set to 1 in TCSRF. At the same time, IRRTFH is set to 1 in IRR2. If IENTFH in IENR2 is 1 at this time, an interrupt request is sent to the CPU.

Toggle output can be provided from the TMOFH pin by means of compare matches, and the output level can be set (high or low) by means of TOLH in TCRF.

b. 8-bit mode (OCRFH/OCRFL)

When CKSH2 is set to 1 in TCRF, OCRFH and OCRFL operate as two independent 8-bit registers. OCRFH contents are compared with TCFH, and OCRFL contents are with TCFL. When the OCRFH (OCRFL) and TCFH (TCFL) values match, CMFH (CMFL) is set to 1 in TCSRF. At the same time, IRRTFH (IRRTFL) is set to 1 in IRR2. If IENTFH (IENTFL) in IENR2 is 1 at this time, an interrupt request is sent to the CPU.



Toggle output can be provided from the TMOFH pin (TMOFL pin) by means of compare matches, and the output level can be set (high or low) by means of TOLH (TOLL) in TCRF.

(3) Timer Control Register F (TCRF)

Bit	7	6	5	4	3	2	1	0
	TOLH	CKSH2	CKSH1	CKSH0	TOLL	CKSL2	CKSL1	CKSL0
Initial value	0	0	0	0	0	0	0	0
Read/Write	W	W	W	W	W	W	W	W

TCRF is an 8-bit write-only register that switches between 16-bit mode and 8-bit mode, selects the input clock from among four internal clock sources or external event input, and sets the output level of the TMOFH and TMOFL pins.

TCRF is initialized to H'00 upon reset.

Bit 7: Toggle output level H (TOLH)

Bit 7 sets the TMOFH pin output level. The output level is effective immediately after this bit is written.

Bit 7 TOLH	Description	
0	Low level	(initial value)
1	High level	

Bits 6 to 4: Clock select H (CKSH2 to CKSH0)

Bits 6 to 4 select the clock input to TCFH from among four internal clock sources or TCFL overflow.

Bit 6 CKSH2	Bit 5 CKSH1	Bit 4 CKSH0	Description
0	0	0	16-bit mode, counting on TCFL overflow signal (initial value)
0	0	1	_
0	1	0	_
0	1	1	Use prohibited
1	0	0	Internal clock: Counting on $\phi/32$
1	0	1	Internal clock: Counting on $\phi/16$
1	1	0	Internal clock: Counting on $\phi/4$
1	1	1	Internal clock: Counting on ϕ w/4

Bit 3: Toggle output level L (TOLL)

Bit 3 sets the TMOFL pin output level. The output level is effective immediately after this bit is written.

Bit 3 TOLL	Description	
0	Low level	(initial value)
1	High level	

Bits 2 to 0: Clock select L (CKSL2 to CKSL0)

Bits 2 to 0 select the clock input to TCFL from among four internal clock sources or external event input.

Bit 2 CKSL2	Bit 1 CKSL1	Bit 0 CKSL0	Description
0	0	0	Counting on external event (TMIF) rising/falling edge*
0	0	1	— (initial value)
0	1	0	_
0	1	1	Use prohibited
1	0	0	Internal clock: Counting on $\phi/32$
1	0	1	Internal clock: Counting on $\phi/16$
1	1	0	Internal clock: Counting on φ/4
1	1	1	Internal clock: Counting on φw/4

Note: * External event edge selection is set by IEG3 in the IRQ edge select register (IEGR). For details, see (1) IRQ Edge Select Register (IEGR) in section 3.3.2, Interrupt Control Registers.

Note that the timer F counter may increment if the setting of IRQ3 in port mode register 1 (PMR1) is changed from 0 to 1 while the TMIF pin is low in order to change the TMIF pin function.

(4) Timer Control/Status Register F (TCSRF)

Bit	7	6	5	4	3	2	1	0
	OVFH	CMFH	OVIEH	CCLRH	OVFL	CMFL	OVIEL	CCLRL
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W*	R/W*	R/W	R/W	R/W*	R/W*	R/W	R/W

Note: * Bits 7, 6, 3, and 2 can only be written with 0, for flag clearing.

TCSRF is an 8-bit read/write register that performs counter clear selection, overflow flag setting, and compare match flag setting, and controls enabling of overflow interrupt requests.

TCSRF is initialized to H'00 upon reset.

Bit 7: Timer overflow flag H (OVFH)

Bit 7 is a status flag indicating that TCFH has overflowed from H'FF to H'00. This flag is set by hardware and cleared by software. It cannot be set by software.

Bit 7 OVFH	Description	
0	Clearing condition: After reading OVFH = 1, cleared by writing 0 to OVFH	(initial value)
1	Setting condition: Set when TCFH overflows from H'FF to H'00	

Bit 6: Compare match flag H (CMFH)

Bit 6 is a status flag indicating that TCFH has matched OCRFH. This flag is set by hardware and cleared by software. It cannot be set by software.

Bit 6 CMFH	Description	
0	Clearing condition: After reading CMFH = 1, cleared by writing 0 to CMFH	(initial value)
1	Setting condition: Set when the TCFH value matches the OCRFH value	

Bit 5: Timer overflow interrupt enable H (OVIEH)

Bit 5 selects enabling or disabling of interrupt generation when TCFH overflows.

Bit 5 OVIEH	Description	
OVILII	Description	
0	TCFH overflow interrupt request is disabled	(initial value)
1	TCFH overflow interrupt request is enabled	



Bit 4: Counter clear H (CCLRH)

In 16-bit mode, bit 4 selects whether TCF is cleared when TCF and OCRF match.

In 8-bit mode, bit 4 selects whether TCFH is cleared when TCFH and OCRFH match.

Bit 4 CCLRH	Description	
0	16-bit mode: TCF clearing by compare match is disabled	_
	8-bit mode: TCFH clearing by compare match is disabled	(initial value)
1	16-bit mode: TCF clearing by compare match is enabled	
	8-bit mode: TCFH clearing by compare match is enabled	

Bit 3: Timer overflow flag L (OVFL)

Bit 3 is a status flag indicating that TCFL has overflowed from H'FF to H'00. This flag is set by hardware and cleared by software. It cannot be set by software.

Bit 3 OVFL	Description	
0	Clearing condition:	(initial value)
	After reading OVFL = 1, cleared by writing 0 to OVFL	
1	Setting condition:	
	Set when TCFL overflows from H'FF to H'00	

Bit 2: Compare match flag L (CMFL)

Bit 2 is a status flag indicating that TCFL has matched OCRFL. This flag is set by hardware and cleared by software. It cannot be set by software.

Bit 2 CMFL	Description	
0	Clearing condition:	(initial value)
	After reading CMFL = 1, cleared by writing 0 to CMFL	
1	Setting condition:	
	Set when the TCFL value matches the OCRFL value	

Bit 1: Timer overflow interrupt enable L (OVIEL)

Bit 1 selects enabling or disabling of interrupt generation when TCFL overflows.

Bit 1 OVIEL	Description	
0	TCFL overflow interrupt request is disabled	(initial value)
1	TCFL overflow interrupt request is enabled	

Bit 0: Counter clear L (CCLRL)

Bit 0 selects whether TCFL is cleared when TCFL and OCRFL match.

Bit 0 CCLRL	Description	
0	TCFL clearing by compare match is disabled	(initial value)
1	TCFL clearing by compare match is enabled	

(5) Clock Stop Register 1 (CKSTPR1)

Bit	7	6	5	4	3	2	1	0
	_	S31CKSTP	S32CKSTP	ADCKSTP	TGCKSTP	TFCKSTP	TCCKSTP	TACKSTP
Initial value	1	1	1	1	1	1	1	1
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W

CKSTPR1 is an 8-bit read/write register that performs module standby mode control for peripheral modules. Only the bit relating to timer F is described here. For details of the other bits, see the sections on the relevant modules.

Bit 2: Timer F module standby mode control (TFCKSTP)

Bit 2 controls setting and clearing of module standby mode for timer F.

TFCKSTP	Description	
0	Timer F is set to module standby mode	
1	Timer F module standby mode is cleared	(initial value)



9.4.3 CPU Interface

TCF and OCRF are 16-bit read/write registers, but the CPU is connected to the on-chip peripheral modules by an 8-bit data bus. When the CPU accesses these registers, it therefore uses an 8-bit temporary register (TEMP).

In 16-bit mode, TCF read/write access and OCRF write access must be performed 16 bits at a time (using two consecutive byte-size MOV instructions), and the upper byte must be accessed before the lower byte. Data will not be transferred correctly if only the upper byte or only the lower byte is accessed.

In 8-bit mode, there are no restrictions on the order of access.

(1) Write Access

Write access to the upper byte results in transfer of the upper-byte write data to TEMP. Next, write access to the lower byte results in transfer of the data in TEMP to the upper register byte, and direct transfer of the lower-byte write data to the lower register byte.

Figure 9.4 shows an example in which H'AA55 is written to TCF.

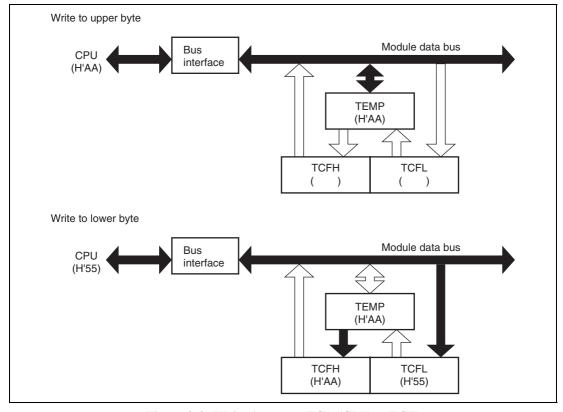


Figure 9.4 Write Access to TCR (CPU \rightarrow TCF)

(2) Read Access

In access to TCF, when the upper byte is read the upper-byte data is transferred directly to the CPU and the lower-byte data is transferred to TEMP. Next, when the lower byte is read, the lower-byte data in TEMP is transferred to the CPU.

In access to OCRF, when the upper byte is read the upper-byte data is transferred directly to the CPU. When the lower byte is read, the lower-byte data is transferred directly to the CPU.

Figure 9.5 shows an example in which TCF is read when it contains H'AAFF.

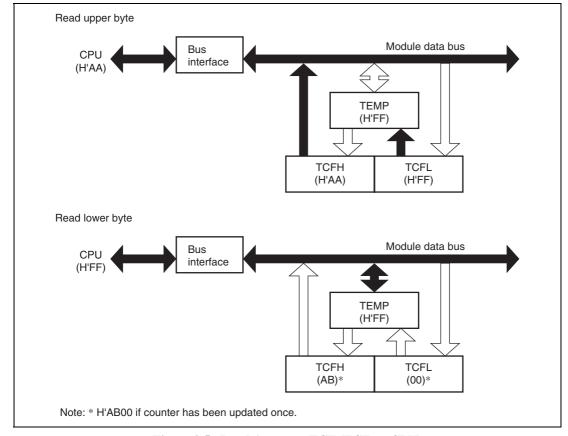


Figure 9.5 Read Access to TCF (TCF \rightarrow CPU)

9.4.4 Operation

Timer F is a 16-bit counter that increments on each input clock pulse. The timer F value is constantly compared with the value set in output compare register F, and the counter can be cleared, an interrupt requested, or port output toggled, when the two values match. Timer F can also function as two independent 8-bit timers.

(1) Timer F Operation

Timer F has two operating modes, 16-bit timer mode and 8-bit timer mode. The operation in each of these modes is described below.

a. Operation in 16-bit timer mode

When CKSH2 is cleared to 0 in timer control register F (TCRF), timer F operates as a 16-bit timer.

Following a reset, timer counter F (TCF) is initialized to H'0000, output compare register F (OCRF) to H'FFFF, and timer control register F (TCRF) and timer control/status register F (TCSRF) to H'00. The counter starts incrementing on external event (TMIF) input. The external event edge selection is set by IEG3 in the IRQ edge select register (IEGR).

The timer F operating clock can be selected from four internal clocks or an external clock by means of bits CKSL2 to CKSL0 in TCRF.

OCRF contents are constantly compared with TCF, and when both values match, CMFH is set to 1 in TCSRF. If IENTFH in IENR2 is 1 at this time, an interrupt request is sent to the CPU, and at the same time, TMOFH pin output is toggled. If CCLRH in TCSRF is 1, TCF is cleared. TMOFH pin output can also be set by TOLH in TCRF.

When TCF overflows from H'FFFF to H'0000, OVFH is set to 1 in TCSRF. If OVIEH in TCSRF and IENTFH in IENR2 are both 1, an interrupt request is sent to the CPU.

b. Operation in 8-bit timer mode

When CKSH2 is set to 1 in TCRF, TCF operates as two independent 8-bit timers, TCFH and TCFL. The TCFH/TCFL input clock is selected by CKSH2 to CKSH0/CKSL2 to CKSL0 in TCRF.

When the OCRFH/OCRFL and TCFH/TCFL values match, CMFH/CMFL is set to 1 in TCSRF. If IENTFH/IENTFL in IENR2 is 1, an interrupt request is sent to the CPU, and at the same time, TMOFH pin/TMOFL pin output is toggled. If CCLRH/CCLRL in TCSRF is 1, TCFH/TCFL is cleared. TMOFH pin/TMOFL pin output can also be set by TOLH/TOLL in TCRF.

When TCFH/TCFL overflows from H'FF to H'00, OVFH/OVFL is set to 1 in TCSRF. If OVIEH/OVIEL in TCSRF and IENTFH/IENTFL in IENR2 are both 1, an interrupt request is sent to the CPU.



(2) TCF Increment Timing

TCF is incremented by clock input (internal clock or external event input).

a. Internal clock operation

Bits CKSH2 to CKSH0 or CKSL2 to CKSL0 in TCRF select one of four internal clock sources $(\phi/32, \phi/16, \phi/4, \text{ or }\phi\text{w}/4)$ created by dividing the system clock $(\phi \text{ or }\phi\text{w})$.

b. External event operation

External event input is selected by clearing CKSL2 to 0 in TCRF. TCF can increment on either the rising or falling edge of external event input. External event edge selection is set by IEG3 in the interrupt controller's IEGR register. An external event pulse width of at least 2 system clocks (ϕ) is necessary. Shorter pulses will not be counted correctly.

(3) TMOFH/TMOFL Output Timing

In TMOFH/TMOFL output, the value set in TOLH/TOLL in TCRF is output. The output is toggled by the occurrence of a compare match. Figure 9.6 shows the output timing.

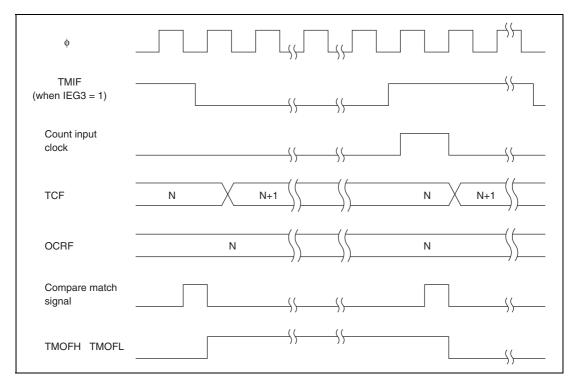


Figure 9.6 TMOFH/TMOFL Output Timing

(4) TCF Clear Timing

TCF can be cleared by a compare match with OCRF.

(5) Timer Overflow Flag (OVF) Set Timing

OVF is set to 1 when TCF overflows from H'FFFF to H'0000.

(6) Compare Match Flag set Timing

The compare match flag (CMFH or CMFL) is set to 1 when the TCF and OCRF values match. The compare match signal is generated in the last state during which the values match (when TCF is updated from the matching value to a new value). When TCF matches OCRF, the compare match signal is not generated until the next counter clock.

(7) Timer F Operation Modes

Timer F operation modes are shown in table 9.10.

Table 9.10 Timer F Operation Modes

Operation Mode	Reset	Active	Sleep	Watch	Sub- active	Sub- sleep	Standby	Module Standby
TCF	Reset	Functions	Functions	Functions/ Halted*	Functions/ Halted*	Functions/ Halted*	Halted	Halted
OCRF	Reset	Functions	Held	Held	Functions	Held	Held	Held
TCRF	Reset	Functions	Held	Held	Functions	Held	Held	Held
TCSRF	Reset	Functions	Held	Held	Functions	Held	Held	Held

Note: * When $\phi_w/4$ is selected as the TCF internal clock in active mode or sleep mode, since the system clock and internal clock are mutually asynchronous, synchronization is maintained by a synchronization circuit. This results in a maximum count cycle error of $1/\phi$ (s). When the counter is operated in subactive mode, watch mode, or subsleep mode, $\phi_w/4$ must be selected as the internal clock. The counter will not operate if any other internal clock is selected.



9.4.5 Application Notes

The following types of contention and operation can occur when timer F is used.

(1) 16-bit Timer Mode

In toggle output, TMOFH pin output is toggled when all 16 bits match and a compare match signal is generated. If a TCRF write by a MOV instruction and generation of the compare match signal occur simultaneously, TOLH data is output to the TMOFH pin as a result of the TCRF write. TMOFL pin output is unstable in 16-bit mode, and should not be used; the TMOFL pin should be used as a port pin.

If an OCRFL write and compare match signal generation occur simultaneously, the compare match signal is invalid. However, if the written data and the counter value match, a compare match signal will be generated at that point. As the compare match signal is output in synchronization with the TCFL clock, a compare match will not result in compare match signal generation if the clock is stopped.

Compare match flag CMFH is set when all 16 bits match and a compare match signal is generated. Compare match flag CMFL is set if the setting conditions for the lower 8 bits are satisfied.

When TCF overflows, OVFH is set. OVFL is set if the setting conditions are satisfied when the lower 8 bits overflow. If a TCFL write and overflow signal output occur simultaneously, the overflow signal is not output.

(2) 8-bit Timer Mode

a. TCFH, OCRFH

In toggle output, TMOFH pin output is toggled when a compare match occurs. If a TCRF write by a MOV instruction and generation of the compare match signal occur simultaneously, TOLH data is output to the TMOFH pin as a result of the TCRF write.

If an OCRFH write and compare match signal generation occur simultaneously, the compare match signal is invalid. However, if the written data and the counter value match, a compare match signal will be generated at that point. The compare match signal is output in synchronization with the TCFH clock.

If a TCFH write and overflow signal output occur simultaneously, the overflow signal is not output.



b. TCFL, OCRFL

In toggle output, TMOFL pin output is toggled when a compare match occurs. If a TCRF write by a MOV instruction and generation of the compare match signal occur simultaneously, TOLL data is output to the TMOFL pin as a result of the TCRF write.

If an OCRFL write and compare match signal generation occur simultaneously, the compare match signal is invalid. However, if the written data and the counter value match, a compare match signal will be generated at that point. As the compare match signal is output in synchronization with the TCFL clock, a compare match will not result in compare match signal generation if the clock is stopped.

If a TCFL write and overflow signal output occur simultaneously, the overflow signal is not output.

(3) Clear Timer FH, Timer FL Interrupt Request Flags (IRRTFH, IRRTFL), Timer Overflow Flags H, L (OVFH, OVFL) and Compare Match Flags H, L (CMFH, CMFL)

When ϕ w/4 is selected as the internal clock, "Interrupt factor generation signal" will be operated with ϕ w and the signal will be outputted with ϕ w width. And, "Overflow signal" and "Compare match signal" are controlled with 2 cycles of ϕ w signals. Those signals are outputted with 2 cycles width of ϕ w (figure 9.7)

In active (high-speed, medium-speed) mode, even if you cleared interrupt request flag during the term of validity of "Interrupt factor generation signal", same interrupt request flag is set. (figure 9.7 (1)) And, you cannot be cleared timer overflow flag and compare match flag during the term of validity of "Overflow signal" and "Compare match signal".

For interrupt request flag is set right after interrupt request is cleared, interrupt process to one time timer FH, timer FL interrupt might be repeated. (figure 9.7 (2)) Therefore, to definitely clear interrupt request flag in active (high-speed, medium-speed) mode, clear should be processed after the time that calculated with below (1) formula. And, to definitely clear timer overflow flag and compare match flag, clear should be processed after read timer control status register F (TCSRF) after the time that calculated with below (1) formula. For ST of (1) formula, please substitute the longest number of execution states in used instruction. (10 states of RTE instruction when MULXU, DIVXU instruction is not used, 14 states when MULXU, DIVXU instruction is used) In subactive mode, there are not limitation for interrupt request flag, timer overflow flag, and compare match flag clear.



The term of validity of "Interrupt factor generation signal"

- = 1 cycle of ϕ w + waiting time for completion of executing instruction
- + interrupt time synchronized with $\phi = 1/\phi w + ST \times (1/\phi) + (2/\phi)$ (second)....(1)

ST: Executing number of execution states

Method 1 is recommended to operate for time efficiency.

Method 1

- 1. Prohibit interrupt in interrupt handling routine (set IENFH, IENFL to 0).
- 2. After program process returned normal handling, clear interrupt request flags (IRRTFH, IRRTFL) after more than that calculated with (1) formula.
- 3. After read timer control status register F (TCSRF), clear timer overflow flags (OVFH, OVFL) and compare match flags (CMFH, CMFL).
- 4. Operate interrupt permission (set IENFH, IENFL to 1).

Method 2

- 1. Set interrupt handling routine time to more than time that calculated with (1) formula.
- 2. Clear interrupt request flags (IRRTFH, IRRTFL) at the end of interrupt handling routine.
- 3. After read timer control status register F (TCSRF), clear timer overflow flags (OVFH, OVFL) and compare match flags (CMFH, CMFL).

All above attentions are also applied in 16-bit mode and 8-bit mode.

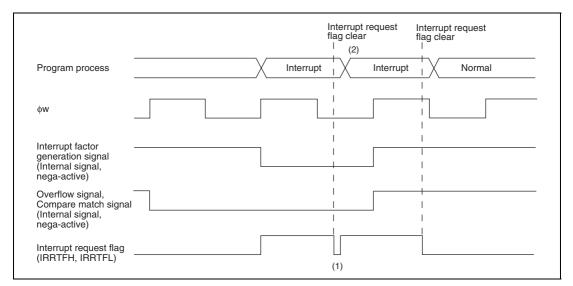


Figure 9.7 Clear Interrupt Request Flag when Interrupt Factor Generation Signal is Valid

(4) Timer Counter (TCF) Read/Write

When ϕ w/4 is selected as the internal clock in active (high-speed, medium-speed) mode, write on TCF is impossible. And, when read TCF, as the system clock and internal clock are mutually asynchronous, TCF synchronizes with synchronization circuit. This results in a maximum TCF read value error of ± 1 .

When read/write TCF in active (high-speed, medium-speed) mode is needed, please select internal clock except for ϕ w/4 before read/write.

In subactive mode, even ϕ w/4 is selected as the internal clock, normal read/write TCF is possible.

9.5 Timer G

9.5.1 Overview

Timer G is an 8-bit timer with dedicated input capture functions for the rising/falling edges of pulses input from the input capture input pin (input capture input signal). High-frequency component noise in the input capture input signal can be eliminated by a noise canceler, enabling accurate measurement of the input capture input signal duty cycle. If input capture input is not set, timer G functions as an 8-bit interval timer.

(1) Features

Features of timer G are given below.

- Choice of four internal clock sources ($\phi/64$, $\phi/32$, $\phi/2$, $\phi w/4$)
- · Dedicated input capture functions for rising and falling edges
- Level detection at counter overflow
 It is possible to detect whether overflow occurred when the input capture input signal was high or when it was low.
- Selection of whether or not the counter value is to be cleared at the input capture input signal rising edge, falling edge, or both edges
- Two interrupt sources: one input capture, one overflow. The input capture input signal rising or falling edge can be selected as the interrupt source.
- A built-in noise canceler eliminates high-frequency component noise in the input capture input signal.
- Watch mode, subactive mode and subsleep mode operation is possible when φw/4 is selected as the internal clock.
- Use of module standby mode enables this module to be placed in standby mode independently when not used.

(2) Block Diagram

Figure 9.8 shows a block diagram of timer G.

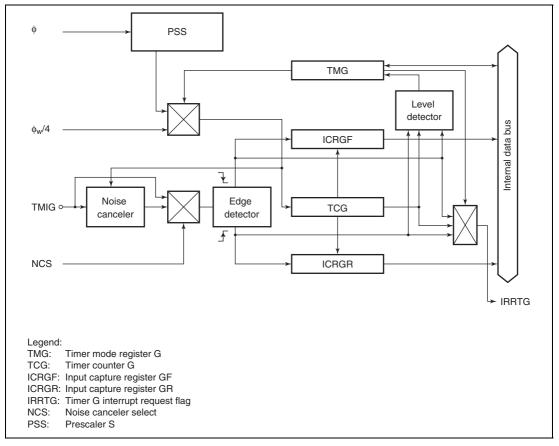


Figure 9.8 Block Diagram of Timer G

(3) Pin Configuration

Table 9.11 shows the timer G pin configuration.

Table 9.11 Pin Configuration

Name	Abbr.	I/O	Function
Input capture input	TMIG	Input	Input capture input pin



(4) Register Configuration

Table 9.12 shows the register configuration of timer G.

Table 9.12 Timer G Registers

Name	Abbr.	R/W	Initial Value	Address
Timer control register G	TMG	R/W	H'00	H'FFBC
Timer counter G	TCG	_	H'00	_
Input capture register GF	ICRGF	R	H'00	H'FFBD
Input capture register GR	ICRGR	R	H'00	H'FFBE
Clock stop register 1	CKSTPR1	R/W	H'FF	H'FFFA

9.5.2 Register Descriptions

(1) Timer Counter (TCG)

Bit	7	6	5	4	3	2	1	0
	TCG7	TCG6	TCG5	TCG4	TCG3	TCG2	TCG1	TCG0
Initial value	0	0	0	0	0	0	0	0
Read/Write	_	_	_	_	_	_	_	_

TCG is an 8-bit up-counter which is incremented by clock input. The input clock is selected by bits CKS1 and CKS0 in TMG.

TMIG in PMR1 is set to 1 to operate TCG as an input capture timer, or cleared to 0 to operate TCG as an interval timer*. In input capture timer operation, the TCG value can be cleared by the rising edge, falling edge, or both edges of the input capture input signal, according to the setting made in TMG.

When TCG overflows from H'FF to H'00, if OVIE in TMG is 1, IRRTG is set to 1 in IRR2, and if IENTG in IENR2 is 1, an interrupt request is sent to the CPU.

For details of the interrupt, see section 3.3, Interrupts.

TCG cannot be read or written by the CPU. It is initialized to H'00 upon reset.

Note: * An input capture signal may be generated when TMIG is modified.

(2) Input Capture Register GF (ICRGF)

Bit	7	6	5	4	3	2	1	0
	ICRGF7	ICRGF6	ICRGF5	ICRGF4	ICRGF3	ICRGF2	ICRGF1	ICRGF0
Initial value	0	0	0	0	0	0	0	0
Read/Write	R	R	R	R	R	R	R	R

ICRGF is an 8-bit read-only register. When a falling edge of the input capture input signal is detected, the current TCG value is transferred to ICRGF. If IIEGS in TMG is 1 at this time, IRRTG is set to 1 in IRR2, and if IENTG in IENR2 is 1, an interrupt request is sent to the CPU.

For details of the interrupt, see section 3.3, Interrupts.

To ensure dependable input capture operation, the pulse width of the input capture input signal must be at least 2ϕ or $2\phi_{\text{SUR}}$ (when the noise canceler is not used).

ICRGF is initialized to H'00 upon reset.

(3) Input Capture Register GR (ICRGR)

Bit	7	6	5	4	3	2	1	0
	ICRGR7	ICRGR6	ICRGR5	ICRGR4	ICRGR3	ICRGR2	ICRGR1	ICRGR0
Initial value	0	0	0	0	0	0	0	0
Read/Write	R	R	R	R	R	R	R	R

ICRGR is an 8-bit read-only register. When a rising edge of the input capture input signal is detected, the current TCG value is transferred to ICRGR. If IIEGS in TMG is 1 at this time, IRRTG is set to 1 in IRR2, and if IENTG in IENR2 is 1, an interrupt request is sent to the CPU.

For details of the interrupt, see section 3.3, Interrupts.

To ensure dependable input capture operation, the pulse width of the input capture input signal must be at least 2ϕ or $2\phi_{\text{SUR}}$ (when the noise canceler is not used).

ICRGR is initialized to H'00 upon reset.



(4) Timer Mode Register G (TMG)

Bit:	7	6	5	4	3	2	1	0
	OVFH	OVFL	OVIE	IIEGS	CCLR1	CCLR0	CKS1	CKS0
Initial value:	0	0	0	0	0	0	0	0
Read/Write:	R/W*	R/W*	R/W	R/W	R/W	R/W	R/W	R/W

Note: * Bits 7 and 6 can only be written with 0, for flag clearing.

TMG is an 8-bit read/write register that performs TCG clock selection from four internal clock sources, counter clear selection, and edge selection for the input capture input signal interrupt request, controls enabling of overflow interrupt requests, and also contains the overflow flags.

TMG is initialized to H'00 upon reset.

Bit 7: Timer overflow flag H (OVFH)

Bit 7 is a status flag indicating that TCG has overflowed from H'FF to H'00 when the input capture input signal is high. This flag is set by hardware and cleared by software. It cannot be set by software.

Bit 7 OVFH	Description	
0	Clearing condition:	(initial value)
	After reading OVFH = 1, cleared by writing 0 to OVFH	
1	Setting condition:	
	Set when TCG overflows from H'FF to H'00	

Bit 6: Timer overflow flag L (OVFL)

Bit 6 is a status flag indicating that TCG has overflowed from H'FF to H'00 when the input capture input signal is low, or in interval operation. This flag is set by hardware and cleared by software. It cannot be set by software.

Bit 6 OVFL	Description	
0	Clearing condition:	(initial value)
	After reading OVFL = 1, cleared by writing 0 to OVFL	
1	Setting condition:	
	Set when TCG overflows from H'FF to H'00	

Bit 5: Timer overflow interrupt enable (OVIE)

Bit 5 selects enabling or disabling of interrupt generation when TCG overflows.

Bit 5 OVIE	Description	
0	TCG overflow interrupt request is disabled	(initial value)
1	TCG overflow interrupt request is enabled	

Bit 4: Input capture interrupt edge select (IIEGS)

Bit 4 selects the input capture input signal edge that generates an interrupt request.

Bit 4 IIEGS	Description	
0	Interrupt generated on rising edge of input capture input signal	(initial value)
1	Interrupt generated on falling edge of input capture input signal	

Bits 3 and 2: Counter clear 1 and 0 (CCLR1, CCLR0)

Bits 3 and 2 specify whether or not TCG is cleared by the rising edge, falling edge, or both edges of the input capture input signal.

Bit 3 CCLR1	Bit 2 CCLR0	Description	
0	0	TCG clearing is disabled	(initial value)
0	1	TCG cleared by falling edge of input capture input signal	_
1	0	TCG cleared by rising edge of input capture input signal	_
1	1	TCG cleared by both edges of input capture input signal	

Bits 1 and 0: Clock select (CKS1, CKS0)

Bits 1 and 0 select the clock input to TCG from among four internal clock sources.

Bit 1 CKS1	Bit 0 CKS0	Description	
0	0	Internal clock: Counting on φ/64	(initial value)
0	1	Internal clock: Counting on \$\phi/32\$	
1	0	Internal clock: Counting on φ/2	
1	1	Internal clock: Counting on \psiw/4	

(5) Clock Stop Register 1 (CKSTPR1)

Bit	7	6	5	4	3	2	1	0
	_	S31CKSTP	S32CKSTP	ADCKSTP	TGCKSTP	TFCKSTP	TCCKSTP	TACKSTP
Initial value	1	1	1	1	1	1	1	1
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W

CKSTPR1 is an 8-bit read/write register that performs module standby mode control for peripheral modules. Only the bit relating to timer G is described here. For details of the other bits, see the sections on the relevant modules.

Bit 3: Timer G module standby mode control (TGCKSTP)

Bit 3 controls setting and clearing of module standby mode for timer G.

TGCKSTP	Description	
0	Timer G is set to module standby mode	_
1	Timer G module standby mode is cleared	(initial value)

9.5.3 Noise Canceler

The noise canceler consists of a digital low-pass filter that eliminates high-frequency component noise from the pulses input from the input capture input pin. The noise canceler is set by NCS* in PMR3.

Figure 9.9 shows a block diagram of the noise canceler.

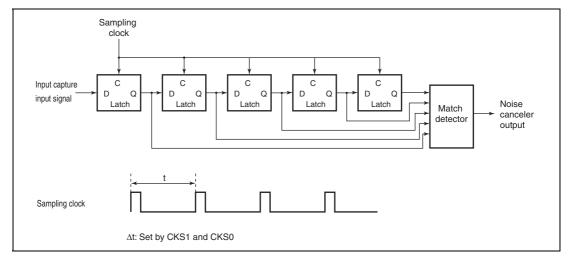


Figure 9.9 Noise Canceler Block Diagram

The noise canceler consists of five latch circuits connected in series and a match detector circuit. When the noise cancellation function is not used (NCS = 0), the system clock is selected as the sampling clock. When the noise cancellation function is used (NCS = 1), the sampling clock is the internal clock selected by CKS1 and CKS0 in TMG, the input capture input is sampled on the rising edge of this clock, and the data is judged to be correct when all the latch outputs match. If all the outputs do not match, the previous value is retained. After a reset, the noise canceler output is initialized when the falling edge of the input capture input signal has been sampled five times. Therefore, after making a setting for use of the noise cancellation function, a pulse with at least five times the width of the sampling clock is a dependable input capture signal. Even if noise cancellation is not used, an input capture input signal pulse width of at least 2ϕ or $2\phi_{\text{SUB}}$ is necessary to ensure that input capture operations are performed properly

Note: * An input capture signal may be generated when the NCS bit is modified.

Figure 9.10 shows an example of noise canceler timing.

In this example, high-level input of less than five times the width of the sampling clock at the input capture input pin is eliminated as noise.

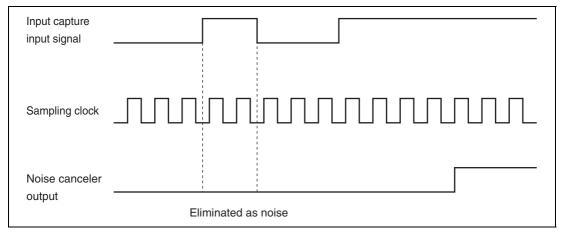


Figure 9.10 Noise Canceler Timing (Example)

9.5.4 Operation

Timer G is an 8-bit timer with built-in input capture and interval functions.

(1) Timer G Functions

Timer G is an 8-bit up-counter with two functions, an input capture timer function and an interval timer function.

The operation of these two functions is described below.

a. Input capture timer operation

When the TMIG bit is set to 1 in port mode register 1 (PMR1), timer G functions as an input capture timer*.

In a reset, timer mode register G (TMG), timer counter G (TCG), input capture register GF (ICRGF), and input capture register GR (ICRGR) are all initialized to H'00.

Following a reset, TCG starts incrementing on the φ/64 internal clock.

The input clock can be selected from four internal clock sources by bits CKS1 and CKS0 in TMG.

When a rising edge/falling edge is detected in the input capture signal input from the TMIG pin, the TCG value at that time is transferred to ICRGR/ICRGF. When the edge selected by IIEGS in TMG is input, IRRTG is set to 1 in IRR2, and if the IENTG bit in IENR2 is 1 at this time, an interrupt request is sent to the CPU. For details of the interrupt, see section 3.3, Interrupts.

TCG can be cleared by a rising edge, falling edge, or both edges of the input capture signal, according to the setting of bits CCLR1 and CCLR0 in TMG. If TCG overflows when the input capture signal is high, the OVFH bit is set in TMG; if TCG overflows when the input capture signal is low, the OVFL bit is set in TMG. If the OVIE bit in TMG is 1 when these bits are set, IRRTG is set to 1 in IRR2, and if the IENTG bit in IENR2 is 1, timer G sends an interrupt request to the CPU. For details of the interrupt, see section 3.3, Interrupts.

Timer G has a built-in noise canceler that enables high-frequency component noise to be eliminated from pulses input from the TMIG pin. For details, see section 9.5.3, Noise Canceler.

Note: * An input capture signal may be generated when TMIG is modified.

b. Interval timer operation

When the TMIG bit is cleared to 0 in PMR1, timer G functions as an interval timer. Following a reset, TCG starts incrementing on the $\phi/64$ internal clock. The input clock can be selected from four internal clock sources by bits CKS1 and CKS0 in TMG. TCG increments on the selected clock, and when it overflows from H'FF to H'00, the OVFL bit is set to 1 in TMG. If the OVIE bit in TMG is 1 at this time, IRRTG is set to 1 in IRR2, and if the IENTG bit in IENR2 is 1, timer G sends an interrupt request to the CPU. For details of the interrupt, see section 3.3, Interrupts.

(2) Increment Timing

TCG is incremented by internal clock input. Bits CKS1 and CKS0 in TMG select one of four internal clock sources ($\phi/64$, $\phi/32$, $\phi/2$, or $\phi w/4$) created by dividing the system clock (ϕ) or watch clock (ϕw).

(3) Input Capture Input Timing

a. Without noise cancellation function

For input capture input, dedicated input capture functions are provided for rising and falling edges.

Figure 9.11 shows the timing for rising/falling edge input capture input.

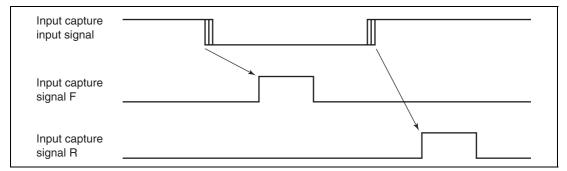


Figure 9.11 Input Capture Input Timing (without Noise Cancellation Function)

b. With noise cancellation function

When noise cancellation is performed on the input capture input, the passage of the input capture signal through the noise canceler results in a delay of five sampling clock cycles from the input capture input signal edge.

Figure 9.12 shows the timing in this case.



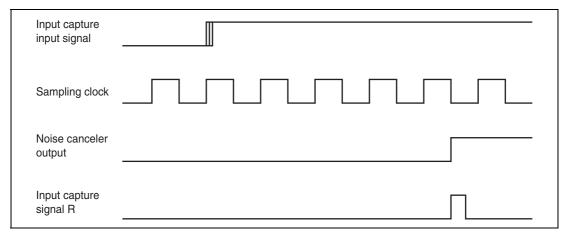


Figure 9.12 Input Capture Input Timing (with Noise Cancellation Function)

(4) Timing of Input Capture by Input Capture Input

Figure 9.13 shows the timing of input capture by input capture input

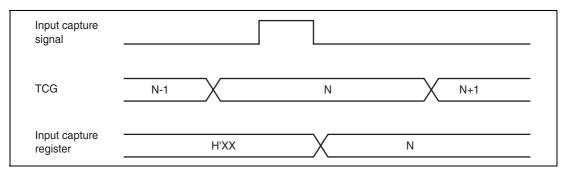


Figure 9.13 Timing of Input Capture by Input Capture Input

(5) TGC Clear Timing

TCG can be cleared by the rising edge, falling edge, or both edges of the input capture input signal.

Figure 9.14 shows the timing for clearing by both edges.

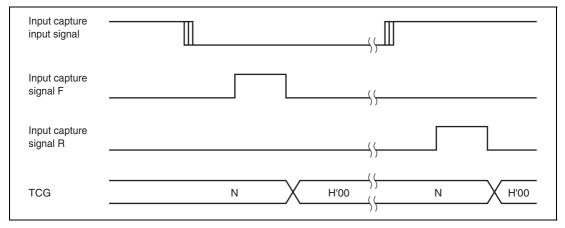


Figure 9.14 TCG Clear Timing

(6) Timer G Operation Modes

Timer G operation modes are shown in table 9.13.

Table 9.13 Timer G Operation Modes

Opera	ation Mode	Reset	Active	Sleep	Watch	Sub- active	Sub- sleep	Standby	Module Standby
TCG	Input capture	Reset	Functions*	Functions*	Functions/ halted*	Functions/ halted*	Functions/ halted*	Halted	Halted
	Interval	Reset	Functions*	Functions*	Functions/ halted*	Functions/ halted*	Functions/ halted*	Halted	Halted
ICRG	F	Reset	Functions*	Functions*	Functions/ halted*	Functions/ halted*	Functions/ halted*	Held	Held
ICRG	R	Reset	Functions*	Functions*	Functions halted*	Functions/ halted*	Functions/ halted*	Held	Held
TMG		Reset	Functions	Held	Held	Functions	Held	Held	Held

Note: * When φw/4 is selected as the TCG internal clock in active mode or sleep mode, since the system clock and internal clock are mutually asynchronous, synchronization is maintained by a synchronization circuit. This results in a maximum count cycle error of 1/φ (s). When φw/4 is selected as the TCG internal clock in watch mode, TCG and the noise canceler operate on the φw/4 internal clock without regard to the φ_{SUB} subclock (φw/8, φw/4, φw/2). Note that when another internal clock is selected, TCG and the noise canceler do not operate, and input of the input capture input signal does not result in input capture.

To be operated Timer G in subactive mode or subsleep mode, select ϕ w/4 for internal clock of TCG and also select ϕ w/2 for sub clock ϕ _{SUB}. When another internal clock is selected and when another sub clock (ϕ w/8, ϕ w/4) is selected, TCG and noise canceler do not operate.

9.5.5 Application Notes

(1) Internal Clock Switching and TCG Operation

Depending on the timing, TCG may be incremented by a switch between difference internal clock sources. Table 9.14 shows the relation between internal clock switchover timing (by write to bits CKS1 and CKS0) and TCG operation.

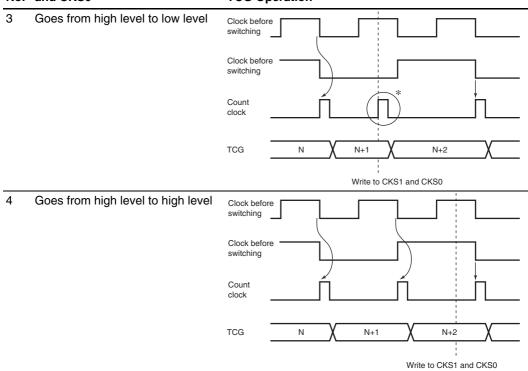
When TCG is internally clocked, an increment pulse is generated on detection of the falling edge of an internal clock signal, which is divided from the system clock (ϕ) or subclock (ϕ w). For this reason, in a case like No. 3 in table 9.14 where the switch is from a high clock signal to a low clock signal, the switchover is seen as a falling edge, causing TCG to increment.

Table 9.14 Internal Clock Switching and TCG Operation

No.	Clock Levels Before and After Modifying Bits CKS1 and CKS0	TCG Operation
1	Goes from low level to low level	Clock before switching
		Clock after switching
		Count clock
		TCG N N+1
		Write to CKS1 and CKS0
2	Goes from low level to high level	Clock before switching
		Clock before switching
		Count clock
		TCG N N+1 N+2 N+2
		Write to CKS1 and CKS0

Clock Levels Before and After Modifying Bits CKS1 No. and CKS0

TCG Operation



Note: * The switchover is seen as a falling edge, and TCG is incremented.

(2) Notes on Port Mode Register Modification

The following points should be noted when a port mode register is modified to switch the input capture function or the input capture input noise canceler function.

• Switching input capture input pin function

Note that when the pin function is switched by modifying TMIG in port mode register 1 (PMR1), which performs input capture input pin control, an edge will be regarded as having been input at the pin even though no valid edge has actually been input. Input capture input signal input edges, and the conditions for their occurrence, are summarized in table 9.15.

Table 9.15 Input Capture Input Signal Input Edges Due to Input Capture Input Pin Switching, and Conditions for Their Occurrence

Input Capture Input Signal Input Edge	Conditions
Generation of rising edge	When TMIG is modified from 0 to 1 while the TMIG pin is high
	When NCS is modified from 0 to 1 while the TMIG pin is high, then TMIG is modified from 0 to 1 before the signal is sampled five times by the noise canceler
Generation of falling edge	When TMIG is modified from 1 to 0 while the TMIG pin is high
	When NCS is modified from 0 to 1 while the TMIG pin is low, then TMIG is modified from 0 to 1 before the signal is sampled five times by the noise canceler
	When NCS is modified from 0 to 1 while the TMIG pin is high, then TMIG is modified from 1 to 0 after the signal is sampled five times by the noise canceler

Note: When the P1₃ pin is not set as an input capture input pin, the timer G input capture input signal is low.

Switching input capture input noise canceler function

When performing noise canceler function switching by modifying NCS in port mode register 3 (PMR3), which controls the input capture input noise canceler, TMIG should first be cleared to 0. Note that if NCS is modified without first clearing TMIG, an edge will be regarded as having been input at the pin even though no valid edge has actually been input. Input capture input signal input edges, and the conditions for their occurrence, are summarized in table 9.16.



Table 9.16 Input Capture Input Signal Input Edges Due to Noise Canceler Function Switching, and Conditions for Their Occurrence

Input Capture Input Signal Input Edge	Conditions
Generation of rising edge	When the TMIG pin level is switched from low to high while TMIG is set to 1, then NCS is modified from 0 to 1 before the signal is sampled five times by the noise canceler
Generation of falling edge	When the TMIG pin level is switched from high to low while TMIG is set to 1, then NCS is modified from 1 to 0 before the signal is sampled five times by the noise canceler

When the pin function is switched and an edge is generated in the input capture input signal, if this edge matches the edge selected by the input capture interrupt select (IIEGS) bit, the interrupt request flag will be set to 1. The interrupt request flag should therefore be cleared to 0 before use. Figure 9.15 shows the procedure for port mode register manipulation and interrupt request flag clearing. When switching the pin function, set the interrupt-disabled state before manipulating the port mode register, then, after the port mode register operation has been performed, wait for the time required to confirm the input capture input signal as an input capture signal (at least two system clocks when the noise canceler is not used; at least five sampling clocks when the noise canceler is used), before clearing the interrupt enable flag to 0. There are two ways of preventing interrupt request flag setting when the pin function is switched: by controlling the pin level so that the conditions shown in tables 9.15 and 9.16 are not satisfied, or by setting the opposite of the generated edge in the IIEGS bit in TMG.

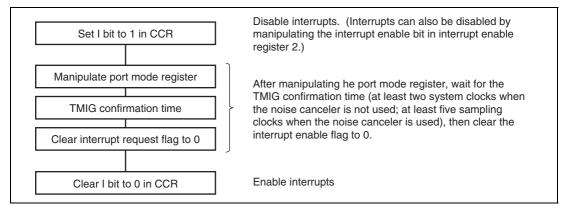


Figure 9.15 Port Mode Register Manipulation and Interrupt Enable Flag Clearing Procedure

9.5.6 Timer G Application Example

Using timer G, it is possible to measure the high and low widths of the input capture input signal as absolute values. For this purpose, CCLR1 and CCLR0 should both be set to 1 in TMG.

Figure 9.16 shows an example of the operation in this case.

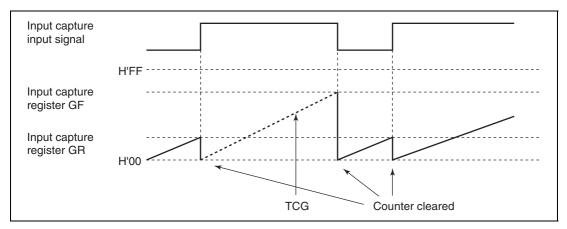


Figure 9.16 Timer G Application Example

9.6 Watchdog Timer

9.6.1 Overview

The watchdog timer has an 8-bit counter that is incremented by an input clock. If a system runaway allows the counter value to overflow before being rewritten, the watchdog timer can reset the chip internally.

(1) Features

Features of the watchdog timer are given below.

- Incremented by internal clock source (φ/8192 or φw/32).
- A reset signal is generated when the counter overflows. The overflow period can be set from from 1 to 256 times 8192/φ or 32/φw (from approximately 4 ms to 1000 ms when φ = 2.00 MHz).
- Use of module standby mode enables this module to be placed in standby mode independently
 when not used.

(2) Block Diagram

Figure 9.17 shows a block diagram of the watchdog timer.

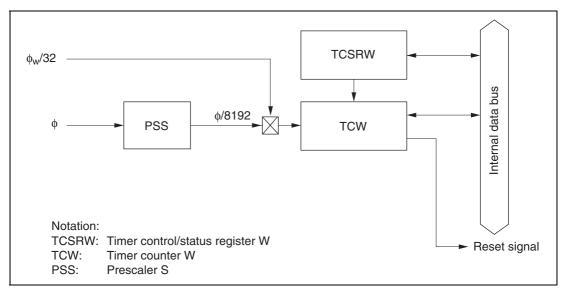


Figure 9.17 Block Diagram of Watchdog Timer

(3) Register Configuration

Table 9.17 shows the register configuration of the watchdog timer.

Table 9.17 Watchdog Timer Registers

Name	Abbr.	R/W	Initial Value	Address
Timer control/status register W	TCSRW	R/W	H'AA	H'FFB2
Timer counter W	TCW	R/W	H'00	H'FFB3
Clock stop register 2	CKSTPR2	R/W	H'FF	H'FFFB
Port mode register 3	PMR3	R/W	H'04	H'FFCA

9.6.2 Register Descriptions

(1) Timer Control/Status Register W (TCSRW)

Bit	7	6	5	4	3	2	1	0
	B6WI	TCWE	B4WI	TCSRWE	B2WI	WDON	B0WI	WRST
Initial value	1	0	1	0	1	0	1	0
Read/Write	R	R/(W)*	R	R/(W)*	R	R/(W)*	R	R/(W)*

Note: * Write is permitted only under certain conditions, which are given in the descriptions of the individual bits.

TCSRW is an 8-bit read/write register that controls write access to TCW and TCSRW itself, controls watchdog timer operations, and indicates operating status.

Bit 7: Bit 6 write inhibit (B6WI)

Bit 7 controls the writing of data to bit 6 in TCSRW.

Bit 7 B6WI	Description	
0	Bit 6 is write-enabled	
1	Bit 6 is write-protected	(initial value)

This bit is always read as 1. Data written to this bit is not stored.



Bit 6: Timer counter W write enable (TCWE)

Bit 6 controls the writing of data to TCW.

Bit 6 TCWE	Description	
0	Data cannot be written to TCW	(initial value)
1	Data can be written to TCW	

Bit 5: Bit 4 write inhibit (B4WI)

Bit 5 controls the writing of data to bit 4 in TCSRW.

Bit 5 B4WI	Description	
0	Bit 4 is write-enabled	
1	Bit 4 is write-protected	(initial value)

This bit is always read as 1. Data written to this bit is not stored.

Bit 4: Timer control/status register W write enable (TCSRWE)

Bit 4 controls the writing of data to TCSRW bits 2 and 0.

Bit 4 TCSRWE	Description	
0	Data cannot be written to bits 2 and 0	(initial value)
1	Data can be written to bits 2 and 0	

Bit 3: Bit 2 write inhibit (B2WI)

Bit 3 controls the writing of data to bit 2 in TCSRW.

Bit 3		
B2WI	Description	
0	Bit 2 is write-enabled	
1	Bit 2 is write-protected	(initial value)

This bit is always read as 1. Data written to this bit is not stored.



Bit 2: Watchdog timer on (WDON)

Bit 2 enables watchdog timer operation.

Bit 2		
WDON	Description	
0	Watchdog timer operation is disabled	(initial value)
	Clearing condition:	
	Reset, or when TCSRWE = 1 and 0 is written in both B2WI and WDO	N
1	Watchdog timer operation is enabled	
	Setting condition:	
	When TCSRWE = 1 and 0 is written in B2WI and 1 is written in WDOI	V

Counting starts when this bit is set to 1, and stops when this bit is cleared to 0.

Bit 1: Bit 0 write inhibit (B0WI)

Bit 1 controls the writing of data to bit 0 in TCSRW.

Bit 1 B0WI	Description	
0	Bit 0 is write-enabled	_
1	Bit 0 is write-protected	(initial value)

This bit is always read as 1. Data written to this bit is not stored.

Bit 0: Watchdog timer reset (WRST)

Bit 0 indicates that TCW has overflowed, generating an internal reset signal. The internal reset signal generated by the overflow resets the entire chip. WRST is cleared to 0 by a reset from the RES pin, or when software writes 0.



Bit 0 WRST	Description
0	Clearing condition:
	Reset by RES pin
	When TCSRWE = 1, and 0 is written in both B0WI and WRST
1	Setting condition:
	When TCW overflows and an internal reset signal is generated

(2) Timer Counter W (TCW)

Bit	7	6	5	4	3	2	1	0
	TCW7	TCW6	TCW5	TCW4	TCW3	TCW2	TCW1	TCW0
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W							

TCW is an 8-bit read/write up-counter, which is incremented by internal clock input. The input clock is $\phi/8192$ or $\phi w/32$. The TCW value can always be written or read by the CPU.

When TCW overflows from H'FF to H'00, an internal reset signal is generated and WRST is set to 1 in TCSRW. Upon reset, TCW is initialized to H'00.

(3) Clock Stop Register 2 (CKSTPR2)

Bit	7	6	5	4	3	2	1	0
	_	_	_	_	AECKSTP	WDCKSTP	PWCKSTP	LDCKSTP
Initial value	1	1	1	1	1	1	1	1
Read/Write	_	_		_	R/W	R/W	R/W	R/W

CKSTPR2 is an 8-bit read/write register that performs module standby mode control for peripheral modules. Only the bit relating to the watchdog timer is described here. For details of the other bits, see the sections on the relevant modules.

Bit 2: Watchdog timer module standby mode control (WDCKSTP)

Bit 2 controls setting and clearing of module standby mode for the watchdog timer.

WDCKSTP Description

0	Watchdog timer is set to module standby mode	
1	Watchdog timer module standby mode is cleared	(initial value)

Note: WDCKSTP is valid when the WDON bit is cleared to 0 in timer control/status register W (TCSRW). If WDCKSTP is set to 0 while WDON is set to 1 (during watchdog timer operation), 0 will be set in WDCKSTP but the watchdog timer will continue its watchdog function and will not enter module standby mode. When the watchdog function ends and WDON is cleared to 0 by software, the WDCKSTP setting will become valid and the watchdog timer will enter module standby mode.

(4) Port Mode Register 3 (PMR3)

Bit	7	6	5	4	3	2	1	0
	AEVL	AEVH	WDCKS	NCS	IRQ0	_	UD	PWM
Initial value	0	0	0	0	0	1	0	0
Read/Write	R/W	R/W	R/W	R/W	R/W	_	R/W	R/W

PMR3 is an 8-bit read/write register, mainly controlling the selection of pin functions for port 3 pins. Only the bit relating to the watchdog timer is described here. For details of the other bits, see section 8, I/O Ports.

Bit 5: Watchdog timer source clock select (WDCKS)

WDCKS	Description	
0	φ/8192 selected	(initial value)
1	φw/32 selected	



9.6.3 Timer Operation

The watchdog timer has an 8-bit counter (TCW) that is incremented by clock input (ϕ /8192 or ϕ w/32). The input clock is selected by bit WDCKS in port mode register 3 (PMR3): ϕ /8192 is selected when WDCKS is cleared to 0, and ϕ w/32 when set to 1. When TCSRWE = 1 in TCSRW, if 0 is written in B2WI and 1 is simultaneously written in WDON, TCW starts counting up. When the TCW count reaches H'FF, the next clock input causes the watchdog timer to overflow, and an internal reset signal is generated one reference clock (ϕ or ϕ_{SUB}) cycle later. The internal reset signal is output for 512 clock cycles of the ϕ_{OSC} clock. It is possible to write to TCW, causing TCW to count up from the written value. The overflow period can be set in the range from 1 to 256 input clocks, depending on the value written in TCW.

Figure 9.18 shows an example of watchdog timer operations.

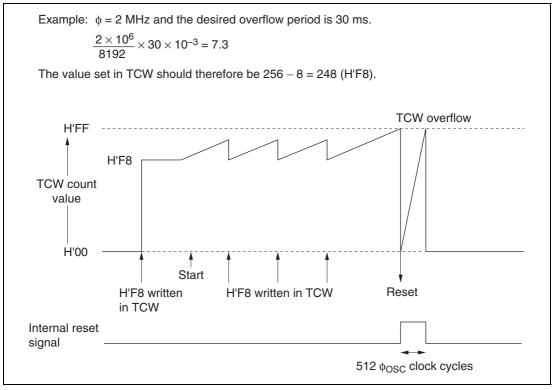


Figure 9.18 Typical Watchdog Timer Operations (Example)

9.6.4 Watchdog Timer Operation States

Table 9.18 summarizes the watchdog timer operation states.

Table 9.18 Watchdog Timer Operation States

Operation Mode	Reset	Active	Sleep	Watch	Sub- active	Sub- sleep	Standby	Module Standby
TCW	Reset	Functions	Functions	Halted	Functions/ Halted*	Halted	Halted	Halted
TCSRW	Reset	Functions	Functions	Retained	Functions/ Halted*	Retained	Retained	Retained

Note: * Functions when $\phi w/32$ is selected as the input clock.

9.7 Asynchronous Event Counter (AEC)

9.7.1 Overview

The asynchronous event counter is incremented by external event clock input.

(1) Features

Features of the asynchronous event counter are given below.

• Can count asynchronous events

Can count external events input asynchronously without regard to the operation of base clocks ϕ and $\phi_{\text{\tiny SUIR}}$.

The counter has a 16-bit configuration, enabling it to count up to 65536 (2¹⁶) events.

- Can also be used as two independent 8-bit event counter channels.
- Counter resetting and halting of the count-up function controllable by software
- Automatic interrupt generation on detection of event counter overflow
- Use of module standby mode enables this module to be placed in standby mode independently when not used.

(2) Block Diagram

Figure 9.19 shows a block diagram of the asynchronous event counter.

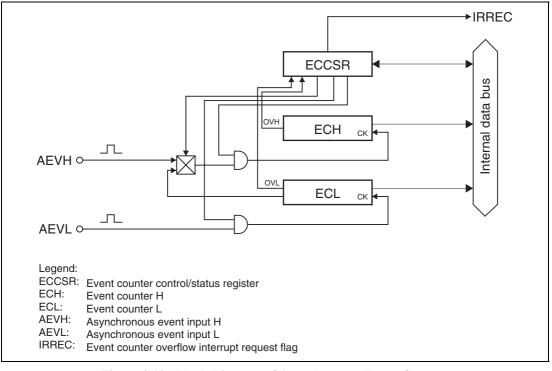


Figure 9.19 Block Diagram of Asynchronous Event Counter

(3) Pin Configuration

Table 9.19 shows the asynchronous event counter pin configuration.

Table 9.19 Pin Configuration

Name	Abbr.	I/O	Function
Asynchronous event input H	AEVH	Input	Event input pin for input to event counter H
Asynchronous event input L	AEVL	Input	Event input pin for input to event counter L

(4) Register Configuration

Table 9.20 shows the register configuration of the asynchronous event counter.

Table 9.20 Asynchronous Event Counter Registers

Name	Abbr.	R/W	Initial Value	Address
Event counter control/status register	ECCSR	R/W	H'00	H'FF95
Event counter H	ECH	R	H'00	H'FF96
Event counter L	ECL	R	H'00	H'FF97
Clock stop register 2	CKSTPR2	R/W	H'FF	H'FFFB

9.7.2 Register Descriptions

(1) Event Counter Control/Status Register (ECCSR)

Bit	7	6	5	4	3	2	1	0
	OVH	OVL	_	CH2	CUEH	CUEL	CRCH	CRCL
Initial Value	0	0	0	0	0	0	0	0
Read/Write	R/W*	R/W*	R/W	R/W	R/W	R/W	R/W	R/W

Note: * Bits 7 and 6 can only be written with 0, for flag clearing.

ECCSR is an 8-bit read/write register that controls counter overflow detection, counter resetting, and halting of the count-up function.

ECCSR is initialized to H'00 upon reset.

Bit 7: Counter overflow flag H (OVH)

Bit 7 is a status flag indicating that ECH has overflowed from H'FF to H'00. This flag is set when ECH overflows. It is cleared by software but cannot be set by software. OVH is cleared by reading it when set to 1, then writing 0.

When ECH and ECL are used as a 16-bit event counter with CH2 cleared to 0, OVH functions as a status flag indicating that the 16-bit event counter has overflowed from H'FFFF to H'0000.

Bit 7		
OVH	Description	
0	ECH has not overflowed	(initial value)
	Clearing condition:	
	After reading OVH = 1, cleared by writing 0 to OVH	
1	ECH has overflowed	
	Setting condition:	
	Set when ECH overflows from H'FF to H'00	

Bit 6: Counter overflow flag L (OVL)

Bit 6 is a status flag indicating that ECL has overflowed from H'FF to H'00. This flag is set when ECL overflows. It is cleared by software but cannot be set by software. OVL is cleared by reading it when set to 1, then writing 0.

Bit 6 OVL	Description	
0	ECL has not overflowed	(initial value)
	Clearing condition:	
	After reading OVL = 1, cleared by writing 0 to OVL	
1	ECL has overflowed	
	Setting condition:	
	Set when ECL overflows from H'FF to H'00 while CH2 is set to 1	

Bit 5: Reserved bit

Bit 5 is reserved; it can be read and written, and is initialized to 0 upon reset.



Bit 4: Channel select (CH2)

Bit 4 selects whether ECH and ECL are used as a single-channel 16-bit event counter or as two independent 8-bit event counter channels. When CH2 is cleared to 0, ECH and ECL function as a 16-bit event counter which is incremented each time an event clock is input to the AEVL pin as asynchronous event input. In this case, the overflow signal from ECL is selected as the ECH input clock. When CH2 is set to 1, ECH and ECL function as independent 8-bit event counters which are incremented each time an event clock is input to the AEVH or AEVL pin, respectively, as asynchronous event input.

Bit 4 CH2	Description
0	ECH and ECL are used together as a single-channel 16-bit event counter (initial value)
1	ECH and ECL are used as two independent 8-bit event counter channels

Bit 3: Count-up enable H (CUEH)

Bit 3 enables event clock input to ECH. When 1 is written to this bit, event clock input is enabled and increments the counter. When 0 is written to this bit, event clock input is disabled and the ECH value is held. The AEVH pin or the ECL overflow signal can be selected as the event clock source by bit CH2.

Bit 3 CUEH	Description	
0	ECH event clock input is disabled	(initial value)
	ECH value is held	
1	ECH event clock input is enabled	

Bit 2: Count-up enable L (CUEL)

Bit 2 enables event clock input to ECL. When 1 is written to this bit, event clock input is enabled and increments the counter. When 0 is written to this bit, event clock input is disabled and the ECL value is held.

Bit 2 CUEL	Description	
0	ECL event clock input is disabled ECL value is held	(initial value)
1	ECL event clock input is enabled	

Bit 1: Counter reset control H (CRCH)

Bit 1 controls resetting of ECH. When this bit is cleared to 0, ECH is reset. When 1 is written to this bit, the counter reset is cleared and the ECH count-up function is enabled.

Bit 1 CRCH	Description	
0	ECH is reset	(initial value)
1	ECH reset is cleared and count-up function is enabled	

Bit 0: Counter reset control L (CRCL)

Bit 0 controls resetting of ECL. When this bit is cleared to 0, ECL is reset. When 1 is written to this bit, the counter reset is cleared and the ECL count-up function is enabled.

Bit 0 CRCL	Description	
0	ECL is reset	(initial value)
1	ECL reset is cleared and count-up function is enabled	

(2) Event Counter H (ECH)

Bit	7	6	5	4	3	2	1	0
	ECH7	ECH6	ECH5	ECH4	ECH3	ECH2	ECH1	ECH0
Initial Value	0	0	0	0	0	0	0	0
Read/Write	R	R	R	R	R	R	R	R

ECH is an 8-bit read-only up-counter that operates either as an independent 8-bit event counter or as the upper 8-bit up-counter of a 16-bit event counter configured in combination with ECL. Either the external asynchronous event AEVH pin or the overflow signal from lower 8-bit counter ECL can be selected as the input clock source by bit CH2. ECH can be cleared to H'00 by software, and is also initialized to H'00 upon reset.



(3) Event Counter L (ECL)

ECL is an 8-bit read-only up-counter that operates either as an independent 8-bit event counter or as the lower 8-bit up-counter of a 16-bit event counter configured in combination with ECH. The event clock from the external asynchronous event AEVL pin is used as the input clock source. ECL can be cleared to H'00 by software, and is also initialized to H'00 upon reset.

Bit	7	6	5	4	3	2	1	0
	ECL7	ECL6	ECL5	ECL4	ECL3	ECL2	ECL1	ECL0
Initial Value	0	0	0	0	0	0	0	0
Read/Write	R	R	R	R	R	R	R	R

(4) Clock Stop Register 2 (CKSTPR2)

Bit	7	6	5	4	3	2	1	0
	_		_	_	AECKSTP	WDCKSTP	PWCKSTP	LDCKSTP
Initial value	1	1	1	1	1	1	1	1
Read/Write	_	_	_	_	R/W	R/W	R/W	R/W

CKSTPR2 is an 8-bit read/write register that performs module standby mode control for peripheral modules. Only the bit relating to the asynchronous event counter is described here. For details of the other bits, see the sections on the relevant modules.

Bit 3: Asynchronous event counter module standby mode control (AECKSTP)

Bit 3 controls setting and clearing of module standby mode for the asynchronous event counter.

AECKSTP	Description	
0	Asynchronous event counter is set to module standby mode	
1	Asynchronous event counter module standby mode is cleared	(initial value)

9.7.3 Operation

(1) 16-bit Event Counter Operation

When bit CH2 is cleared to 0 in ECCSR, ECH and ECL operate as a 16-bit event counter. Figure 9.20 shows an example of the software processing when ECH and ECL are used as a 16-bit event counter.

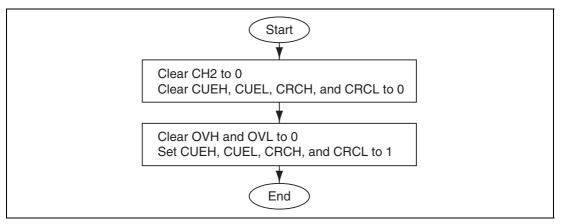


Figure 9.20 Example of Software Processing when Using ECH and ECL as 16-Bit Event Counter

As CH2 is cleared to 0 by a reset, ECH and ECL operate as a 16-bit event counter after a reset. They can also be used as a 16-bit event counter by carrying out the software processing shown in the example in figure 9.20. The operating clock source is asynchronous event input from the AEVL pin. When the next clock is input after the count value reaches H'FF in both ECH and ECL, ECH and ECL overflow from H'FFFF to H'0000, the OVH flag is set to 1 in ECCSR, the ECH and ECL count values each return to H'00, and counting up is restarted. When overflow occurs, the IRREC bit is set to 1 in IRR2. If the IENEC bit in IENR2 is 1 at this time, an interrupt request is sent to the CPU.

(2) 8-bit Event Counter Operation

When bit CH2 is set to 1 in ECCSR, ECH and ECL operate as independent 8-bit event counters. Figure 9.21 shows an example of the software processing when ECH and ECL are used as 8-bit event counters.

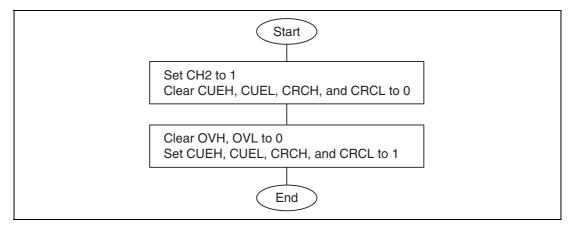


Figure 9.21 Example of Software Processing when Using ECH and ECL as 8-Bit Event Counters

ECH and ECL can be used as 8-bit event counters by carrying out the software processing shown in the example in figure 9.21. The 8-bit event counter operating clock source is asynchronous event input from the AEVH pin for ECH, and asynchronous event input from the AEVL pin for ECL. When the next clock is input after the ECH count value reaches H'FF, ECH overflows, the OVH flag is set to 1 in ECCSR, the ECH count value returns to H'00, and counting up is restarted. Similarly, when the next clock is input after the ECL count value reaches H'FF, ECL overflows, the OVL flag is set to 1 in ECCSR, the ECL count value returns to H'00, and counting up is restarted. When overflow occurs, the IRREC bit is set to 1 in IRR2. If the IENEC bit in IENR2 is 1 at this time, an interrupt request is sent to the CPU.

9.7.4 Asynchronous Event Counter Operation Modes

Asynchronous event counter operation modes are shown in table 9.21.

Table 9.21 Asynchronous Event Counter Operation Modes

Operation Mode	Reset	Active	Sleep	Watch	Sub- active	Sub- sleep	Standby	Module Standby
ECCSR	Reset	Functions	Functions	Held*	Functions	Functions	Held*	Held
ECH	Reset	Functions	Functions	Functions*	Functions	Functions	Functions*	Halted
ECL	Reset	Functions	Functions	Functions*	Functions	Functions	Functions*	Halted

Note: * When an asynchronous external event is input, the counter increments but the counter overflow H/L flags are not affected.

9.7.5 Application Notes

- 1. When reading the values in ECH and ECL, the correct value will not be returned if the event counter increments during the read operation. Therefore, if the counter is being used in the 8-bit mode, clear bits CUEH and CUEL in ECCSR to 0 before reading ECH or ECL. If the counter is being used in the 16-bit mode, clear CUEL only to 0 before reading ECH or ECL.
- The maximum clock frequency to be input to the AEVH and AEVL pins is 16 MHz. In addition, ensure that the high and low widths of the clock are at least 32 ns. The duty cycle is immaterial.

Mode			Maximum AEVH/AEVL Pin Input Clock Frequency
16-bit mode			$V_{cc} = 2.7 \text{ to } 5.5 \text{ V}/16 \text{ MHz}$
8-bit mode	Active (high-speed), sleep (high-speed)		_
8-bit mode	Active (medium-speed), sleep (medium-speed)	(φ/16)	2 · f _{osc}
		(¢/32)	f_{osc}
		(¢/64)	$1/2 \cdot f_{osc}$
	$f_{OSC} = 1 \text{ MHz to } 16 \text{ MHz}$	(¢/128)	$1/4 \cdot f_{OSC}$
8-bit mode	Watch, subactive, subsleep, standby	(\psi w/2)	1000 kHz
		$(\phi w/4)$	500 kHz
	$\phi_{W} = 32.768 \text{ kHz or } 38.4 \text{ kHz}$	$(\phi w/8)$	250 kHz

3. When AEC uses with 16-bit mode, set CUEH in ECCSR to "1" first, set CRCH in ECCSR to "1" second, or set both CUEH and CRCH to "1" at same time before clock entry. While AEC is operating on 16-bit mode, do not change CUEH. Otherwise, ECH will be miscounted up. Also, to reset the counter, clear CRCH and CRCL to 0 simultaneously or clear CRCL and CRCH to 0 sequentially, in that order.

Section 10 Serial Communication Interface

10.1 Overview

This LSI is provided with two serial communication interfaces, SCI3-1 and SCI3-2. These two SCIs have identical functions. In this manual, the generic term SCI3 is used to refer to both SCIs.

Serial communication interface 3 (SCI3) can carry out serial data communication in either asynchronous or synchronous mode.

10.1.1 Features

Features of SCI3 are listed below.

- Choice of asynchronous or synchronous mode for serial data communication
 - Asynchronous mode

Serial data communication is performed asynchronously, with synchronization provided character by character. In this mode, serial data can be exchanged with standard asynchronous communication LSIs such as a Universal Asynchronous Receiver/Transmitter (UART) or Asynchronous Communication Interface Adapter (ACIA).

There is a choice of 16 data transfer formats.

Data length	7, 8, 5 bits
Stop bit length	1 or 2 bits
Parity	Even, odd, or none
Receive error detection	Parity, overrun, and framing errors
Break detection	Break detected by reading the $\mbox{RXD}_{\mbox{\tiny 3x}}$ pin level directly when a framing error occurs

— Synchronous mode

Serial data communication is synchronized with a clock. In his mode, serial data can be exchanged with another LSI that has a synchronous communication function.

Data length	8 bits
Receive error detection	Overrun errors

Full-duplex communication

Separate transmission and reception units are provided, enabling transmission and reception to be carried out simultaneously. The transmission and reception units are both double-buffered, allowing continuous transmission and reception.

- On-chip baud rate generator, allowing any desired bit rate to be selected
- Choice of an internal or external clock as the transmit/receive clock source
- Six interrupt sources: transmit end, transmit data empty, receive data full, overrun error, framing error, and parity error

10.1.2 Block Diagram

Figure 10.1 shows a block diagram of SCI3.

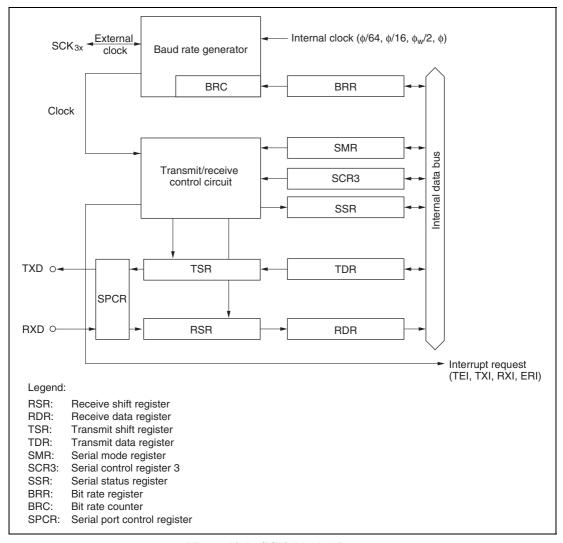


Figure 10.1 SCI3 Block Diagram

10.1.3 Pin Configuration

Table 10.1 shows the SCI3 pin configuration.

Table 10.1 Pin Configuration

Name	Abbr.	I/O	Function
SCI3 clock	SCK₃ _x	I/O	SCI3 clock input/output
SCI3 receive data input	$RXD_{\scriptscriptstyle 3x}$	Input	SCI3 receive data input
SCI3 transmit data output	TXD _{3x}	Output	SCI3 transmit data output

10.1.4 Register Configuration

Table 10.2 shows the SCI3 register configuration.

Table 10.2 Registers

Name	Abbr.	R/W	Initial Value	Address
Serial mode register	SMR	R/W	H'00	H'FFA8/FF98
Bit rate register	BRR	R/W	H'FF	H'FFA9/FF99
Serial control register 3	SCR3	R/W	H'00	H'FFAA/FF9A
Transmit data register	TDR	R/W	H'FF	H'FFAB/FF9B
Serial status register	SSR	R/W	H'84	H'FFAC/FF9C
Receive data register	RDR	R	H'00	H'FFAD/FF9D
Transmit shift register	TSR	Protected	_	_
Receive shift register	RSR	Protected	_	_
Bit rate counter	BRC	Protected	_	_
Clock stop register 1	CKSTPR1	R/W	H'FF	H'FFFA
Serial port control register	SPCR	R/W	H'C0	H'FF91

10.2 Register Descriptions

10.2.1 Receive Shift Register (RSR)

Bit	7	6	5	4	3	2	1	0
Read/Write	_	_	_	_	_	_	_	_

RSR is a register used to receive serial data. Serial data input to RSR from the RXD_{3x} pin is set in the order in which it is received, starting from the LSB (bit 0), and converted to parallel data. When one byte of data is received, it is transferred to RDR automatically.

RSR cannot be read or written directly by the CPU.

10.2.2 Receive Data Register (RDR)

Bit	7	6	5	4	3	2	1	0
	RDR7	RDR6	RDR5	RDR4	RDR3	RDR2	RDR1	RDR0
Initial value	0	0	0	0	0	0	0	0
Read/Write	R	R	R	R	R	R	R	R

RDR is an 8-bit register that stores received serial data.

When reception of one byte of data is finished, the received data is transferred from RSR to RDR, and the receive operation is completed. RSR is then able to receive data. RSR and RDR are double-buffered, allowing consecutive receive operations.

RDR is a read-only register, and cannot be written by the CPU.

RDR is initialized to H'00 upon reset, and in standby, module standby or watch mode.

10.2.3 Transmit Shift Register (TSR)

Bit	7	6	5	4	3	2	1	0	
Read/Write		_	_	_	_	_	_	_	

TSR is a register used to transmit serial data. Transmit data is first transferred from TDR to TSR, and serial data transmission is carried out by sending the data to the TXD_{3x} pin in order, starting from the LSB (bit 0). When one byte of data is transmitted, the next byte of transmit data is transferred to TDR, and transmission started, automatically. Data transfer from TDR to TSR is not performed if no data has been written to TDR (if bit TDRE is set to 1 in the serial status register (SSR)).

TSR cannot be read or written directly by the CPU.

10.2.4 Transmit Data Register (TDR)

Bit	7	6	5	4	3	2	1	0
	TDR7	TDR6	TDR5	TDR4	TDR3	TDR2	TDR1	TDR0
Initial value	1	1	1	1	1	1	1	1
Read/Write	R/W							

TDR is an 8-bit register that stores transmit data. When TSR is found to be empty, the transmit data written in TDR is transferred to TSR, and serial data transmission is started. Continuous transmission is possible by writing the next transmit data to TDR during TSR serial data transmission.

TDR can be read or written by the CPU at any time.

TDR is initialized to H'FF upon reset, and in standby, module standby, or watch mode.

10.2.5 Serial Mode Register (SMR)

Bit	7	6	5	4	3	2	1	0
	СОМ	CHR	PE	PM	STOP	MP	CKS1	CKS0
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W

SMR is an 8-bit register used to set the serial data transfer format and to select the clock source for the baud rate generator.

SMR can be read or written by the CPU at any time.

SMR is initialized to H'00 upon reset, and in standby, module standby, or watch mode.

Bit 7: Communication mode (COM)

Bit 7 selects whether SCI3 operates in asynchronous mode or synchronous mode.

Bit 7		
COM	Description	
0	Asynchronous mode	(initial value)
1	Synchronous mode	

Bit 6: Character length (CHR)

Bit 6 selects either 7 or 8 bits as the data length to be used in asynchronous mode. In synchronous mode the data length is always 8 bits, irrespective of the bit 6 setting.

Bit 6		
CHR	Description	
0	8-bit data/5-bit data*2	(initial value)
1	7-bit data*1/5-bit data*2	

Notes: 1. When 7-bit data is selected, the MSB (bit 7) of TDR is not transmitted.

2. When 5-bit data is selected, set both PE and MP to 1. The three most significant bits (bits 7, 6, and 5) of TDR are not transmitted.

Bit 5: Parity enable (PE)

Bit 5 selects whether a parity bit is to be added during transmission and checked during reception in asynchronous mode. In synchronous mode parity bit addition and checking is not performed, irrespective of the bit 5 setting.

Bit 5 PE	Description	
0	Parity bit addition and checking disabled*2	(initial value)
1	Parity bit addition and checking enabled*1*2	

- Notes: 1. When PE is set to 1, even or odd parity, as designated by bit PM, is added to transmit data before it is sent, and the received parity bit is checked against the parity designated by bit PM.
 - 2. For the case where 5-bit data is selected, see table 10.11.

Bit 4: Parity mode (PM)

Bit 4 selects whether even or odd parity is to be used for parity addition and checking. The PM bit setting is only valid in asynchronous mode when bit PE is set to 1, enabling parity bit addition and checking. The PM bit setting is invalid in synchronous mode, and in asynchronous mode if parity bit addition and checking is disabled.

Bit 4		
PM	Description	
0	Even parity*1	(initial value)
1	Odd parity*2	

- Notes: 1. When even parity is selected, a parity bit is added in transmission so that the total number of 1 bits in the transmit data plus the parity bit is an even number; in reception, a check is carried out to confirm that the number of 1 bits in the receive data plus the parity bit is an even number.
 - 2. When odd parity is selected, a parity bit is added in transmission so that the total number of 1 bits in the transmit data plus the parity bit is an odd number; in reception, a check is carried out to confirm that the number of 1 bits in the receive data plus the parity bit is an odd number.

Bit 3: Stop bit length (STOP)

Bit 3 selects 1 bit or 2 bits as the stop bit length is asynchronous mode. The STOP bit setting is only valid in asynchronous mode. When synchronous mode is selected the STOP bit setting is invalid since stop bits are not added.

Bit 3		
STOP	Description	
0	1 stop bit*1	(initial value)
1	2 stop bits*2	

Notes: 1. In transmission, a single 1 bit (stop bit) is added at the end of a transmit character.

2. In transmission, two 1 bits (stop bits) are added at the end of a transmit character.

In reception, only the first of the received stop bits is checked, irrespective of the STOP bit setting. If the second stop bit is 1 it is treated as a stop bit, but if 0, it is treated as the start bit of the next transmit character.

Bit 2: 5-Bit Communication (MP)

Bit 2 can set the 5-bit communication format. When this bit is set to 1, the 5-bit communication format is enabled. When writing 1 to this bit, always write 1 to bit 5 (RE) at the same time.

Bit 2 MP	Description			
0	7/8-bit communication format*	(initial value)		
1	5-bit communication format*			
Notes	For the case where F hit data is calcuted assetable 10.11			

Note: * For the case where 5-bit data is selected, see table 10.11.

Bits 1 and 0: Clock select 1, 0 (CKS1, CKS0)

Bits 1 and 0 choose $\phi/64$, $\phi/16$, $\phi w/2$, or ϕ as the clock source for the baud rate generator.

For the relation between the clock source, bit rate register setting, and baud rate, see section 10.2.8, Bit rate register (BRR).

Bit 1 CKS1	Bit 0 CKS0	Description	
0	0	φ clock	(initial value)
0	1	φ w/2 clock*1/φ w clock*2	
1	0	φ/16 clock	
1	1	φ/64 clock	

Notes: 1. ϕ w/2 clock in active (medium-speed/high-speed) mode and sleep mode

2. φ w clock in subactive mode and subsleep mode
In subactive or subsleep mode, SCI3 can be operated when CPU clock is φw/2 only.

10.2.6 Serial Control Register 3 (SCR3)

Bit	7	6	5	4	3	2	1	0
	TIE	RIE	TE	RE	MPIE	TEIE	CKE1	CKE0
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W

SCR3 is an 8-bit register for selecting transmit or receive operation, the asynchronous mode clock output, interrupt request enabling or disabling, and the transmit/receive clock source.

SCR3 can be read or written by the CPU at any time.

SCR3 is initialized to H'00 upon reset, and in standby, module standby or watch mode.

Bit 7: Transmit interrupt enable (TIE)

Bit 7 selects enabling or disabling of the transmit data empty interrupt request (TXI) when transmit data is transferred from the transmit data register (TDR) to the transmit shift register (TSR), and bit TDRE in the serial status register (SSR) is set to 1.

TXI can be released by clearing bit TDRE or bit TIE to 0.

Bit 7 TIE	Description	
0	Transmit data empty interrupt request (TXI) disabled	(initial value)
1	Transmit data empty interrupt request (TXI) enabled	

Bit 6: Receive interrupt enable (RIE)

Bit 6 selects enabling or disabling of the receive data full interrupt request (RXI) and the receive error interrupt request (ERI) when receive data is transferred from the receive shift register (RSR) to the receive data register (RDR), and bit RDRF in the serial status register (SSR) is set to 1. There are three kinds of receive error: overrun, framing, and parity.

RXI can be released by clearing bit RDRF or the FER, PER, or OER error flag to 0, or by clearing bit RIE to 0.

Bit 6 RIE	Description	
0	Receive data full interrupt request (RXI) and receive error interrupt request (ERI) disabled	(initial value)
1	Receive data full interrupt request (RXI) and receive error interrupt request (ERI) enabled	

Bit 5: Transmit enable (TE)

Bit 5 selects enabling or disabling of the start of transmit operation.

Bit 5 TE	Description	
0	Transmit operation disabled*1 (TXD pin is I/O port)	(initial value)
1	Transmit operation enabled*2 (TXD pin is transmit data pin)	

Notes: 1. Bit TDRE in SSR is fixed at 1.

 When transmit data is written to TDR in this state, bit TDR in SSR is cleared to 0 and serial data transmission is started. Be sure to carry out serial mode register (SMR) settings, and setting of bit SPC31 or SPC32 in SPCR, to decide the transmission format before setting bit TE to 1.

Bit 4: Receive enable (RE)

Bit 4 selects enabling or disabling of the start of receive operation.

Bit 4 RE	Description	
0	Receive operation disabled*1 (RXD pin is I/O port)	(initial value)
1	Receive operation enabled*2 (RXD pin is receive data pin)	

Notes: 1. Note that the RDRF, FER, PER, and OER flags in SSR are not affected when bit RE is cleared to 0, and retain their previous state.

In this state, serial data reception is started when a start bit is detected in asynchronous mode or serial clock input is detected in synchronous mode. Be sure to carry out serial mode register (SMR) settings to decide the reception format before setting bit RE to 1.

Bit 3: Reserved (MPIE)

Bit 3 is reserved.



Bit 2: Transmit end interrupt enable (TEIE)

Bit 2 selects enabling or disabling of the transmit end interrupt request (TEI) if there is no valid transmit data in TDR when MSB data is to be sent.

Bit 2 TEIE	Description	
0	Transmit end interrupt request (TEI) disabled	(initial value)
1	Transmit end interrupt request (TEI) enabled*	

Note: * TEI can be released by clearing bit TDRE to 0 and clearing bit TEND to 0 in SSR, or by clearing bit TEIE to 0.

Bits 1 and 0: Clock enable 1 and 0 (CKE1, CKE0)

Bits 1 and 0 select the clock source and enabling or disabling of clock output from the SCK_{3x} pin. The combination of CKE1 and CKE0 determines whether the SCK_{3x} pin functions as an I/O port, a clock output pin, or a clock input pin.

The CKE0 bit setting is only valid in case of internal clock operation (CKE1 = 0) in asynchronous mode. In synchronous mode, or when external clock operation is used (CKE1 = 1), bit CKE0 should be cleared to 0.

After setting bits CKE1 and CKE0, set the operating mode in the serial mode register (SMR).

For details on clock source selection, see table 10.9 in 10.3.1, Overview.

Bit 1	Bit 0	Description					
CKE1	CKE0	Communication Mode	Clock Source	SCK₃x Pin Function			
0	0	Asynchronous	Internal clock	I/O port*1			
		Synchronous	Internal clock	Serial clock output*1			
0	1	Asynchronous	Internal clock	Clock output*2			
		Synchronous	Reserved				
1	0	Asynchronous	External clock	Clock input*3			
		Synchronous	External clock	Serial clock input			
1	1	Asynchronous	Reserved				
		Synchronous	Reserved				

Notes: 1. Initial value

- 2. A clock with the same frequency as the bit rate is output.
- 3. Input a clock with a frequency 16 times the bit rate.

10.2.7 Serial Status Register (SSR)

Bit	7	6	5	4	3	2	1	0
	TDRE	RDRF	OER	FER	PER	TEND	MPBR	MPBT
Initial value	1	0	0	0	0	1	0	0
Read/Write	R/(W)*	R/(W)*	R/(W)*	R/(W)*	R/(W)*	R	R	R/W

Note: * Only a write of 0 for flag clearing is possible.

SSR is an 8-bit register containing status flags that indicate the operational status of SCI3.

SSR can be read or written by the CPU at any time, but only a write of 1 is possible to bits TDRE, RDRF, OER, PER, and FER. In order to clear these bits by writing 0, 1 must first be read.

Bits TEND and MPBR are read-only bits, and cannot be modified.

SSR is initialized to H'84 upon reset, and in standby, module standby, or watch mode.



Bit 7: Transmit data register empty (TDRE)

Bit 7 indicates that transmit data has been transferred from TDR to TSR.

Bit 7 TDRE	Description	
0	Transmit data written in TDR has not been transferred to TSR	
	Clearing conditions:	
	After reading TDRE = 1, cleared by writing 0 to TDRE	
	When data is written to TDR by an instruction	
1	Transmit data has not been written to TDR, or transmit data written in TDR has b transferred to TSR	een
	Setting conditions:	
	When bit TE in SCR3 is cleared to 0	
	When data is transferred from TDR to TSR (initial va	ılue)

Bit 6: Receive data register full (RDRF)

Bit 6 indicates that received data is stored in RDR.

Bit 6 RDRF	Description	
0	•	<u>~</u>
U	(manual salah	e)
	Clearing conditions:	
	After reading RDRF = 1, cleared by writing 0 to RDRF	
	When RDR data is read by an instruction	
1	There is receive data in RDR	
	Setting condition:	
	When reception ends normally and receive data is transferred from RSR to RDR	

Note: If an error is detected in the receive data, or if the RE bit in SCR3 has been cleared to 0, RDR and bit RDRF are not affected and retain their previous state.

Note that if data reception is completed while bit RDRF is still set to 1, an overrun error (OER) will result and the receive data will be lost.

Bit 5: Overrun error (OER)

Bit 5 indicates that an overrun error has occurred during reception.

OER	Description
0	Reception in progress or completed*1 (initial value)
	Clearing condition:
	After reading OER = 1, cleared by writing 0 to OER
1	An overrun error has occurred during reception*2
	Setting condition:
	When reception is completed with RDRF set to 1
NI-t d	When his DC is CODO is also and to O his OCD is not effected and action its annuity

- Notes: 1. When bit RE in SCR3 is cleared to 0, bit OER is not affected and retains its previous state.
 - RDR retains the receive data it held before the overrun error occurred, and data received after the error is lost. Reception cannot be continued with bit OER set to 1, and in synchronous mode, transmission cannot be continued either.

Bit 4: Framing error (FER)

Bit 4 indicates that a framing error has occurred during reception in asynchronous mode.

Bit 4 FER	Description
0	Reception in progress or completed*1 (initial value)
	Clearing condition:
	After reading FER = 1, cleared by writing 0 to FER
1	A framing error has occurred during reception
	Setting condition:
	When the stop bit at the end of the receive data is checked for a value of 1 at the end of reception, and the stop bit is 0^{*2}

- Notes: 1. When bit RE in SCR3 is cleared to 0, bit FER is not affected and retains its previous state.
 - 2. Note that, in 2-stop-bit mode, only the first stop bit is checked for a value of 1, and the second stop bit is not checked. When a framing error occurs the receive data is transferred to RDR but bit RDRF is not set. Reception cannot be continued with bit FER set to 1. In synchronous mode, neither transmission nor reception is possible when bit FER is set to 1.

Bit 3: Parity error (PER)

Bit 3 indicates that a parity error has occurred during reception with parity added in asynchronous mode.

Bit 3		
PER	Description	
0	Reception in progress or completed*1	(initial value)
	Clearing condition:	
	After reading PER = 1, cleared by writing 0 to PER	
1	A parity error has occurred during reception*2	_
	Setting condition:	
	When the number of 1 bits in the receive data plus parity bit does not m designated by bit PM in the serial mode register (SMR)	natch the parity

Notes: 1. When bit RE in SCR3 is cleared to 0, bit PER is not affected and retains its previous state.

 Receive data in which it a parity error has occurred is still transferred to RDR, but bit RDRF is not set. Reception cannot be continued with bit PER set to 1. In synchronous mode, neither transmission nor reception is possible when bit FER is set to 1.

Bit 2: Transmit end (TEND)

Bit 2 indicates that bit TDRE is set to 1 when the last bit of a transmit character is sent.

Bit 2 is a read-only bit and cannot be modified.

Bit 2 TEND	Description	
0	Transmission in progress	
	Clearing conditions:	
	After reading TDRE = 1, cleared by writing 0 to TDRE	
	When data is written to TDR by an instruction	
1	Transmission ended	(initial value)
	Setting conditions:	
	When bit TE in SCR3 is cleared to 0	
	When bit TDRE is set to 1 when the last bit of a transmit character is s	ent

Bit 1: Reserved (MPBR)

Bit 1 is reserved.

It is a read-only bit and cannot be modified.

Bit 0: Reserved (MPBT)

Bit 0 is reserved.

The write value should always be 0.

10.2.8 Bit Rate Register (BRR)

Bit	7	6	5 4		3 2		1	0	
	BRR7	BRR6	BRR5	BRR4	BRR3	BRR2	BRR1	BRR0	
Initial value	1	1	1	1	1	1	1	1	
Read/Write	R/W								

BRR is an 8-bit register that designates the transmit/receive bit rate in accordance with the baud rate generator operating clock selected by bits CKS1 and CKS0 of the serial mode register (SMR).

BRR can be read or written by the CPU at any time.

BRR is initialized to H'FF upon reset, and in standby, module standby, or watch mode.

Table 10.3 shows examples of BRR settings in asynchronous mode. The values shown are for active (high-speed) mode.

Table 10.3 Examples of BRR Settings for Various Bit Rates (Asynchronous Mode) (1)

osc

	32.8 kHz			38.4 kHz			2 MHz			2.	4576 I	MHz	4 MHz		
Bit Rate			Error			Error			Error			Error			Error
(bit/s)	n	N	(%)	n	N	(%)	n	N	(%)	n	N	(%)	n	N	(%)
110			e used,	_	_	_	_	_	_	2	21	-0.83	_	_	_
150	as error exceeds 3%			0	3	0	2	12	0.16	3	3	0	2	25	0.16
200				0	2	0	0	155	0.16	3	2	0	_	_	_
250	_			_	_	_	0	124	0	0	153	-0.26	0	249	0
300	_			0	1	0	0	103	0.16	3	1	0	2	12	0.16
600	_			0	0	0	0	51	0.16	3	0	0	0	103	0.16
1200	_			_	_	_	0	25	0.16	2	1	0	0	51	0.16
2400	_			_	_	_	0	12	0.16	2	0	0	0	25	0.16
4800	_			_	_	_	_	_	_	0	7	0	0	12	0.16
9600	_			_	_	_	_	_	_	0	3	0	_	_	_
19200				_	_		_		_	0	1	0		_	_
31250				_	_		0	0	0	_	_	_	0	1	0
38400				_	_		_		_	0	0	0		_	_

Table 10.3 Examples of BRR Settings for Various Bit Rates (Asynchronous Mode) (2)

			0	SC			
		10 MI	Ηz		16 M	6 MHz	
Bit Rate			Error			Error	
(bit/s)	n	N	(%)	n	N	(%)	
110	2	88	-0.25	2	141	0.03	
150	2	64	0.16	2	103	0.16	
200	2	48	-0.35	2	77	0.16	
250	2	38	0.16	2	62	-0.79	
300	_	_	_	2	51	0.16	
600	_	_	_	2	25	0.16	
1200	0	129	0.16	0	207	0.16	
2400	0	64	0.16	0	103	0.16	
4800	_	_	_	0	51	0.16	
9600	_		_	0	25	0.16	
19200	_	_	_	0	12	0.16	
31250	0	4	0	0	7	0	
38400	_	_	_	_	_	_	

Notes: 1. The setting should be made so that the error is not more than 1%.

2. The value set in BRR is given by the following equation:

$$N = \frac{OSC}{(64 \times 2^{2n} \times B)} - 1$$

where B: Bit rate (bit/s)

N: Baud rate generator BRR setting $(0 \le N \le 255)$

OSC: Value of ϕ_{OSC} (Hz)

n: Baud rate generator input clock number (n = 0, 2, or 3)(The relation between n and the clock is shown in table 10.4.)

3. The error in table 10.3 is the value obtained from the following equation, rounded to two decimal places.

Error (%) =
$$\frac{\text{B (rate obtained from n, N, OSC)} - \text{R(bit rate in left-hand column in table 10.3.)}}{\text{R (bit rate in left-hand column in table 10.3.)}} \times 100$$

Table 10.4 Relation between n and Clock

SMR Se	attina

			_	
n	Clock	CKS1	CKS0	
0	ф	0	0	
0	$\phi_{w}/2^{*1}/\phi_{w}^{*2}$	0	1	
2	φ/16	1	0	
3	ф/64	1	1	

Notes: 1. ϕ w/2 clock in active (medium-speed/high-speed) mode and sleep mode

2. φ w clock in subactive mode and subsleep mode
In subactive or subsleep mode, SCI3 can be operated when CPU clock is φw/2 only.

Table 10.5 shows the maximum bit rate for each frequency. The values shown are for active (high-speed) mode.

Table 10.5 Maximum Bit Rate for Each Frequency (Asynchronous Mode)

			Setting
OSC (MHz)	Maximum Bit Rate (bit/s)	n	N
0.0384*	600	0	0
2	31250	0	0
2.4576	38400	0	0
4	62500	0	0
10	156250	0	0
16	250000	0	0

^{*:} When SMR is set up to CKS1 = "0", CKS0 = "1".

Table 10.6 shows examples of BRR settings in synchronous mode. The values shown are for active (high-speed) mode.



Table 10.6 Examples of BRR Settings for Various Bit Rates (Synchronous Mode)

osc

Bit Rate	;	38.4 k	Hz		2 MH	z		4 MH	lz		10 MH	Ηz		16 M	Hz
(bit/s)	n	N	Error	n	N	Error	n	N	Error	n	N	Error	n	N	Error
200	0	23	0	_	_	_	_	_	_	_	_	_	_	_	_
250			_	_	_	_	2	124	0	_	_	_	3	124	0
300	2	0	0	_	_	_	_	_		_	_	_	_	_	_
500				_	_	_	_		_	_	_	_	2	249	0
1K				0	249	0	_	_	_	_	_	_	2	124	0
2.5K				0	99	0	0	199	0	_	_	_	2	49	0
5K				0	49	0	0	99	0	0	249	0	2	24	0
10K				0	24	0	0	49	0	0	124	0	0	199	0
25K				0	9	0	0	19	0	0	49	0	0	79	0
50K				0	4	0	0	9	0	0	24	0	0	39	0
100K				_	_	_	0	4	0	_	_	_	0	19	0
250K				0	0	0	0	1	0	0	4	0	0	7	0
500K							0	0	0	_	_	_	0	3	0
1M										_	_	_	0	1	0

Blank: Cannot be set.

—: A setting can be made, but an error will result.

Notes: The value set in BRR is given by the following equation:

$$N = \frac{OSC}{(8 \times 2^{2n} \times B)} - 1$$

where B: Bit rate (bit/s)

N: Baud rate generator BRR setting $(0 \le N \le 255)$

OSC: Value of ϕ_{OSC} (Hz)

n: Baud rate generator input clock number (n = 0, 2, or 3)
(The relation between n and the clock is shown in table 10.7.)

Table 10.7 Relation between n and Clock

		SMR Setting				
n	Clock	CKS1	CKS0			
0	ф	0	0			
0	$\phi_{w}/2^{*1}/\phi_{w}^{*2}$	0	1			
2	ф/16	1	0			
3	φ/64	1	1			

Notes: 1. φw/2 clock in active (medium-speed/high-speed) mode and sleep mode

2. φw clock in subactive mode and subsleep mode
In subactive or subsleep mode, SCI3 can be operated when CPU clock is φw/2 only.

10.2.9 Clock Stop Register 1 (CKSTPR1)

Bit	7	6	5	4	3	2	1	0
	_	S31CKSTP	S32CKSTP	ADCKSTP	TGCKSTP	TFCKSTP	TCCKSTP	TACKSTP
Initial value	1	1	1	1	1	1	1	1
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W

CKSTPR1 is an 8-bit read/write register that performs module standby mode control for peripheral modules. Only the bits relating to SCI3 are described here. For details of the other bits, see the sections on the relevant modules.

Bit 6: SCI31 module standby mode control (S31CKSTP)

Bit 6 controls setting and clearing of module standby mode for SCI31.

S31CKSTP Description

0	SCI31 is set to module standby mode	
1	SCI31 module standby mode is cleared	(initial value)

Note: All SCI31 register is initialized in module standby mode.



Bit 5: SCI32 module standby mode control (S32CKSTP)

Bit 5 controls setting and clearing of module standby mode for SCI32.

S32CKSTP Description

0	SCI32 is set to module standby mode	
1	SCI32 module standby mode is cleared	(initial value)

Note: All SCI32 register is initialized in module standby mode.

10.2.10 Serial Port Control Register (SPCR)

Bit	7	6	5	4	3	2	1	0
	_	_	SPC32	SPC31	SCINV3	SCINV2	SCINV1	SCINV0
Initial value	1	1	0	0	0	0	0	0
Read/Write	_	_	R/W	R/W	R/W	R/W	R/W	R/W

SPCR is an 8-bit readable/writable register that performs RXD₃₁, RXD₃₂, TXD₃₁, and TXD₃₂ pin input/output data inversion switching. SPCR is initialized to H'C0 by a reset.

Bits 7 and 6: Reserved bits

Bits 7 and 6 are reserved; they are always read as 1 and cannot be modified.

Bit 5: P4₂/TXD₃₂ pin function switch (SPC32)

This bit selects whether pin P4₂/TXD₃₂ is used as P4₂ or as TXD₃₂.

Bit 5 SPC32	Description	
0	Functions as P4 ₂ I/O pin	(initial value)
1	Functions as TXD ₃₂ output pin*	

Note: $\,\,^*\,\,$ Set the TE bit in SCR3 after setting this bit to 1.

Bit 4: P3₅/TXD₃₁ pin function switch (SPC31)

This bit selects whether pin P3₅/TXD₃₁ is used as P3₅ or as TXD₃₁.

Bit 4 SPC31	Description	
0	Functions as P3 ₅ I/O pin	(initial value)
1	Functions as TXD ₃₁ output pin*	
Note: *	Set the TE bit in SCR3 after setting this bit to 1.	

Bit 3: TXD₃₂ pin output data inversion switch

Bit 3 specifies whether or not TXD₃₂ pin output data is to be inverted.

Bit 3 SCINV3	Description	
0	TXD ₃₂ output data is not inverted	(initial value)
1	TXD ₃₂ output data is inverted	

Bit 2: RXD₃₂ pin input data inversion switch

Bit 2 specifies whether or not RXD_{32} pin input data is to be inverted.

Bit 2 SCINV2	Description	
0	RXD ₃₂ input data is not inverted	(initial value)
1	RXD ₃₂ input data is inverted	

Bit 1: TXD₃₁ pin output data inversion switch

Bit 1 specifies whether or not TXD₃₁ pin output data is to be inverted.

Bit 1 SCINV1	Description	
0	TXD ₃₁ output data is not inverted	(initial value)
1	TXD ₃₁ output data is inverted	

Bit 0: RXD₃₁ pin input data inversion switch

Bit 0 specifies whether or not $RXD_{_{\rm 31}}$ pin input data is to be inverted.

Bit 0 SCINV0	Description	
0	RXD ₃₁ input data is not inverted	(initial value)
1	RXD ₃₁ input data is inverted	

10.3 Operation

10.3.1 Overview

SCI3 can perform serial communication in two modes: asynchronous mode in which synchronization is provided character by character, and synchronous mode in which synchronization is provided by clock pulses. The serial mode register (SMR) is used to select asynchronous or synchronous mode and the data transfer format, as shown in table 10.8.

The clock source for SCI3 is determined by bit COM in SMR and bits CKE1 and CKE0 in SCR3, as shown in table 10.9.

(1) Synchronous Mode

- Choice of 5-, 7-, or 8-bit data length
- Choice of parity addition, and addition of 1 or 2 stop bits. (The combination of these parameters determines the data transfer format and the character length.)
- Framing error (FER), parity error (PER), overrun error (OER), and break detection during reception
- Choice of internal or external clock as the clock source
 - When internal clock is selected: SCI3 operates on the baud rate generator clock, and a clock with the same frequency as the bit rate can be output.
 - When external clock is selected: A clock with a frequency 16 times the bit rate must be input. (The on-chip baud rate generator is not used.)

(2) Synchronous Mode

- Data transfer format: Fixed 8-bit data length
- Overrun error (OER) detection during reception
- Choice of internal or external clock as the clock source
 - When internal clock is selected: SCI3 operates on the baud rate generator clock, and a serial clock is output.
 - When external clock is selected: The on-chip baud rate generator is not used, and SCI3 operates on the input serial clock.



Table 10.8 SMR Settings and Corresponding Data Transfer Formats

		SMR				Dat	ta Transfer Format			
Bit 7 COM	Bit 6 CHR	Bit 2 MP	Bit 5 PE	Bit 3 STOP	Mode	Data Length	Parity Bit	Stop Bit Length		
0	0	0	0	0	Asynchronous	8-bit data	No	1 bit		
				1	mode			2 bits		
			1	0	_		Yes	1 bit		
				1	_			2 bits		
	1	_	0	0	_	7-bit data	No	1 bit		
				1	_			2 bits		
			1	0	_		Yes	1 bit		
				1	_			2 bits		
	0	1	1	1 (0	0	Setting			
				1	prohibited					
			1	0	Asynchronous	5-bit data	No	1 bit		
				1	mode			2 bits		
	1		0	0	Setting					
				1	prohibited					
			1	0	Asynchronous	5-bit data	Yes	1 bit		
				1	mode			2 bits		
1	*	0	*	*	Synchronous mode	8-bit data	No	No		

*: Don't care

Table 10.9 SMR and SCR3 Settings and Clock Source Selection

SMR	SC	CR3					
Bit 7	Bit 1	Bit 0	•	Transmit/Receive Clock			
СОМ	CKE1	CKE0	Mode	Clock Source	SCK _{3x} Pin Function		
0	0	0	Asynchronous	Internal	I/O port (SCK _{3x} pin not used)		
		1	mode		Outputs clock with same frequency as bit rate		
	1	0	•	External	Outputs clock with frequency 16 times bit rate		
1	0	0	Synchronous	Internal	Outputs serial clock		
	1	0	mode	External	Inputs serial clock		
0	1	1	Reserved (Do not specify these combinations)				
1	0	1	-				
1	1	1	•				

(3) Interrupts and Continuous Transmission/Reception

SCI3 can carry out continuous reception using RXI and continuous transmission using TXI. These interrupts are shown in table 10.10.

Table 10.10 Transmit/Receive Interrupts

Interrupt	Flags	Interrupt Request Conditions	Notes
RXI	RDRF RIE	When serial reception is performed normally and receive data is transferred from RSR to RDR, bit RDRF is set to 1, and if bit RIE is set to 1 at this time, RXI is enabled and an interrupt is requested. (See figure 10.2(a).)	The RXI interrupt routine reads the receive data transferred to RDR and clears bit RDRF to 0. Continuous reception can be performed by repeating the above operations until reception of the next RSR data is completed.
TXI	TDRE TIE	When TSR is found to be empty (on completion of the previous transmission) and the transmit data placed in TDR is transferred to TSR, bit TDRE is set to 1. If bit TIE is set to 1 at this time, TXI is enabled and an interrupt is requested. (See figure 10.2(b).)	The TXI interrupt routine writes the next transmit data to TDR and clears bit TDRE to 0. Continuous transmission can be performed by repeating the above operations until the data transferred to TSR has been transmitted.
TEI	TEND TEIE	When the last bit of the character in TSR is transmitted, if bit TDRE is set to 1, bit TEND is set to 1. If bit TEIE is set to 1 at this time, TEI is enabled and an interrupt is requested. (See figure 10.2(c).)	TEI indicates that the next transmit data has not been written to TDR when the last bit of the transmit character in TSR is sent.

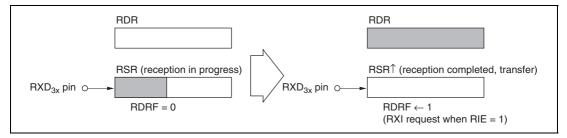


Figure 10.2 (a) RDRF Setting and RXI Interrupt

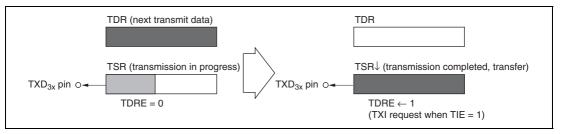


Figure 10.2 (b) TDRE Setting and TXI Interrupt

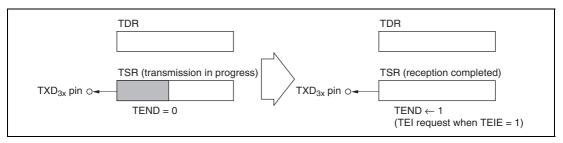


Figure 10.2 (c) TEND Setting and TEI Interrupt

10.3.2 Operation in Asynchronous Mode

In asynchronous mode, serial communication is performed with synchronization provided character by character. A start bit indicating the start of communication and one or two stop bits indicating the end of communication are added to each character before it is sent.

SCI3 has separate transmission and reception units, allowing full-duplex communication. As the transmission and reception units are both double-buffered, data can be written during transmission and read during reception, making possible continuous transmission and reception.

(1) Data Transfer Format

The general data transfer format in asynchronous communication is shown in figure 10.3.

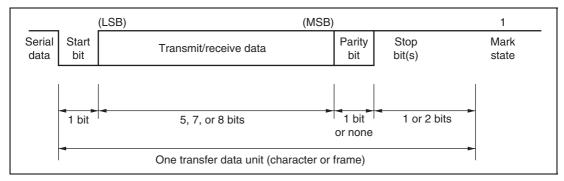


Figure 10.3 Data Format in Asynchronous Communication

In asynchronous communication, the communication line is normally in the mark state (high level). SCI3 monitors the communication line and when it detects a space (low level), identifies this as a start bit and begins serial data communication.

One transfer data character consists of a start bit (low level), followed by transmit/receive data (LSB-first format, starting from the least significant bit), a parity bit (high or low level), and finally one or two stop bits (high level).

In asynchronous mode, synchronization is performed by the falling edge of the start bit during reception. The data is sampled on the 8th pulse of a clock with a frequency 16 times the bit period, so that the transfer data is latched at the center of each bit.

Table 10.11 shows the 16 data transfer formats that can be set in asynchronous mode. The format is selected by the settings in the serial mode register (SMR).

Table 10.11 Data Transfer Formats (Asynchronous Mode)

SMR				Serial Data Transfer Format and Frame Length
CHR	PE	MP	STOP	1 2 3 4 5 6 7 8 9 10 11 12
0	0	0	0	S 8-bit data STOP
0	0	0	1	S 8-bit data STOP STOP
0	0	1	0	Setting prohibited
0	0	1	1	Setting prohibited
0	1	0	0	S 8-bit data P STOP
0	1	0	1	S 8-bit data P STOP STOP
0	1	1	0	S 5-bit data STOP
0	1	1	1	S 5-bit data STOP STOP
1	0	0	0	S 7-bit data STOP
1	0	0	1	S 7-bit data STOP STOP
1	0	1	0	Setting prohibited
1	0	1	1	Setting prohibited
1	1	0	0	S 7-bit data P STOP
1	1	0	1	S 7-bit data P STOP STOP
1	1	1	0	S 5-bit data P STOP
1	1	1	1	S 5-bit data P STOP STOP

Legend:

S: Start bit STOP: Stop bit P: Parity bit



(2) Clock

Either an internal clock generated by the baud rate generator or an external clock input at the SCK_{3x} pin can be selected as the SCI3 transmit/receive clock. The selection is made by means of bit COM in SMR and bits SCE1 and CKE0 in SCR3. See table 10.9 for details on clock source selection.

When an external clock is input at the SCK_{3x} pin, the clock frequency should be 16 times the bit rate.

When SCI3 operates on an internal clock, the clock can be output at the SCK_{3x} pin. In this case the frequency of the output clock is the same as the bit rate, and the phase is such that the clock rises at the center of each bit of transmit/receive data, as shown in figure 10.4.

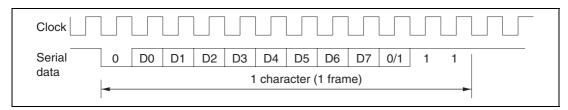


Figure 10.4 Phase Relationship between Output Clock and Transfer Data (Asynchronous Mode) (8-bit data, parity, 2 stop bits)

(3) Data Transfer Operations

SCI3 initialization: Before data is transferred on SCI3, bits TE and RE in SCR3 must first be cleared to 0, and then SCI3 must be initialized as follows.

Note: If the operation mode or data transfer format is changed, bits TE and RE must first be cleared to 0.

When bit TE is cleared to 0, bit TDRE is set to 1.

Note that the RDRF, PER, FER, and OER flags and the contents of RDR are retained when RE is cleared to 0.

When an external clock is used in asynchronous mode, the clock should not be stopped during operation, including initialization. When an external clock is used in synchronous mode, the clock should not be supplied during operation, including initialization.

Figure 10.5 shows an example of a flowchart for initializing SCI3.

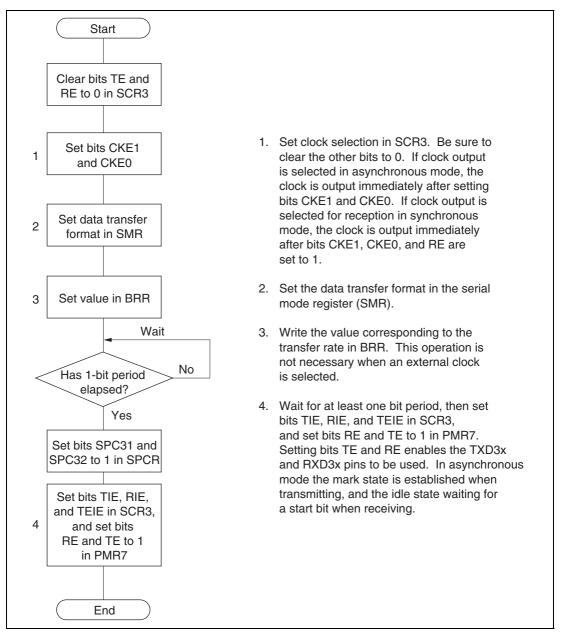


Figure 10.5 Example of SCI3 Initialization Flowchart

Transmitting: Figure 10.6 shows an example of a flowchart for data transmission. This procedure should be followed for data transmission after initializing SCI3.

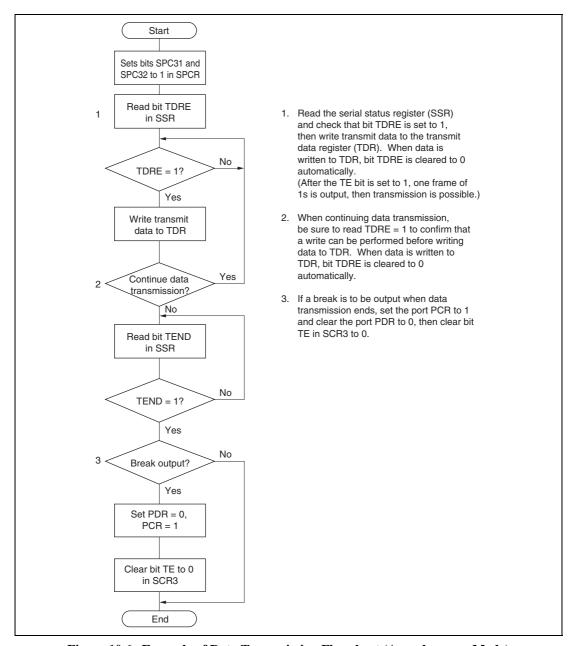


Figure 10.6 Example of Data Transmission Flowchart (Asynchronous Mode)

SCI3 operates as follows when transmitting data.

SCI3 monitors bit TDRE in SSR, and when it is cleared to 0, recognizes that data has been written to TDR and transfers data from TDR to TSR. It then sets bit TDRE to 1 and starts transmitting. If bit TIE in SCR3 is set to 1 at this time, a TXI request is made.

Serial data is transmitted from the TXD3x pin using the relevant data transfer format in table 10.11. When the stop bit is sent, SCI3 checks bit TDRE. If bit TDRE is cleared to 0, SCI3 transfers data from TDR to TSR, and when the stop bit has been sent, starts transmission of the next frame. If bit TDRE is set to 1, bit TEND in SSR bit is set to 1 the mark state, in which 1s are transmitted, is established after the stop bit has been sent. If bit TEIE in SCR3 is set to 1 at this time, a TEI request is made.

Figure 10.12 shows an example of the operation when transmitting in asynchronous mode.

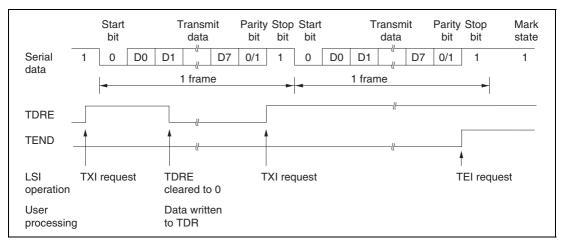


Figure 10.7 Example of Operation when Transmitting in Asynchronous Mode (8-bit data, parity, 1 stop bit)

Receiving: Figure 10.8 shows an example of a flowchart for data reception. This procedure should be followed for data reception after initializing SCI3.

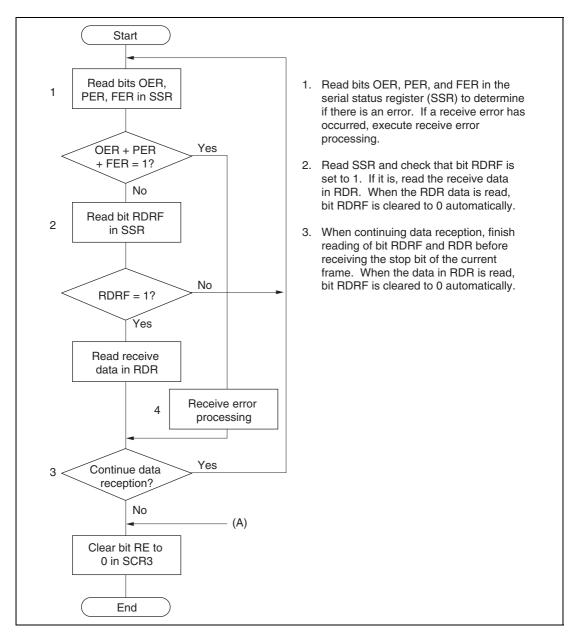


Figure 10.8 Example of Data Reception Flowchart (Asynchronous Mode)

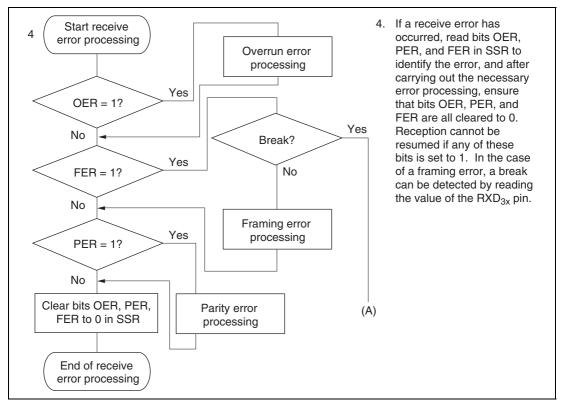


Figure 10.8 Example of Data Reception Flowchart (Asynchronous Mode) (cont)

SCI3 operates as follows when receiving data.

SCI3 monitors the communication line, and when it detects a 0 start bit, performs internal synchronization and begins reception. Reception is carried out in accordance with the relevant data transfer format in table 10.11. The received data is first placed in RSR in LSB-to-MSB order, and then the parity bit and stop bit(s) are received. SCI3 then carries out the following checks.

- · Parity check
 - SCI3 checks that the number of 1 bits in the receive data conforms to the parity (odd or even) set in bit PM in the serial mode register (SMR).
- Stop bit check
 SCI3 checks that the stop bit is 1. If two stop bits are used, only the first is checked.
- Status check
 SCI3 checks that bit RDRF is set to 0, indicating that the receive data can be transferred from RSR to RDR.

If no receive error is found in the above checks, bit RDRF is set to 1, and the receive data is stored in RDR. If bit RIE is set to 1 in SCR3, an RXI interrupt is requested. If the error checks identify a receive error, bit OER, PER, or FER is set to 1 depending on the kind of error. Bit RDRF retains its state prior to receiving the data. If bit RIE is set to 1 in SCR3, an ERI interrupt is requested.

Table 10.12 shows the conditions for detecting a receive error, and receive data processing.

Note: No further receive operations are possible while a receive error flag is set. Bits OER, FER, PER, and RDRF must therefore be cleared to 0 before resuming reception.

Table 10.12 Receive Error Detection Conditions and Receive Data Processing

Receive Error	Abbr.	Detection Conditions	Receive Data Processing
Overrun error	OER	When the next date receive operation is completed while bit RDRF is still set to 1 in SSR	Receive data is not transferred from RSR to RDR
Framing error	FER	When the stop bit is 0	Receive data is transferred from RSR to RDR
Parity error	PER	When the parity (odd or even) set in SMR is different from that of the received data	Receive data is transferred from RSR to RDR

Figure 10.9 shows an example of the operation when receiving in asynchronous mode.

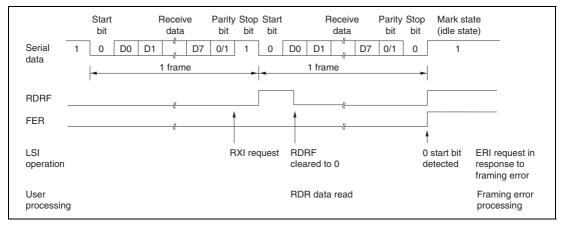


Figure 10.9 Example of Operation when Receiving in Asynchronous Mode (8-bit data, parity, 1 stop bit)

10.3.3 Operation in Synchronous Mode

In synchronous mode, SCI3 transmits and receives data in synchronization with clock pulses. This mode is suitable for high-speed serial communication.

SCI3 has separate transmission and reception units, allowing full-duplex communication with a shared clock.

As the transmission and reception units are both double-buffered, data can be written during transmission and read during reception, making possible continuous transmission and reception.

(1) Data Transfer Format

The general data transfer format in asynchronous communication is shown in figure 10.10.

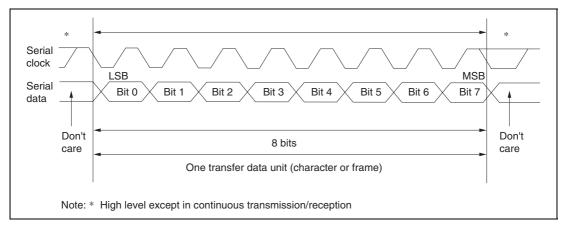


Figure 10.10 Data Format in Synchronous Communication

In synchronous communication, data on the communication line is output from one falling edge of the serial clock until the next falling edge. Data confirmation is guaranteed at the rising edge of the serial clock.

One transfer data character begins with the LSB and ends with the MSB. After output of the MSB, the communication line retains the MSB state.

When receiving in synchronous mode, SCI3 latches receive data at the rising edge of the serial clock.

The data transfer format uses a fixed 8-bit data length.

Parity bits cannot be added.

(2) Clock

Either an internal clock generated by the baud rate generator or an external clock input at the SCK_{3x} pin can be selected as the SCI3 serial clock. The selection is made by means of bit COM in SMR and bits CKE1 and CKE0 in SCR3. See table 10.9 for details on clock source selection.

When SCI3 operates on an internal clock, the serial clock is output at the SCK_{3x} pin. Eight pulses of the serial clock are output in transmission or reception of one character, and when SCI3 is not transmitting or receiving, the clock is fixed at the high level.

(3) Data Transfer Operations

SCI3 initialization: Data transfer on SCI3 first of all requires that SCI3 be initialized as described in 10.3.2.3. SCI3 initialization, and shown in figure 10.5.

Transmitting: Figure 10.11 shows an example of a flowchart for data transmission. This procedure should be followed for data transmission after initializing SCI3.

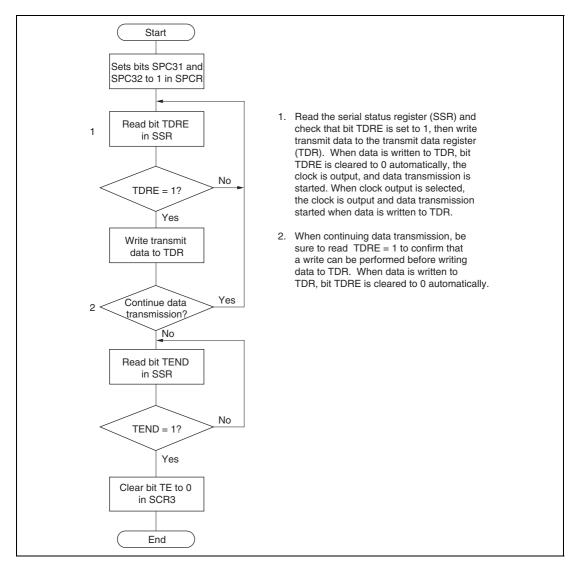


Figure 10.11 Example of Data Transmission Flowchart (Synchronous Mode)

SCI3 operates as follows when transmitting data.

SCI3 monitors bit TDRE in SSR, and when it is cleared to 0, recognizes that data has been written to TDR and transfers data from TDR to TSR. It then sets bit TDRE to 1 and starts transmitting. If bit TIE in SCR3 is set to 1 at this time, a TXI request is made.

When clock output mode is selected, SCI3 outputs 8 serial clock pulses. When an external clock is selected, data is output in synchronization with the input clock.

Serial data is transmitted from the TXD3x pin in order from the LSB (bit 0) to the MSB (bit 7). When the MSB (bit 7) is sent, checks bit TDRE. If bit TDRE is cleared to 0, SCI3 transfers data from TDR to TSR, and starts transmission of the next frame. If bit TDRE is set to 1, SCI3 sets bit TEND to 1 in SSR, and after sending the MSB (bit 7), retains the MSB state. If bit TEIE in SCR3 is set to 1 at this time, a TEI request is made.

After transmission ends, the SCK pin is fixed at the high level.

Note: Transmission is not possible if an error flag (OER, FER, or PER) that indicates the data reception status is set to 1. Check that these error flags are all cleared to 0 before a transmit operation.

Figure 10.12 shows an example of the operation when transmitting in synchronous mode.

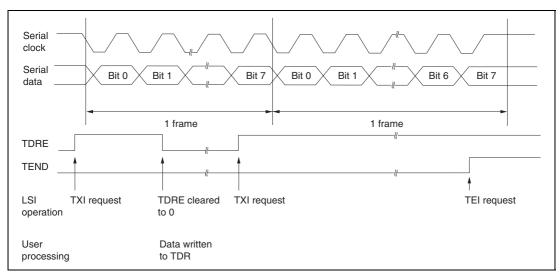


Figure 10.12 Example of Operation when Transmitting in Synchronous Mode

Receiving: Figure 10.13 shows an example of a flowchart for data reception. This procedure should be followed for data reception after initializing SCI3.

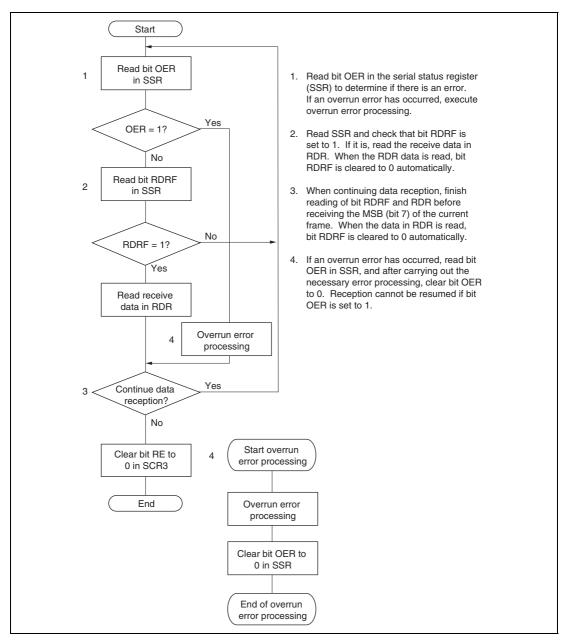


Figure 10.13 Example of Data Reception Flowchart (Synchronous Mode)

SCI3 operates as follows when receiving data.

SCI3 performs internal synchronization and begins reception in synchronization with the serial clock input or output.

The received data is placed in RSR in LSB-to-MSB order.

After the data has been received, SCI3 checks that bit RDRF is set to 0, indicating that the receive data can be transferred from RSR to RDR.

If this check shows that there is no overrun error, bit RDRF is set to 1, and the receive data is stored in RDR. If bit RIE is set to 1 in SCR3, an RXI interrupt is requested. If the check identifies an overrun error, bit OER is set to 1.

Bit RDRF remains set to 1. If bit RIE is set to 1 in SCR3, an ERI interrupt is requested.

See table 10.12 for the conditions for detecting a receive error, and receive data processing.

Note: No further receive operations are possible while a receive error flag is set. Bits OER, FER, PER, and RDRF must therefore be cleared to 0 before resuming reception.

Figure 10.14 shows an example of the operation when receiving in synchronous mode.

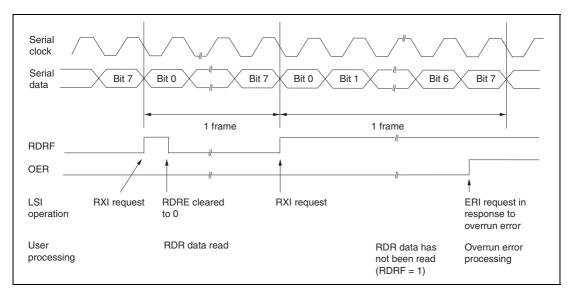


Figure 10.14 Example of Operation when Receiving in Synchronous Mode

Simultaneous transmit/receive: Figure 10.15 shows an example of a flowchart for a simultaneous transmit/receive operation. This procedure should be followed for simultaneous transmission/reception after initializing SCI3.

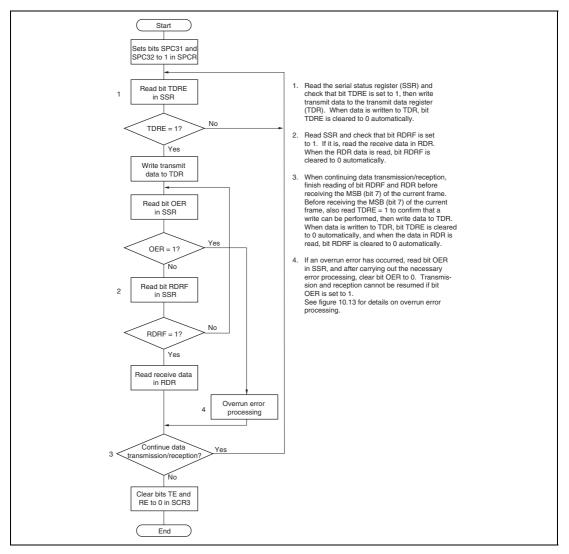


Figure 10.15 Example of Simultaneous Data Transmission/Reception Flowchart (Synchronous Mode)

- Notes: 1. When switching from transmission to simultaneous transmission/reception, check that SCI3 has finished transmitting and that bits TDRE and TEND are set to 1, clear bit TE to 0, and then set bits TE and RE to 1 simultaneously.
 - 2. When switching from reception to simultaneous transmission/reception, check that SCI3 has finished receiving, clear bit RE to 0, then check that bit RDRF and the error flags (OER, FER, and PER) are cleared to 0, and finally set bits TE and RE to 1 simultaneously.

10.4 Interrupts

SCI3 can generate six kinds of interrupts: transmit end, transmit data empty, receive data full, and three receive error interrupts (overrun error, framing error, and parity error). These interrupts have the same vector address.

The various interrupt requests are shown in table 10.13.

Table 10.13 SCI3 Interrupt Requests

Interrupt Abbr.	Interrupt Request	Vector Address
RXI	Interrupt request initiated by receive data full flag (RDRF)	H'0022/H'0024
TXI	Interrupt request initiated by transmit data empty flag (TDRE)	_
TEI	Interrupt request initiated by transmit end flag (TEND)	_
ERI	Interrupt request initiated by receive error flag (OER, FER, PER)	

Each interrupt request can be enabled or disabled by means of bits TIE and RIE in SCR3.

When bit TDRE is set to 1 in SSR, a TXI interrupt is requested. When bit TEND is set to 1 in SSR, a TEI interrupt is requested. These two interrupts are generated during transmission.

The initial value of bit TDRE in SSR is 1. Therefore, if the transmit data empty interrupt request (TXI) is enabled by setting bit TIE to 1 in SCR3 before transmit data is transferred to TDR, a TXI interrupt will be requested even if the transmit data is not ready.

Also, the initial value of bit TEND in SSR is 1. Therefore, if the transmit end interrupt request (TEI) is enabled by setting bit TEIE to 1 in SCR3 before transmit data is transferred to TDR, a TEI interrupt will be requested even if the transmit data has not been sent.

Effective use of these interrupt requests can be made by having processing that transfers transmit data to TDR carried out in the interrupt service routine.

To prevent the generation of these interrupt requests (TXI and TEI), on the other hand, the enable bits for these interrupt requests (bits TIE and TEIE) should be set to 1 after transmit data has been transferred to TDR.

When bit RDRF is set to 1 in SSR, an RXI interrupt is requested, and if any of bits OER, PER, and FER is set to 1, an ERI interrupt is requested. These two interrupt requests are generated during reception.

For further details, see section 3.3, Interrupts.



10.5 Application Notes

The following points should be noted when using SCI3.

(1) Relation between Writes to TDR and Bit TDRE

Bit TDRE in the serial status register (SSR) is a status flag that indicates that data for serial transmission has not been prepared in TDR. When data is written to TDR, bit TDRE is cleared to 0 automatically. When SCI3 transfers data from TDR to TSR, bit TDRE is set to 1.

Data can be written to TDR irrespective of the state of bit TDRE, but if new data is written to TDR while bit TDRE is cleared to 0, the data previously stored in TDR will be lost of it has not yet been transferred to TSR. Accordingly, to ensure that serial transmission is performed dependably, you should first check that bit TDRE is set to 1, then write the transmit data to TDR once only (not two or more times).

(2) Operation when a Number of Receive Errors Occur Simultaneously

If a number of receive errors are detected simultaneously, the status flags in SSR will be set to the states shown in table 10.14. If an overrun error is detected, data transfer from RSR to RDR will not be performed, and the receive data will be lost.

Table 10.14 SSR Status Flag States and Receive Data Transfer

SSR Status Flags			gs	Receive Data Transfer	
RDRF*	OER	FER	PER	$RSR \to RDR$	Receive Error Status
1	1	0	0	X	Overrun error
0	0	1	0	0	Framing error
0	0	0	1	0	Parity error
1	1	1	0	X	Overrun error + framing error
1	1	0	1	X	Overrun error + parity error
0	0	1	1	0	Framing error + parity error
1	1	1	1	X	Overrun error + framing error + parity error

O: Receive data is transferred from RSR to RDR.

X: Receive data is not transferred from RSR to RDR.

Note: * Bit RDRF retains its state prior to data reception. However, note that if RDR is read after an overrun error has occurred in a frame because reading of the receive data in the previous frame was delayed, RDRF will be cleared to 0.

(3) Break Detection and Processing

When a framing error is detected, a break can be detected by reading the value of the RXD_{3x} pin directly. In a break, the input from the RXD_{3x} pin becomes all 0s, with the result that bit FER is set and bit PER may also be set.

SCI3 continues the receive operation even after receiving a break. Note, therefore, that even though bit FER is cleared to 0 it will be set to 1 again.

(4) Mark State and Break Detection

When bit TE is cleared to 0, the TXD_{3x} pin functions as an I/O port whose input/output direction and level are determined by PDR and PCR. This fact can be used to set the TXD_{3x} pin to the mark state, or to detect a break during transmission.

To keep the communication line in the mark state (1 state) until bit TE is set to 1, set PCR = 1 and PDR = 1. Since bit TE is cleared to 0 at this time, the TXD_{3x} pin functions as an I/O port and 1 is output.

To detect a break, clear bit TE to 0 after setting PCR = 1 and PDR = 0.

When bit TE is cleared to 0, the transmission unit is initialized regardless of the current transmission state, the TXD₃, pin functions as an I/O port, and 0 is output from the TXD₃, pin.

(5) Receive Error Flags and Transmit Operation (Synchronous Mode Only)

When a receive error flag (OER, PER, or FER) is set to 1, transmission cannot be started even if bit TDRE is cleared to 0. The receive error flags must be cleared to 0 before starting transmission.

Note also that receive error flags cannot be cleared to 0 even if bit RE is cleared to 0.

(6) Receive Data Sampling Timing and Receive Margin in Asynchronous Mode

In asynchronous mode, SCI3 operates on a basic clock with a frequency 16 times the transfer rate. When receiving, SCI3 performs internal synchronization by sampling the falling edge of the start bit with the basic clock. Receive data is latched internally at the 8th rising edge of the basic clock. This is illustrated in figure 10.16.



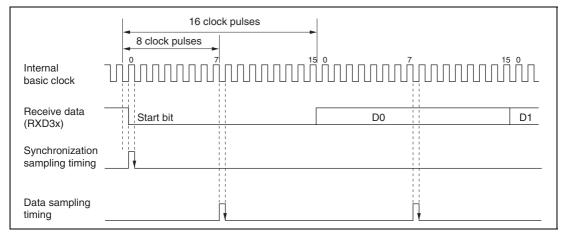


Figure 10.16 Receive Data Sampling Timing in Asynchronous Mode

Consequently, the receive margin in asynchronous mode can be expressed as shown in equation (1).

$$M = \{(0.5 - \frac{1}{2N}) - \frac{D - 0.5}{N} - (L - 0.5) F\} \times 100 [\%]$$
 Equation (1)

where M: Receive margin (%)

N: Ratio of bit rate to clock (N = 16)

D: Clock duty (D = 0.5 to 1.0)

L: Frame length (L = 9 to 12)

F: Absolute value of clock frequency deviation

Substituting 0 for F (absolute value of clock frequency deviation) and 0.5 for D (clock duty) in equation (1), a receive margin of 46.875% is given by equation (2).

When D = 0.5 and F = 0,

$$M = \{0.5 - 1/(2 \times 16)\} \times 100 \, [\%]$$

= 46.875\% Equation (2)

However, this is only a computed value, and a margin of 20% to 30% should be allowed when carrying out system design.

(7) Relation between RDR Reads and Bit RDRF

In a receive operation, SCI3 continually checks the RDRF flag. If bit RDRF is cleared to 0 when reception of one frame ends, normal data reception is completed. If bit RDRF is set to 1, this indicates that an overrun error has occurred.

When the contents of RDR are read, bit RDRF is cleared to 0 automatically. Therefore, if bit RDR is read more than once, the second and subsequent read operations will be performed while bit RDRF is cleared to 0. Note that, when an RDR read is performed while bit RDRF is cleared to 0, if the read operation coincides with completion of reception of a frame, the next frame of data may be read. This is illustrated in figure 10.17.

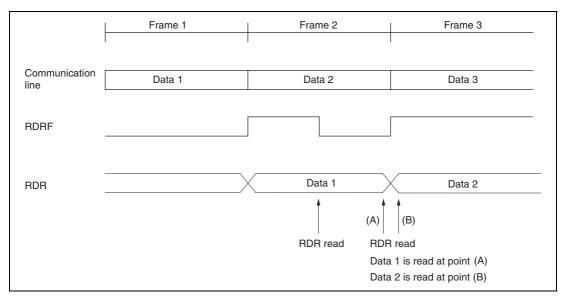


Figure 10.17 Relation between RDR Read Timing and Data

In this case, only a single RDR read operation (not two or more) should be performed after first checking that bit RDRF is set to 1. If two or more reads are performed, the data read the first time should be transferred to RAM, etc., and the RAM contents used. Also, ensure that there is sufficient margin in an RDR read operation before reception of the next frame is completed. To be precise in terms of timing, the RDR read should be completed before bit 7 is transferred in synchronous mode, or before the STOP bit is transferred in asynchronous mode.

(8) Transmit and Receive Operations when Making a State Transition

Make sure that transmit and receive operations have completely finished before carrying out state transition processing.

(9) Switching SCK₃, Function

If pin SCK_{3x} is used as a clock output pin by SCI3 in synchronous mode and is then switched to a general input/output pin (a pin with a different function), the pin outputs a low level signal for half a system clock (ϕ) cycle immediately after it is switched.

This can be prevented by either of the following methods according to the situation.

- a. When an SCK_{3x} function is switched from clock output to non clock-output When stopping data transfer, issue one instruction to clear bits TE and RE to 0 and to set bits CKE1 and CKE0 in SCR3 to 1 and 0, respectively. In this case, bit COM in SMR should be left 1. The above prevents SCK_{3x} from being used as a general input/output pin. To avoid an intermediate level of voltage from being applied to SCK_{3x}, the line connected to SCK_{3x} should be pulled up to the V_{CC} level via a resistor, or supplied with output from an external device.
- b. When an SCK_{3x} function is switched from clock output to general input/output When stopping data transfer,
 - (i) Issue one instruction to clear bits TE and RE to 0 and to set bits CKE1 and CKE0 in SCR3 to 1 and 0, respectively.
 - (ii) Clear bit COM in SMR to 0
 - (iii) Clear bits CKE1 and CKE0 in SCR3 to 0

Note that special care is also needed here to avoid an intermediate level of voltage from being applied to SCK_{3x} .

(10) Setup at Subactive or Subsleep Mode

At subactive or subsleep mode, SCI3 becomes possible use only at CPU clock is $\phi w/2$.

Section 11 14-Bit PWM

11.1 Overview

This LSI is provided with a 14-bit PWM (pulse width modulator) on-chip, which can be used as a D/A converter by connecting a low-pass filter.

11.1.1 Features

Features of the 14-bit PWM are as follows.

- Choice of two conversion periods
 - Any of the following four conversion periods can be chosen:
 - 131,072/ ϕ , with a minimum modulation width of 8/ ϕ (PWCR1 = 1, PWCR0 = 1)
 - $-65,536/\phi$, with a minimum modulation width of $4/\phi$ (PWCR1 = 1, PWCR0 = 0)
 - $-32,768/\phi$, with a minimum modulation width of $2/\phi$ (PWCR1 = 0, PWCR0 = 1)
 - $16,384/\phi$, with a minimum modulation width of $1/\phi$ (PWCR1 = 0, PWCR0 = 0)
- Pulse division method for less ripple
- Use of module standby mode enables this module to be placed in standby mode independently when not used.

11.1.2 Block Diagram

Figure 11.1 shows a block diagram of the 14-bit PWM.

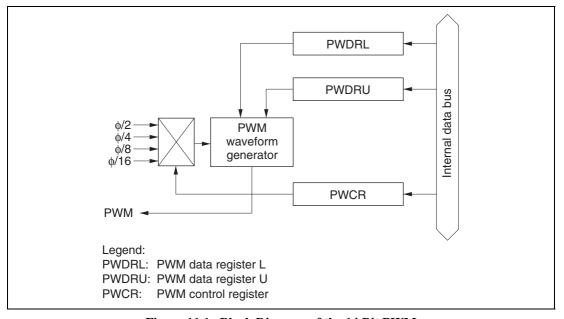


Figure 11.1 Block Diagram of the 14 Bit PWM

11.1.3 Pin Configuration

Table 11.1 shows the output pin assigned to the 14-bit PWM.

Table 11.1 Pin Configuration

Name	Abbr.	I/O	Function
PWM output pin	PWM	Output	Pulse-division PWM waveform output

11.1.4 Register Configuration

Table 11.2 shows the register configuration of the 14-bit PWM.

Table 11.2 Register Configuration

Name	Abbr.	R/W	Initial Value	Address
PWM control register	PWCR	W	H'FC	H'FFD0
PWM data register U	PWDRU	W	H'C0	H'FFD1
PWM data register L	PWDRL	W	H'00	H'FFD2
Clock stop register 2	CKSTPR2	R/W	H'FF	H'FFFB

11.2 Register Descriptions

11.2.1 PWM Control Register (PWCR)

Bit	7	6	5	4	3	2	1	0
	_	_	_	_	_	_	PWCR1	PWCR0
Initial value	1	1	1	1	1	1	0	0
Read/Write	_	_	_	_	_	_	W	W

PWCR is an 8-bit write-only register for input clock selection.

Upon reset, PWCR is initialized to HFC.

Bits 7 to 2: Reserved bits

Bits 7 to 2 are reserved; they are always read as 1, and cannot be modified.

Bits 1 and 0: Clock select 1 and 0 (PWCR1, PWCR0)

Bits 1 and 0 select the clock supplied to the 14-bit PWM. These bits are write-only bits; they are always read as 1.

Bit 1 PWCR1	Bit 0 PWCR0	Description	
0	0	The input clock is $\phi/2$ ($t\phi^*=2/\phi$) The conversion period is 16,384/ ϕ , with a minimum modulation width of $1/\phi$	(initial value)
0	1	The input clock is $\phi/4$ ($t\phi^*=4/\phi$) The conversion period is 32,768/ ϕ , with a minimum modulation width of 2/ ϕ	
1	0	The input clock is $\phi/8$ (t $\phi^*=8/\phi$) The conversion period is 65,536/ ϕ , with a minimum modulation width of $4/\phi$	
1	1	The input clock is $\phi/16$ ($t\phi^*=16/\phi$) The conversion period is 131,072/ ϕ , with a minimum modulation width of $8/\phi$	

Note: * Period of PWM input clock.

11.2.2 PWM Data Registers U and L (PWDRU, PWDRL)

PWDRU								
Bit	7	6	5	4	3	2	1	0
	_	_	PWDRU5	PWDRU4	PWDRU3	PWDRU2	PWDRU1	PWDRU0
Initial value	1	1	0	0	0	0	0	0
Read/Write	_	_	W	W	W	W	W	W
PWDRL								
Bit	7	6	5	4	3	2	1	0
	PWDRL7	PWDRL6	PWDRL5	PWDRL4	PWDRL3	PWDRL2	PWDRL1	PWDRL0
Initial value	0	0	0	0	0	0	0	0
Read/Write	W	W	W	W	W	W	W	W

PWDRU and PWDRL form a 14-bit write-only register, with the upper 6 bits assigned to PWDRU and the lower 8 bits to PWDRL. The value written to PWDRU and PWDRL gives the total high-level width of one PWM waveform cycle.

When 14-bit data is written to PWDRU and PWDRL, the register contents are latched in the PWM waveform generator, updating the PWM waveform generation data. The 14-bit data should always be written in the following sequence:

- 1. Write the lower 8 bits to PWDRL.
- 2. Write the upper 6 bits to PWDRU.

PWDRU and PWDRL are write-only registers. If they are read, all bits are read as 1.

Upon reset, PWDRU and PWDRL are initialized to H'C000.

11.2.3 Clock Stop Register 2 (CKSTPR2)

Bit	7	6	5	4	3	2	1	0
	_		_	_	AECKSTP	WDCKSTP	PWCKSTP	LDCKSTP
Initial value	1	1	1	1	1	1	1	1
Read/Write		_	_	_	R/W	R/W	R/W	R/W

CKSTPR2 is an 8-bit read/write register that performs module standby mode control for peripheral modules. Only the bit relating to the PWM is described here. For details of the other bits, see the sections on the relevant modules.

Bit 1: PWM module standby mode control (PWCKSTP)

Bit 1 controls setting and clearing of module standby mode for the PWM.

PWCKSTP	Description	
0	PWM is set to module standby mode	
1	PWM module standby mode is cleared	(initial value)

11.3 Operation

11.3.1 Operation

When using the 14-bit PWM, set the registers in the following sequence.

- 1. Set bit PWM in port mode register 3 (PMR3) to 1 so that pin P3₀/PWM is designated for PWM output.
- 2. Set bits PWCR1 and PWCR0 in the PWM control register (PWCR) to select a conversion period of $131,072/\phi$ (PWCR1 = 1, PWCR0 = 1), $65,536/\phi$ (PWCR1 = 1, PWCR0 = 0), $32,768/\phi$ (PWCR1 = 0, PWCR0 = 1), or $16,384/\phi$ (PWCR1 = 0, PWCR0 = 0).
- 3. Set the output waveform data in PWM data registers U and L (PWDRU/L). Be sure to write in the correct sequence, first PWDRL then PWDRU. When data is written to PWDRU, the data in these registers will be latched in the PWM waveform generator, updating the PWM waveform generation in synchronization with internal signals.

One conversion period consists of 64 pulses, as shown in figure 11.2. The total of the high-level pulse widths during this period (T_H) corresponds to the data in PWDRU and PWDRL. This relation can be represented as follows.

$$T_{H} = (data \ value \ in \ PWDRU \ and \ PWDRL + 64) \times t_{0}/2$$

where $t\phi$ is the PWM input clock period: $2/\phi$ (PWCR = H'0), $4/\phi$ (PWCR = H'1), $8/\phi$ (PWCR = H'2), or $16/\phi$ (PWCR = H'3).

Example: Settings in order to obtain a conversion period of 32,768 μs:

When PWCR1 = 0 and PWCR0 = 0, the conversion period is $16,384/\phi$, so ϕ must be 0.5 MHz. In this case, tfn = 512 μ s, with $1/\phi$ (resolution) = 2.0 μ s.

When PWCR1 = 0 and PWCR0 = 1, the conversion period is $32,768/\phi$, so ϕ must be 1 MHz. In this case, tfn = $512 \mu s$, with $2/\phi$ (resolution) = $2.0 \mu s$.

When PWCR1 = 1 and PWCR0 = 0, the conversion period is $65,536/\phi$, so ϕ must be 2 MHz. In this case, tfn = $512 \mu s$, with $4/\phi$ (resolution) = $2.0 \mu s$.

Accordingly, for a conversion period of 32,768 μ s, the system clock frequency (ϕ) must be 0.5 MHz, 1 MHz, or 2 MHz.



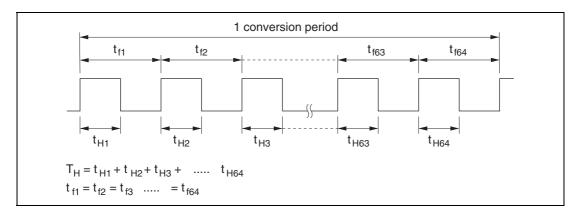


Figure 11.2 PWM Output Waveform

11.3.2 PWM Operation Modes

PWM operation modes are shown in table 11.3.

Table 11.3 PWM Operation Modes

Operation Mode	Reset	Active	Sleep	Watch	Subactive	Subsleep	Standby	Module Standby
PWCR	Reset	Functions	Functions	Held	Held	Held	Held	Held
PWDRU	Reset	Functions	Functions	Held	Held	Held	Held	Held
PWDRL	Reset	Functions	Functions	Held	Held	Held	Held	Held

Section 12 A/D Converter

12.1 Overview

This LSI includes on-chip a resistance-ladder-based successive-approximation analog-to-digital converter, and can convert up to 8 channels of analog input.

12.1.1 Features

The A/D converter has the following features.

- 10-bit resolution
- Eight input channels
- Conversion time: approx. 12.4 µs per channel (at 5 MHz operation)
- Built-in sample-and-hold function
- Interrupt requested on completion of A/D conversion
- A/D conversion can be started by external trigger input
- Use of module standby mode enables this module to be placed in standby mode independently when not used.

12.1.2 Block Diagram

Figure 12.1 shows a block diagram of the A/D converter.

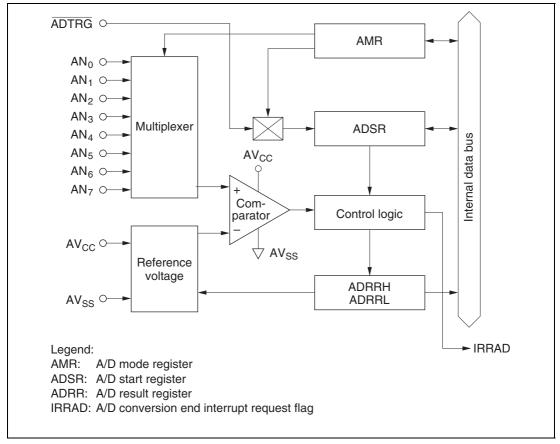


Figure 12.1 Block Diagram of the A/D Converter

12.1.3 Pin Configuration

Table 12.1 shows the A/D converter pin configuration.

Table 12.1 Pin Configuration

Name	Abbr.	I/O	Function
Analog power supply	AV _{cc}	Input	Power supply and reference voltage of analog part
Analog ground	AV _{ss}	Input	Ground and reference voltage of analog part
Analog input 0	AN _o	Input	Analog input channel 0
Analog input 1	AN ₁	Input	Analog input channel 1
Analog input 2	AN ₂	Input	Analog input channel 2
Analog input 3	AN ₃	Input	Analog input channel 3
Analog input 4	AN ₄	Input	Analog input channel 4
Analog input 5	AN ₅	Input	Analog input channel 5
Analog input 6	AN ₆	Input	Analog input channel 6
Analog input 7	AN ₇	Input	Analog input channel 7
External trigger input	ADTRG	Input	External trigger input for starting A/D conversion

12.1.4 Register Configuration

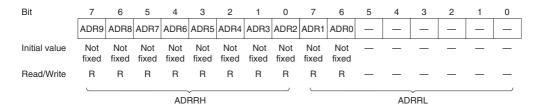
Table 12.2 shows the A/D converter register configuration.

Table 12.2 Register Configuration

Name	Abbr.	R/W	Initial Value	Address
A/D mode register	AMR	R/W	H'30	H'FFC6
A/D start register	ADSR	R/W	H'7F	H'FFC7
A/D result register H	ADRRH	R	Not fixed	H'FFC4
A/D result register L	ADRRL	R	Not fixed	H'FFC5
Clock stop register 1	CKSTPR1	R/W	H'FF	H'FFFA

12.2 Register Descriptions

12.2.1 A/D Result Registers (ADRRH, ADRRL)



ADRRH and ADRRL together comprise a 16-bit read-only register for holding the results of analog-to-digital conversion. The upper 8 bits of the data are held in ADRRH, and the lower 2 bits in ADRRL.

ADRRH and ADRRL can be read by the CPU at any time, but the ADRRH and ADRRL values during A/D conversion are not fixed. After A/D conversion is complete, the conversion result is stored as 10-bit data, and this data is held until the next conversion operation starts.

ADRRH and ADRRL are not cleared on reset.

12.2.2 A/D Mode Register (AMR)

Bit	7	6	5	4	3	2	1	0
	CKS	TRGE	_	_	CH3	CH2	CH1	CH0
Initial value	0	0	1	1	0	0	0	0
Read/Write	R/W	R/W	_	_	R/W	R/W	R/W	R/W

AMR is an 8-bit read/write register for specifying the A/D conversion speed, external trigger option, and the analog input pins.

Upon reset, AMR is initialized to H'30.



Bit 7: Clock select (CKS)

Bit 7 sets the A/D conversion speed.

Bit 7		Conversion Time	ne (Active (High-Speed) Mode)*
CKS	Conversion Period	φ = 1 MHz	φ = 5 MHz
0	62/φ (initial value)	62 µs	12.4 µs
1	31/φ	31 µs	_

Note: * For information on conversion time settings for which operation is guaranteed, see section 16, Electrical Characteristics.

Bit 6: External trigger select (TRGE)

Bit 6 enables or disables the start of A/D conversion by external trigger input.

Bit 6 TRGE	Description	
0	Disables start of A/D conversion by external trigger	(initial value)
1	Enables start of A/D conversion by rising or falling edge of external trip $\overline{\text{ADTRG}}^*$	gger at pin

Note: * The external trigger (ADTRG) edge is selected by bit IEG4 of IEGR. See (1) IRQ Edge Select Register (IEGR) in section 3.3.2, Interrupt Control Registers, for details.

Bits 5 and 4: Reserved bits

Bits 5 and 4 are reserved; they are always read as 1, and cannot be modified.

Bits 3 to 0: Channel select (CH3 to CH0)

Bits 3 to 0 select the analog input channel.

The channel selection should be made while bit ADSF is cleared to 0.

Bit 3 CH3	Bit 2 CH2	Bit 1 CH1	Bit 0 CH0	Analog Input Channel	
0	0	*	*	No channel selected	(initial value)
0	1	0	0	AN0	
0	1	0	1	AN1	_
0	1	1	0	AN2	_
0	1	1	1	AN3	
1	0	0	0	AN4	
1	0	0	1	AN5	_
1	0	1	0	AN6	
1	0	1	1	AN7	
1	1	*	*	Setting prohibited	

Note: * Don't care

12.2.3 A/D Start Register (ADSR)

Bit	7	6	5	4	3	2	1	0	
	ADSF	_	_	_	_	_	_	_	
Initial value	0	1	1	1	1	1	1	1	_
Read/Write	R/W	_	_	_		_			

The A/D start register (ADSR) is an 8-bit read/write register for starting and stopping A/D conversion.

A/D conversion is started by writing 1 to the A/D start flag (ADSF) or by input of the designated edge of the external trigger signal, which also sets ADSF to 1. When conversion is complete, the converted data is set in ADRRH and ADRRL, and at the same time ADSF is cleared to 0.

Bit 7: A/D start flag (ADSF)

Bit 7 controls and indicates the start and end of A/D conversion.

Bit 7 ADSF	Description	
0	Read: Indicates the completion of A/D conversion	(initial value)
	Write: Stops A/D conversion	
1	Read: Indicates A/D conversion in progress	
	Write: Starts A/D conversion	

Bits 6 to 0: Reserved bits

Bits 6 to 0 are reserved; they are always read as 1, and cannot be modified.

12.2.4 Clock Stop Register 1 (CKSTPR1)

Bit	7	6	5	4	3	2	1	0
	_	S31CKSTP	S32CKSTP	ADCKSTP	TGCKSTP	TFCKSTP	TCCKSTP	TACKSTP
Initial value	1	1	1	1	1	1	1	1
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W

CKSTPR1 is an 8-bit read/write register that performs module standby mode control for peripheral modules. Only the bit relating to the A/D converter is described here. For details of the other bits, see the sections on the relevant modules.

Bit 4: A/D converter module standby mode control (ADCKSTP)

Bit 4 controls setting and clearing of module standby mode for the A/D converter.

ADCKSTP	Description	
0	A/D converter is set to module standby mode	
1	A/D converter module standby mode is cleared	(initial value)

12.3 Operation

12.3.1 A/D Conversion Operation

The A/D converter operates by successive approximations, and yields its conversion result as 10-bit data.

A/D conversion begins when software sets the A/D start flag (bit ADSF) to 1. Bit ADSF keeps a value of 1 during A/D conversion, and is cleared to 0 automatically when conversion is complete.

The completion of conversion also sets bit IRRAD in interrupt request register 2 (IRR2) to 1. An A/D conversion end interrupt is requested if bit IENAD in interrupt enable register 2 (IENR2) is set to 1.

If the conversion time or input channel needs to be changed in the A/D mode register (AMR) during A/D conversion, bit ADSF should first be cleared to 0, stopping the conversion operation, in order to avoid malfunction.

12.3.2 Start of A/D Conversion by External Trigger Input

The A/D converter can be made to start A/D conversion by input of an external trigger signal. External trigger input is enabled at pin ADTRG when bit IRQ4 in PMR1 is set to 1 and bit TRGE in AMR is set to 1. Then when the input signal edge designated in bit IEG4 of interrupt edge select register (IEGR) is detected at pin ADTRG, bit ADSF in ADSR will be set to 1, starting A/D conversion.

Figure 12.2 shows the timing.

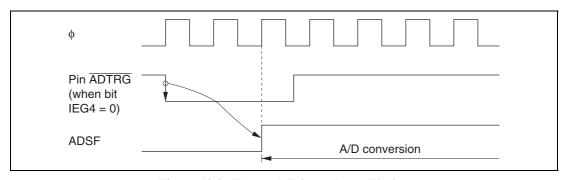


Figure 12.2 External Trigger Input Timing

12.3.3 A/D Converter Operation Modes

A/D converter operation modes are shown in table 12.3.

Table 12.3 A/D Converter Operation Modes

Operation Mode	Reset	Active	Sleep	Watch	Subactive	Subsleep	Standby	Module Standby
AMR	Reset	Functions	Functions	Held	Held	Held	Held	Held
ADSR	Reset	Functions	Functions	Held	Held	Held	Held	Held
ADRRH	Held*	Functions	Functions	Held	Held	Held	Held	Held
ADRRL	Held*	Functions	Functions	Held	Held	Held	Held	Held

Note: * Undefined in a power-on reset.

12.4 Interrupts

When A/D conversion ends (ADSF changes from 1 to 0), bit IRRAD in interrupt request register 2 (IRR2) is set to 1.

A/D conversion end interrupts can be enabled or disabled by means of bit IENAD in interrupt enable register 2 (IENR2).

For further details see section 3.3, Interrupts.

12.5 Typical Use

An example of how the A/D converter can be used is given below, using channel 1 (pin AN1) as the analog input channel. Figure 12.3 shows the operation timing.

- Bits CH3 to CH0 of the A/D mode register (AMR) are set to 0101, making pin AN₁ the analog input channel. A/D interrupts are enabled by setting bit IENAD to 1, and A/D conversion is started by setting bit ADSF to 1.
- 2. When A/D conversion is complete, bit IRRAD is set to 1, and the A/D conversion result is stored is stored in ADRRH and ADRRL. At the same time ADSF is cleared to 0, and the A/D converter goes to the idle state.
- 3. Bit IENAD = 1, so an A/D conversion end interrupt is requested.
- 4. The A/D interrupt handling routine starts.
- 5. The A/D conversion result is read and processed.
- 6. The A/D interrupt handling routine ends.

If ADSF is set to 1 again afterward, A/D conversion starts and steps 2 to 6 take place.

Figures 12.4 and 12.5 show flow charts of procedures for using the A/D converter.



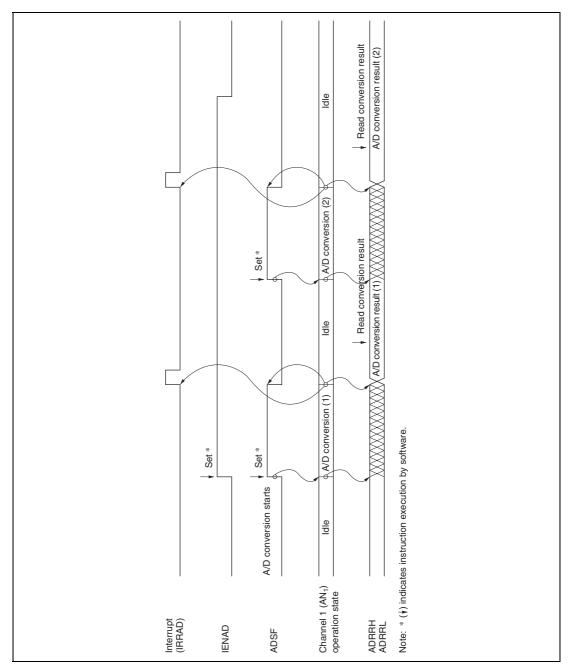


Figure 12.3 Typical A/D Converter Operation Timing

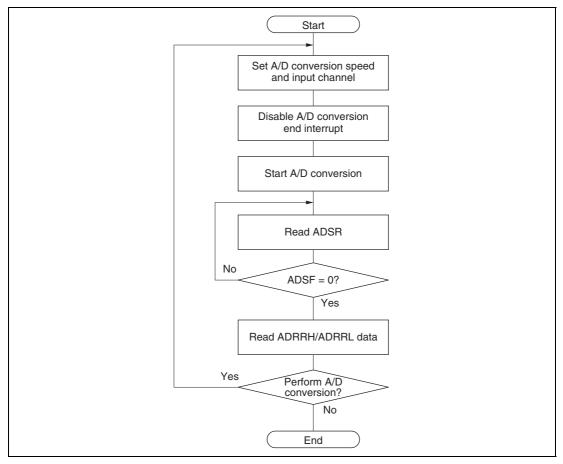


Figure 12.4 Flow Chart of Procedure for Using A/D Converter (Polling by Software)

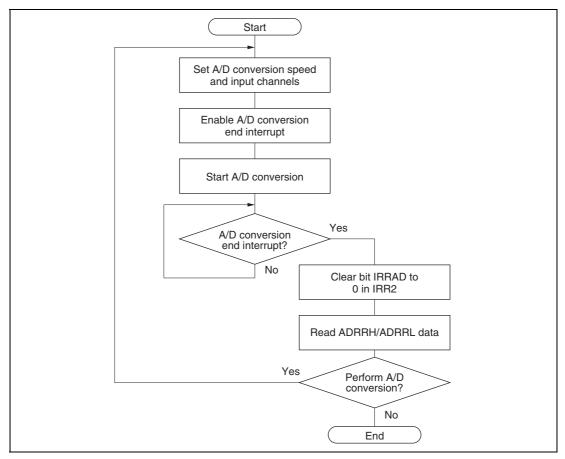


Figure 12.5 Flow Chart of Procedure for Using A/D Converter (Interrupts Used)

12.6 Application Notes

12.6.1 Application Notes

- Data in ADRRH and ADRRL should be read only when the A/D start flag (ADSF) in the A/D start register (ADSR) is cleared to 0.
- Changing the digital input signal at an adjacent pin during A/D conversion may adversely affect conversion accuracy.
- When A/D conversion is started after clearing module standby mode, wait for 10 φ clock cycles before starting.
- In active mode or sleep mode, analog power supply current (AI_{STOPI}) flows into the ladder resistance even when the A/D converter is not operating. Therefore, if the A/D converter is not used, it is recommended that AV_{cc} be connected to the system power supply and the ADCKSTP(A/D converter module standby mode control) bit be cleared to 0 in clock stop register 1 (CKSTPR1).

12.6.2 Permissible Signal Source Impedance

This LSI's analog input is designed such that conversion precision is guaranteed for an input signal for which the signal source impedance is $10~k\Omega$ or less. This specification is provided to enable the A/D converter's sample-and-hold circuit input capacitance to be charged within the sampling time; if the sensor output impedance exceeds $10~k\Omega$, charging may be insufficient and it may not be possible to guarantee A/D conversion precision. However, a large capacitance provided externally, the input load will essentially comprise only the internal input resistance of $10~k\Omega$, and the signal source impedance is ignored. However, as a low-pass filter effect is obtained in this case, it may not be possible to follow an analog signal with a large differential coefficient (e.g., $5~mV/\mu s$ or greater) (see figure 12.6). When converting a high-speed analog signal, a low-impedance buffer should be inserted.



12.6.3 Influences on Absolute Precision

Adding capacitance results in coupling with GND, and therefore noise in GND may adversely affect absolute precision. Be sure to make the connection to an electrically stable GND.

Care is also required to ensure that filter circuits do not interfere with digital signals or act as antennas on the mounting board.

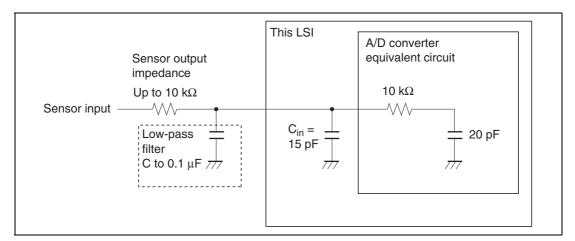


Figure 12.6 Analog Input Circuit Example



Section 13 LCD Controller/Driver

13.1 Overview

This LSI has an on-chip segment type LCD control circuit, LCD driver, and power supply circuit, enabling it to directly drive an LCD panel.

13.1.1 Features

(1) Features

Features of the LCD controller/driver are given below.

• Display capacity

Duty Cycle	Internal Driver
Static	32 seg
1/2	32 seg
1/3	32 seg
1/4	32 seg

- LCD RAM capacity
 - 8 bits \times 32 bytes (256 bits)
- Word access to LCD RAM
- All eight segment output pins can be used individually as port pins.
- Common output pins not used because of the duty cycle can be used for common double-buffering (parallel connection).
- Display possible in operating modes other than standby mode
- Choice of 11 frame frequencies
- Built-in power supply split-resistance, supplying LCD drive power
- Use of module standby mode enables this module to be placed in standby mode independently when not used.
- A or B waveform selectable by software

13.1.2 Block Diagram

Figure 13.1 shows a block diagram of the LCD controller/driver.

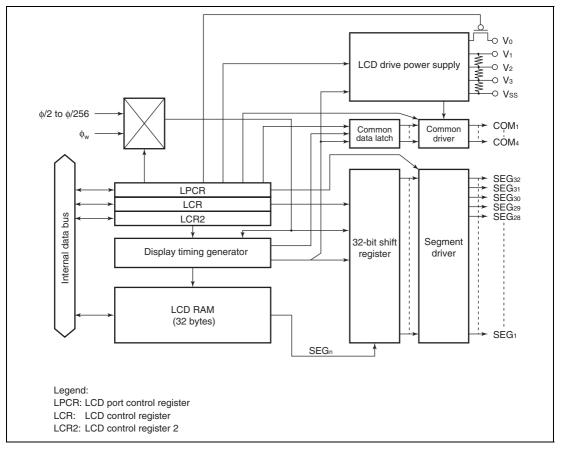


Figure 13.1 Block Diagram of LCD Controller/Driver

13.1.3 Pin Configuration

Table 13.1 shows the LCD controller/driver pin configuration.

Table 13.1 Pin Configuration

Name	Abbr.	1/0	Function
Segment output pins	SEG ₃₂ to SEG ₁	Output	LCD segment drive pins All pins are multiplexed as port pins (setting programmable)
Common output pins	COM ₄ to COM ₁	Output	LCD common drive pins Pins can be used in parallel with static or 1/2 duty
LCD power supply pins	V ₀ , V ₁ , V ₂ , V ₃	_	Used when a bypass capacitor is connected externally, and when an external power supply circuit is used

13.1.4 Register Configuration

Table 13.2 shows the register configuration of the LCD controller/driver.

Table 13.2 LCD Controller/Driver Registers

Name	Abbr.	R/W	Initial Value	Address
LCD port control register	LPCR	R/W	H'00	H'FFC0
LCD control register	LCR	R/W	H'80	H'FFC1
LCD control register 2	LCR2	R/W	H'60	H'FFC2
LCD RAM	_	R/W	Undefined	H'F740, H'F75F
Clock stop register 2	CKSTPR2	R/W	H'FF	H'FFFB

13.2 Register Descriptions

13.2.1 LCD Port Control Register (LPCR)

Bit	7	6	5	4	3	2	1	0
	DTS1	DTS0	CMX	_	SGS3	SGS2	SGS1	SGS0
Initial value	0	0	0	0	0	0	0	0
Read/Write	R/W	R/W	R/W	W	R/W	R/W	R/W	R/W

LPCR is an 8-bit read/write register which selects the duty cycle and LCD driver pin functions.

LPCR is initialized to H'00 upon reset.

Bits 7 to 5: Duty cycle select 1 and 0 (DTS1, DTS0), common function select (CMX)

The combination of DTS1 and DTS0 selects static, 1/2, 1/3, or 1/4 duty. CMX specifies whether or not the same waveform is to be output from multiple pins to increase the common drive power when not all common pins are used because of the duty setting.

Bit 7 DTS1	Bit 6 DTS0	Bit 5 CMX	Duty Cycle	Common Drivers	Notes
0	0	0	Static	COM, (initial value)	Do not use COM ₄ , COM ₃ , and COM ₂ .
		1	_	COM ₄ to COM ₁	COM ₄ , COM ₃ , and COM ₂ output the same waveform as COM ₁ .
0	1	0	1/2 duty	COM ₂ to COM ₁	Do not use COM ₄ and COM ₃ .
		1		COM ₄ to COM ₁	COM ₄ outputs the same waveform as COM ₃ , and COM ₂ outputs the same waveform as COM ₁ .
1	0	0	1/3 duty	COM ₃ to COM ₁	Do not use COM ₄ .
		1		COM ₄ to COM ₁	Do not use COM ₄ .
1	1	0	1/4 duty	COM ₄ to COM ₁	_
		1	_		

Bit 4: Reserved

Bit 4 is reserved. It can only be written with 0.

Bits 3 to 0: Segment driver select 3 to 0 (SGS3 to SGS0)

Bits 3 to 0 select the segment drivers to be used.

				Fun				
Bit 3 SGS3	Bit 2 SGS2	Bit 1 SGS1	Bit 0 SGS0	SEG ₃₂ to SEG ₂₅	SEG ₂₄ to SEG ₁₇	SEG ₁₆ to SEG ₉	SEG ₈ to SEG ₁	 Notes
0	0	0	0	Port	Port	Port	Port	(initial value)
0	0	0	1	Port	Port	Port	Port	
0	0	1	*	SEG	Port	Port	Port	_
0	1	0	*	SEG	SEG	Port	Port	_
0	1	1	*	SEG	SEG	SEG	Port	_
1	*	*	*	SEG	SEG	SEG	SEG	_

*: Don't care

13.2.2 LCD Control Register (LCR)

Bit	7	6	5	4	3	2	1	0
	_	PSW	ACT	DISP	CKS3	CKS2	CKS1	CKS0
Initial value	1	0	0	0	0	0	0	0
Read/Write	_	R/W	R/W	R/W	R/W	R/W	R/W	R/W

LCR is an 8-bit read/write register which performs LCD drive power supply on/off control and display data control, and selects the frame frequency.

LCR is initialized to H'80 upon reset.

Bit 7: Reserved bit

Bit 7 is reserved; it is always read as 1 and cannot be modified.

Bit 6: LCD drive power supply on/off control (PSW)

Bit 6 can be used to turn the LCD drive power supply off when LCD display is not required in a power-down mode, or when an external power supply is used. When the ACT bit is cleared to 0, or in standby mode, the LCD drive power supply is turned off regardless of the setting of this bit.

Bit 6 PSW	Description	
0	LCD drive power supply off	(initial value)
1	LCD drive power supply on	

Bit 5: Display function activate (ACT)

Bit 5 specifies whether or not the LCD controller/driver is used. Clearing this bit to 0 halts operation of the LCD controller/driver. The LCD drive power supply is also turned off, regardless of the setting of the PSW bit. However, register contents are retained.

Bit 5 ACT	Description	
0	LCD controller/driver operation halted	(initial value)
1	LCD controller/driver operates	



Bit 4: Display data control (DISP)

Bit 4 specifies whether the LCD RAM contents are displayed or blank data is displayed regardless of the LCD RAM contents.

Bit 4 DISP	Description	
0	Blank data is displayed	(initial value)
1	LCD RAM data is display	

Bits 3 to 0: Frame frequency select 3 to 0 (CKS3 to CKS0)

Bits 3 to 0 select the operating clock and the frame frequency. In subactive mode, watch mode, and subsleep mode, the system clock (ϕ) is halted, and therefore display operations are not performed if one of the clocks from ϕ /2 to ϕ /256 is selected. If LCD display is required in these modes, ϕ w, ϕ w/2, or ϕ w/4 must be selected as the operating clock.

Bit 3	Bit 2	Bit 1	Bit 0		e Frequency*2	
CKS3	CKS2	CKS1	CKS0	Operating Clock	φ = 2 MHz	φ = 250 kHz*1
0	*	0	0	φ w	128 Hz*³ (initia	l value)
0	*	0	1	φw/2	64 Hz*3	
0	*	1	*	φ w /4	32 Hz*3	
1	0	0	0	φ/2	_	244 Hz
1	0	0	1	φ/4	977 Hz	122 Hz
1	0	1	0	φ/8	488 Hz	61 Hz
1	0	1	1	ф/16	244 Hz	30.5 Hz
1	1	0	0	φ/32	122 Hz	_
1	1	0	1	ф/64	61 Hz	_
1	1	1	0	ф/128	30.5 Hz	_
1	1	1	1	ф/256	_	_

*: Don't care

Notes: 1. This is the frame frequency in active (medium-speed, ϕ osc/16) mode when ϕ = 2 MHz.

- 2. When 1/3 duty is selected, the frame frequency is 4/3 times the value shown.
- 3. This is the frame frequency when $\phi w = 32.768 \text{ kHz}$.

13.2.3 LCD Control Register 2 (LCR2)

Bit	7	6	5	4	3	2	1	0
	LCDAB	_	_	_	CDS3	CDS2	CDS1	CDS0
Initial value	0	1	1	0	0	0	0	0
Read/Write	R/W		_	R/W	R/W	R/W	R/W	R/W

LCR2 is an 8-bit read/write register which controls switching between the A waveform and B waveform, and selects the duty cycle of the charge/discharge pulses which control disconnection of the power supply split-resistance from the power supply circuit.

LCR2 is initialized to H'60 upon reset.

Bit 7: A waveform/B waveform switching control (LCDAB)

Bit 7 specifies whether the A waveform or B waveform is used as the LCD drive waveform.

Bit 7 LCDAB	Description	
0	Drive using A waveform	(initial value)
1	Drive using B waveform	

Bits 6 and 5: Reserved bits

Bits 6 and 5 are reserved; they are always read as 1 and cannot be modified.

Bit 4: Reserved bit

Bit 4 is reserved; it is always read as 0 and must not be written with 1.

Bits 3 to 0: Charge/discharge pulse duty cycle select (CDS3 to CDS0)

Bit 3 CDS3	Bit 2 CDS2	Bit 1 CDS1	Bit 0 CDS0	Duty Cycle	Notes	
0	0	0	0	1	Fixed high	(initial value)
0	0	0	1	1/8		
0	0	1	0	2/8		
0	0	1	1	3/8		
0	1	0	0	4/8		
0	1	0	1	5/8		
0	1	1	0	6/8		
0	1	1	1	0	Fixed low	
1	0	*	*	1/16		
1	1	*	*	1/32		

*: Don't care

Bits 3 to 0 select the duty cycle while the power supply split-resistance is connected to the power supply circuit.

When a 0 duty cycle is selected, the power supply split-resistance is permanently disconnected from the power supply circuit, so power should be supplied to pins V_1 , V_2 , and V_3 by an external circuit.

Figure 13.2 shows the waveform of the charge/discharge pulses. The duty cycle is Tc/Tw.

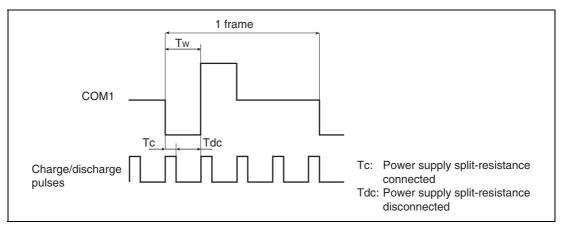


Figure 13.2 Example of A Waveform with 1/2 Duty and 1/2 Bias

13.2.4 Clock Stop Register 2 (CKSTPR2)

Bit	7	6	5	4	3	2	1	0	
	_	_	_	_	AECKSTP	WDCKSTP	PWCKSTP	LDCKSTP	
Initial value	1	1	1	1	1	1	1	1	
Read/Write	_	_	_	_	R/W	R/W	R/W	R/W	

CKSTPR2 is an 8-bit read/write register that performs module standby mode control for peripheral modules. Only the bit relating to the LCD controller/driver is described here. For details of the other bits, see the sections on the relevant modules.

Bit 0: LCD controller/driver module standby mode control (LDCKSTP)

Bit 0 controls setting and clearing of module standby mode for the LCD controller/driver.

Bit 0		
LDCKSTP	Description	
0	LCD controller/driver is set to module standby mode	
1	LCD controller/driver module standby mode is cleared	(initial value)

13.3 Operation

13.3.1 Settings Up to LCD Display

To perform LCD display, the hardware and software related items described below must first be determined.

(1) Hardware Settings

a. Using 1/2 duty

When 1/2 duty is used, interconnect pins V₂ and V₃ as shown in figure 13.3.

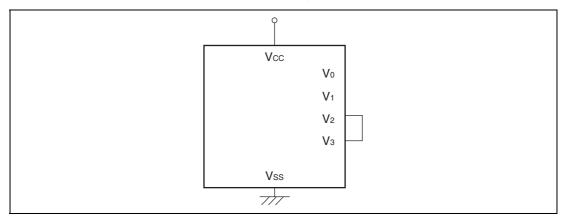


Figure 13.3 Handling of LCD Drive Power Supply when Using 1/2 Duty

b. Large-panel display

As the impedance of the built-in power supply split-resistance is large, it may not be suitable for driving a large panel. If the display lacks sharpness when using a large panel, refer to section 13.3.6, Boosting the LCD Drive Power Supply. When static or 1/2 duty is selected, the common output drive capability can be increased. Set CMX to 1 when selecting the duty cycle. In this mode, with a static duty cycle pins COM₄ to COM₁ output the same waveform, and with 1/2 duty the COM₁ waveform is output from pins COM₂ and COM₁, and the COM₂ waveform is output from pins COM₄ and COM₃.

c. Luminance adjustment function (V₀ pin)

Connecting a resistance between the V_0 and V_1 pins enables the luminance to be adjusted. For details, see section 13.3.3, Luminance Adjustment Function (V_0 Pin).

d. LCD drive power supply setting

With this LSI, there are two ways of providing LCD power: by using the on-chip power supply circuit, or by using an external circuit.

When the on-chip power supply circuit is used for the LCD drive power supply, the V_0 and V_1 pins should be interconnected externally, as shown in figure 13.4 (a).

When an external power supply circuit is used for the LCD drive power supply, connect the external power supply to the V_1 pin, and short the V_0 pin to V_{cc} externally, as shown in figure 13.4 (b).

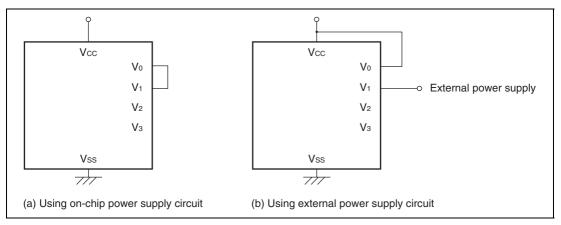


Figure 13.4 Examples of LCD Power Supply Pin Connections

e. Low-power-consumption LCD drive system

Use of a low-power-consumption LCD drive system enables the power consumption required for LCD drive to be optimized. For details, see section 13.3.4, Low-Power-Consumption LCD Drive System.

(2) Software settings

a. Duty selection

Any of four duty cycles—static, 1/2 duty, 1/3 duty, or 1/4 duty—can be selected with bits DTS1 and DTS0.

b. Segment selection

The segment drivers to be used can be selected with bits SGS₃ to SGS₀.

c. Frame frequency selection

The frame frequency can be selected by setting bits CKS₃ to CKS₀. The frame frequency should be selected in accordance with the LCD panel specification. For the clock selection method in watch mode, subactive mode, and subsleep mode, see section 13.3.5, Operation in Power-Down Modes.

d. A or B waveform selection

Either the A or B waveform can be selected as the LCD waveform to be used by means of LCDAB.

13.3.2 Relationship between LCD RAM and Display

The relationship between the LCD RAM and the display segments differs according to the duty cycle. LCD RAM maps for the different duty cycles are shown in figures 13.5 to 13.8.

After setting the registers required for display, data is written to the part corresponding to the duty using the same kind of instruction as for ordinary RAM, and display is started automatically when turned on. Word- or byte-access instructions can be used for RAM setting.

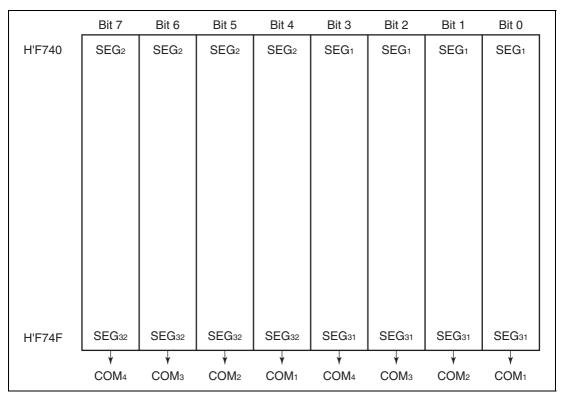


Figure 13.5 LCD RAM Map (1/4 Duty)

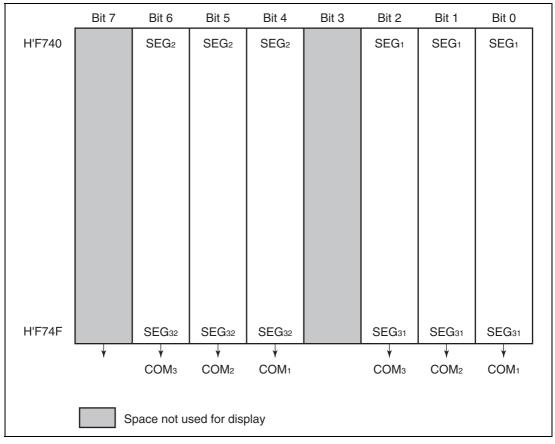


Figure 13.6 LCD RAM Map (1/3 Duty)

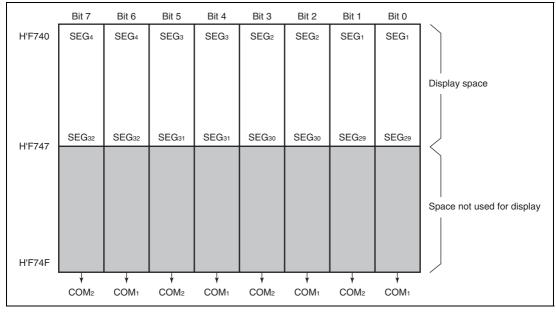


Figure 13.7 LCD RAM Map (1/2 Duty)

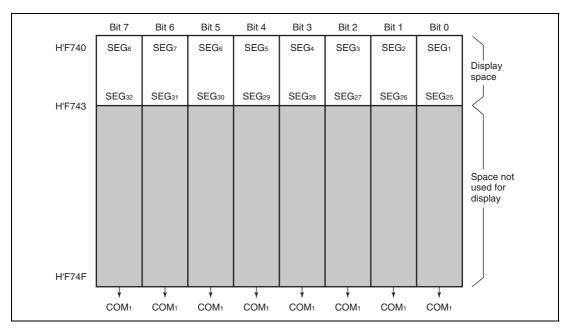


Figure 13.8 LCD RAM Map (Static Mode)

13.3.3 Luminance Adjustment Function (V₀ Pin)

Figure 13.9 shows a detailed block diagram of the LCD drive power supply unit.

The voltage output to the V_0 pin is $V_{\rm cc}$. When either of these voltages is used directly as the LCD drive power supply, the V_0 and V_1 pins should be shorted. Also, connecting a variable resistance, R, between the V_0 and V_1 pins makes it possible to adjust the voltage applied to the V_1 pin, and so to provide luminance adjustment for the LCD panel.

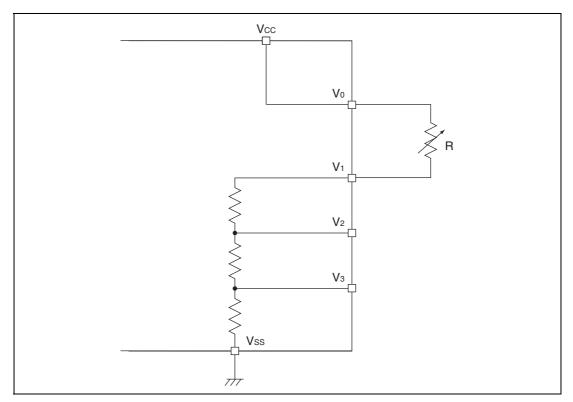


Figure 13.9 LCD Drive Power Supply Unit

13.3.4 Low-Power-Consumption LCD Drive System

The use of the built-in split-resistance is normally the easiest method for implementing the LCD power supply circuit, but since the built-in resistance is fixed, a certain direct current flows constantly from the built-in resistance's V_{cc} to V_{ss} . As this current does not depend on the current dissipation of the LCD panel, if an LCD panel with a small current dissipation is used, a wasteful amount of power will be consumed. This LSI is equipped with a function to minimize this waste of power. Use of this function makes it possible to achieve the optimum power supply circuit for the LCD panel's current dissipation.

(1) Principles

- a. Capacitors are connected as external circuits to LCD power supply pins V1, V2, and V3, as shown in figure 13.10.
- b. The capacitors connected to V1, V2, and V3 are repeatedly charged and discharged in the cycle shown in figure 13.10, maintaining the potentials.
- c. At this time, the charged potential is a potential corresponding to the V1, V2, and V3 pins, respectively. (For example, with 1/3 bias drive, the charge for V2 is 2/3 that of V1, and that for V3 is 1/3 that of V1.)
- d. Power is supplied to the LCD panel by means of the charges accumulated in these capacitors.
- e. The capacitances and charging/discharging periods of these capacitors are therefore determined by the current dissipation of the LCD panel.
- f. The charging and discharging periods can be selected by software.

(2) Example of Operation (with 1/3 bias drive)

- a. During charging period Tc in the figure, the potential is divided among pins V1, V2, and V3 by the built-in split-resistance (the potential of V2 being 2/3 that of V1, and that of V3 being 1/3 that of V1), as shown in figure 13.10, and external capacitors C1, C2, and C3 are charged. The LCD panel is continues to be driven during this time.
- b. In the following discharging period, Tdc, charging is halted and the charge accumulated in each capacitor is discharged, driving the LCD panel.
- c. At this time, a slight voltage drop occurs due to the discharging; optimum values must be selected for the charging period and the capacitor capacitances to ensure that this does not affect the driving of the LCD panel.
- d. In this way, the capacitors connected to V1, V2, and V3 are repeatedly charged and discharged in the cycle shown in figure 13.10, maintaining the potentials and continuously driving the LCD panel.



- e. As can be seen from the above description, the capacitances and charging/discharging periods of the capacitors are determined by the current dissipation of the LCD panel used. The charging/discharging periods can be selected with bits CDS3 to CDS0.
- f. The actual capacitor capacitances and charging/discharging periods must be determined experimentally in accordance with the current dissipation requirements of the LCD panel. An optimum current value can be selected, in contrast to the case in which a direct current flows constantly in the built-in split-resistance.

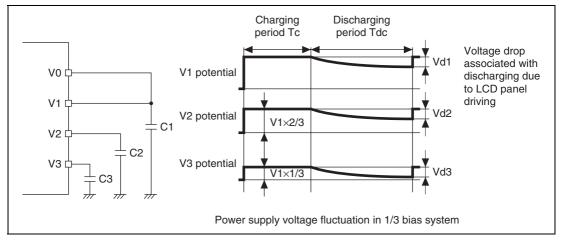


Figure 13.10 Example of Low-Power-Consumption LCD Drive Operation

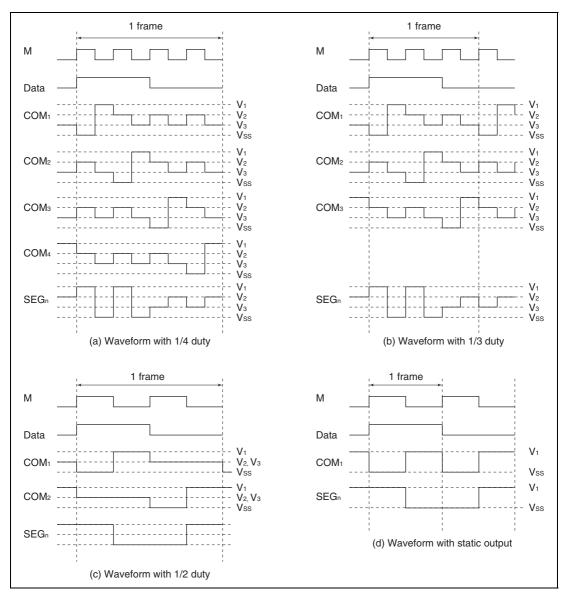


Figure 13.11 Output Waveforms for Each Duty Cycle (A Waveform)

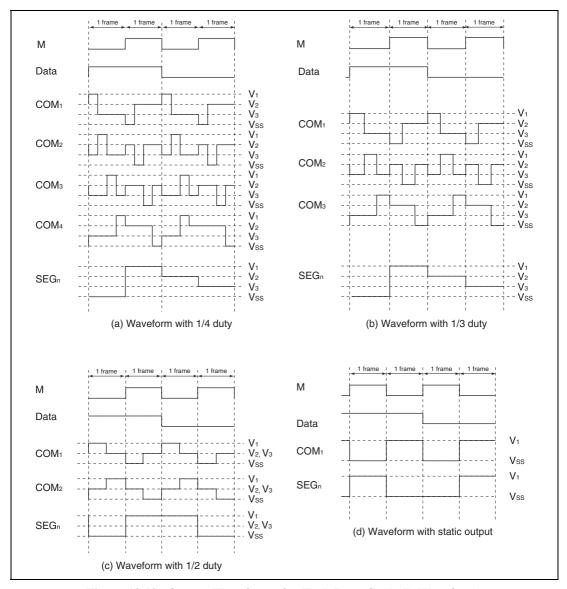


Figure 13.12 Output Waveforms for Each Duty Cycle (B Waveform)

Table 13.3 Output Levels

Data		0	0	1	1	
M		0	1	0	1	
Static	Common output	V ₁	V_{ss}	V ₁	V _{ss}	
	Segment output	V ₁	V _{ss}	V _{ss}	V ₁	
1/2 duty	Common output	V ₂ , V ₃	V ₂ , V ₃	V ₁	V _{ss}	
	Segment output	V ₁	V_{ss}	V _{ss}	V ₁	
1/3 duty	Common output	V ₃	V ₂	V ₁	V _{ss}	
	Segment output	V ₂	V ₃	V _{ss}	V ₁	
1/4 duty	Common output	V ₃	V ₂	V ₁	V _{ss}	
	Segment output	$V_{_2}$	V ₃	V _{ss}	V ₁	

13.3.5 Operation in Power-Down Modes

In this LSI, the LCD controller/driver can be operated even in the power-down modes. The operating state of the LCD controller/driver in the power-down modes is summarized in table 13.4.

In subactive mode, watch mode, and subsleep mode, the system clock oscillator stops, and therefore, unless ϕw , $\phi w/2$, or $\phi w/4$ has been selected by bits CKS3 to CKS0, the clock will not be supplied and display will halt. Since there is a possibility that a direct current will be applied to the LCD panel in this case, it is essential to ensure that ϕw , $\phi w/2$, or $\phi w/4$ is selected. In active (medium-speed) mode, the system clock is switched, and therefore CKS3 to CKS0 must be modified to ensure that the frame frequency does not change.



Table 13.4 Power-Down Modes and Display Operation

Mode		Reset	Active	Sleep	Watch	Sub- active	Sub- sleep	Standby	Module Standby
Clock	ф	Runs	Runs	Runs	Stops	Stops	Stops	Stops	Stops*4
	φw	Runs	Runs	Runs	Runs	Runs	Runs	Stops*1	Stops*4
Display operation	ACT = "0"	Stops	Stops	Stops	Stops	Stops	Stops	Stops*2	Stops
	ACT = "1"	Stops	Functions	Functions	Functions	Functions	Functions	Stops*2	Stops

Notes: 1. The subclock oscillator does not stop, but clock supply is halted.

- 2. The LCD drive power supply is turned off regardless of the setting of the PSW bit.
- 3. Display operation is performed only if ϕw , $\phi w/2$, or $\phi w/4$ is selected as the operating clock.
- 4. The clock supplied to the LCD stops.

13.3.6 Boosting the LCD Drive Power Supply

When a large panel is driven, the on-chip power supply capacity may be insufficient. If the power supply capacity is insufficient when $V_{\rm CC}$ is used as the power supply, the power supply impedance must be reduced. This can be done by connecting bypass capacitors of around 0.1 to 0.3 μF to pins V_{\perp} to V_{\perp} , as shown in figure 13.13, or by adding a split-resistance externally.

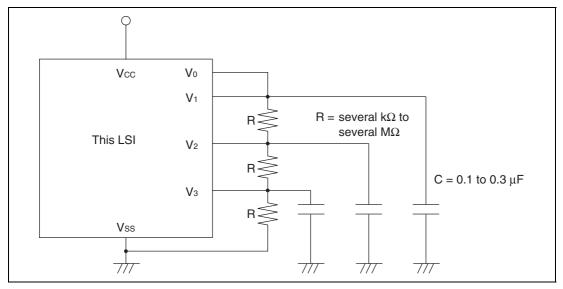


Figure 13.13 Connection of External Split-Resistance

Section 14 Power Supply Circuit

14.1 Overview

This LSI incorporates an internal power supply step-down circuit. Use of this circuit enables the internal power supply to be fixed at a constant level of approximately 3.0 V to 3.2 V, independently of the voltage of the power supply connected to the external V_{cc} pin. As a result, the current consumed when an external power supply is used at 3.0 V or above can be held down to virtually the same low level as when used at approximately 3.0 V. If the external power supply is 3.0 V or below, the internal voltage will be practically the same as the external voltage. It is, of course, also possible to use the same level of external power supply voltage and internal power supply voltage without using the internal power supply step-down circuit.

14.2 When Using Internal Power Supply Step-Down Circuit

Connect the external power supply to the V_{cc} pin, and connect a capacitance of approximately 0.33 μF between CV_{cc} and V_{ss} , as shown in figure 14.1. The internal step-down circuit is made effective simply by adding this external circuit. In the external circuit interface, the external power supply voltage connected to V_{cc} and the GND potential connected to V_{ss} are the reference levels. For example, for port input/output levels, the V_{cc} level is the reference for the high level, and the V_{ss} level is that for the low level. The LCD power supply and A/D converter analog power supply are not affected by the internal step-down current.

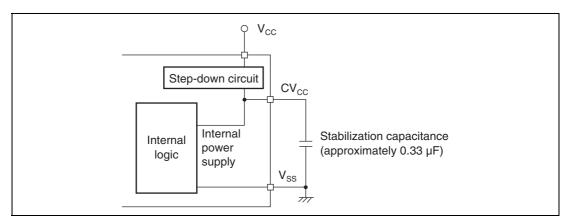


Figure 14.1 Power Supply Connection when Internal Step-Down Circuit is Used

14.3 When Not Using Internal Power Supply Step-Down Circuit

When the internal power supply step-down circuit is not used, connect the external power supply to the CV_{cc} pin and V_{cc} pin, as shown in figure 14.2. The external power supply is then input directly to the internal power supply. The permissible range for the power supply voltage is 2.7 V to 3.6 V. Normally, however, the internal power supply step-down circuit should be used. Operation cannot be guaranteed if a voltage outside this range is input.

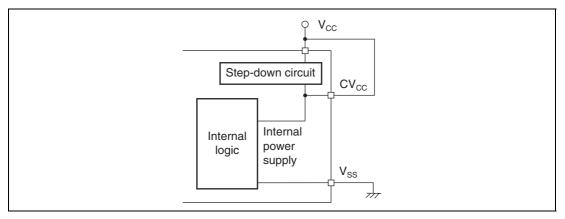


Figure 14.2 Power Supply Connection when Internal Step-Down Circuit is Not Used

Section 15 List of Registers

The register list gives information on the on-chip I/O register addresses, how the register bits are configured, and the register states in each operating mode. The information is given as shown below.

- 1. Register addresses (address order)
- Registers are listed from the lower allocation addresses.
- Registers are classified by functional modules.
- The data bus width is indicated.
- The number of access states is indicated.
- 2. Register bits
- Bit configurations of the registers are described in the same order as the register addresses.
- Reserved bits are indicated by in the bit name column.
- When registers consist of 16 bits, bits are described from the MSB side.
- 3. Register states in each operating mode
- Register states are described in the same order as the register addresses.
- The register states described here are for the basic operating modes. If there is a specific reset for an on-chip peripheral module, refer to the section on that on-chip peripheral module.

15.1 Register Addresses (Address Order)

The data bus width indicates the numbers of bits by which the register is accessed.

The number of access states indicates the number of states based on the specified reference clock.

Register Name	Abbre- viation	Bit No	Address	Module Name	Data Bus Width	Access State
Flash memory control register 1	FLMCR1	8	H'F020	ROM	8	2
Flash memory control register 2	FLMCR2	8	H'F021	ROM	8	2
Flash memory power control register	FLPWCR	8	H'F022	ROM	8	2
Erase block register	EBR	8	H'F023	ROM	8	2
Flash memory enable register	FENR	8	H'F02B	ROM	8	2
Wakeup edge select register	WEGR	8	H'FF90	Interrupts	8	2
Serial port control register	SPCR	8	H'FF91	SCI3	8	2
Subclock output select register	CWOSR	8	H'FF92	Timer A	8	2
Event counter control/status register	ECCSR	8	H'FF95	AEC*1	8	2
Event counter H	ECH	8	H'FF96	AEC*1	8	2
Event counter L	ECL	8	H'FF97	AEC*1	8	2
Serial mode register 31	SMR31	8	H'FF98	SCI31	8	3
Bit rate register 31	BRR31	8	H'FF99	SCI31	8	3
Serial control register 31	SCR31	8	H'FF9A	SCI31	8	3
Transmit data register 31	TDR31	8	H'FF9B	SCI31	8	3
Serial status register 31	SSR31	8	H'FF9C	SCI31	8	3
Receive data register 31	RDR31	8	H'FF9D	SCI31	8	3
Serial mode register 32	SMR32	8	H'FFA8	SCI32	8	3
Bit rate register 32	BRR32	8	H'FFA9	SCI32	8	3
Serial control register 32	SCR32	8	H'FFAA	SCI32	8	3
Transmit data register 32	TDR32	8	H'FFAB	SCI32	8	3
Serial status register 32	SSR32	8	H'FFAC	SCI32	8	3
Receive data register 32	RDR32	8	H'FFAD	SCI32	8	3
Timer mode register A	TMA	8	H'FFB0	Timer A	8	2
Timer counter A	TCA	8	H'FFB1	Timer A	8	2

Register Name	Abbre- viation	Bit No	Address	Module Name	Data Bus Width	Access State
Timer control/status register W	TCSRW	8	H'FFB2	WDT*2	8	2
Timer counter W	TCW	8	H'FFB3	WDT*2	8	2
Timer mode register C	TMC	8	H'FFB4	Timer C	8	2
Timer counter C / Timer load register C	TCC/ TLC	8	H'FFB5	Timer C	8	2
Timer control register F	TCRF	8	H'FFB6	Timer F	8	2
Timer control status register F	TCSRF	8	H'FFB7	Timer F	8	2
8-bit timer counter FH	TCFH	8	H'FFB8	Timer F	8	2
8-bit timer counter FL	TCFL	8	H'FFB9	Timer F	8	2
Output compare register FH	OCRFH	8	H'FFBA	Timer F	8	2
Output compare register FL	OCRFL	8	H'FFBB	Timer F	8	2
Timer mode register G	TMG	8	H'FFBC	Timer G	8	2
Input capture register GF	ICRGF	8	H'FFBD	Timer G	8	2
Input capture register GR	ICRGR	8	H'FFBE	Timer G	8	2
LCD port control register	LPCR	8	H'FFC0	LCD*3	8	2
LCD control register	LCR	8	H'FFC1	LCD*3	8	2
LCD control register 2	LCR2	8	H'FFC2	LCD*3	8	2
A/D result register H	ADRRH	8	H'FFC4	A/D converter	8	2
A/D result register L	ADRRL	8	H'FFC5	A/D converter	8	2
A/D mode register	AMR	8	H'FFC6	A/D converter	8	2
A/D start register	ADSR	8	H'FFC7	A/D converter	8	2
Port mode register 1	PMR1	8	H'FFC8	I/O port	8	2
Port mode register 2	PMR2	8	H'FFC9	I/O port	8	2
Port mode register 3	PMR3	8	H'FFCA	I/O port	8	2
Port mode register 5	PMR5	8	H'FFCC	I/O port	8	2
PWM control register	PWCR	8	H'FFD0	14-bit PWM	8	2
PWM data register U	PWDRU	8	H'FFD1	14-bit PWM	8	2
PWM data register L	PWDRL	8	H'FFD2	14-bit PWM	8	2
Port data register 1	PDR1	8	H'FFD4	I/O port	8	2
Port data register 3	PDR3	8	H'FFD6	I/O port	8	2
Port data register 4	PDR4	8	H'FFD7	I/O port	8	2
Port data register 5	PDR5	8	H'FFD8	I/O port	8	2

Register Name	Abbre- viation	Bit No	Address	Module Name	Data Bus Width	Access State
Port data register 6	PDR6	8	H'FFD9	I/O port	8	2
Port data register 7	PDR7	8	H'FFDA	I/O port	8	2
Port data register 8	PDR8	8	H'FFDB	I/O port	8	2
Port data register A	PDRA	8	H'FFDD	I/O port	8	2
Port data register B	PDRB	8	H'FFDE	I/O port	8	2
Port pull-up control register 1	PUCR1	8	H'FFE0	I/O port	8	2
Port pull-up control register 3	PUCR3	8	H'FFE1	I/O port	8	2
Port pull-up control register 5	PUCR5	8	H'FFE2	I/O port	8	2
Port pull-up control register 6	PUCR6	8	H'FFE3	I/O port	8	2
Port control register 1	PCR1	8	H'FFE4	I/O port	8	2
Port control register 3	PCR3	8	H'FFE6	I/O port	8	2
Port control register 4	PCR4	8	H'FFE7	I/O port	8	2
Port control register 5	PCR5	8	H'FFE8	I/O port	8	2
Port control register 6	PCR6	8	H'FFE9	I/O port	8	2
Port control register 7	PCR7	8	H'FFEA	I/O port	8	2
Port control register 8	PCR8	8	H'FFEB	I/O port	8	2
Port control register A	PCRA	8	H'FFED	I/O port	8	2
System control register 1	SYSCR1	8	H'FFF0	SYSTEM	8	2
System control register 2	SYSCR2	8	H'FFF1	SYSTEM	8	2
IRQ edge select register	IEGR	8	H'FFF2	Interrupts	8	2
Interrupt enable register 1	IENR1	8	H'FFF3	Interrupts	8	2
Interrupt enable register 2	IENR2	8	H'FFF4	Interrupts	8	2
Interrupt request register 1	IRR1	8	H'FFF6	Interrupts	8	2
Interrupt request register 2	IRR2	8	H'FFF7	Interrupts	8	2
Wakeup interrupt request register	IWPR	8	H'FFF9	Interrupts	8	2
Clock stop register 1	CKSTPR1	8	H'FFFA	SYSTEM	8	2
Clock stop register 2	CKSTPR2	8	H'FFFB	SYSTEM	8	2

Notes: 1. AEC: Asynchronous event counter

WDT: Watchdog timer
 LCD: LCD controller/driver



15.2 Register Bits

Register bit names of the on-chip peripheral modules are described below.

Register Abbreviation	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Module Name
FLMCR1	_	SWE	ESU	PSU	EV	PV	E	Р	ROM
FLMCR2	FLER	_	_	_	_	_	_	_	_
FLPWCR	PDWND	_	_	_	_	_	_	_	_
EBR	EB7	EB6	EB5	EB4	EB3	EB2	EB1	EB0	_
FENR	FLSHE	_	_	_	_	_	_	_	_
WEGR	WKEGS7	WKEGS6	WKEGS5	WKEGS4	WKEGS3	WKEGS2	WKEGS1	WKEGS0	Interrupts
SPCR	_	_	SPC32	SPC31	SCINV3	SCINV2	SCINV1	SCINV0	SCI3
CWOSR	_	_	_	_	_	_	_	cwos	Timer A
ECCSR	OVH	OVL	_	CH2	CUEH	CUEL	CRCH	CRCL	AEC*1
ECH	ECH7	ECH6	ECH5	ECH4	ECH3	ECH2	ECH1	ECH0	_
ECL	ECL7	ECL6	ECL5	ECL4	ECL3	ECL2	ECL1	ECL0	=
SMR31	COM31	CHR31	PE31	PM31	STOP31	MP31	CKS311	CKS310	SCI31
BRR31	BRR317	BRR316	BRR315	BRR314	BRR313	BRR312	BRR311	BRR310	_
SCR31	TIE31	RIE31	TE31	RE31	MPIE31	TEIE31	CKE311	CKE310	=
TDR31	TDR317	TDR316	TDR315	TDR314	TDR313	TDR312	TDR311	TDR310	_
SSR31	TDRE31	RDRF31	OER31	FER31	PER31	TEND31	MPBR31	MPBT31	_
RDR31	RDR317	RDR316	RDR315	RDR314	RDR313	RDR312	RDR311	RDR310	=
SMR32	COM32	CHR32	PE32	PM32	STOP32	MP32	CKS321	CKS320	SCI32
BRR32	BRR327	BRR326	BRR325	BRR324	BRR323	BRR322	BRR321	BRR320	_
SCR32	TIE32	RIE32	TE32	RE32	MPIE32	TEIE32	CKE321	CKE320	=
TDR32	TDR327	TDR326	TDR325	TDR324	TDR323	TDR322	TDR321	TDR320	=
SSR32	TDRE32	RDRF32	OER32	FER32	PER32	TEND32	MPBR32	MPBT32	=
RDR32	RDR327	RDR326	RDR325	RDR324	RDR323	RDR322	RDR321	RDR320	_
TMA	TMA7	TMA6	TMA5	_	TMA3	TMA2	TMA1	TMA0	Timer A
TCA	TCA7	TCA6	TCA5	TCA4	TCA3	TCA2	TCA1	TCA0	-
TCSRW	B6WI	TCWE	B4WI	TCSRWE	B2WI	WDON	B0WI	WRST	WDT*2
TCW	TCW7	TCW6	TCW5	TCW4	TCW3	TCW2	TCW1	TCW0	-

Register Abbreviation	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Module Name
TMC	TMC7	TMC6	TMC5	_	_	TMC2	TMC1	TMC0	Timer C
TCC/ TLC	TCC7/ TLC7	TCC6/ TLC6	TCC5/ TLC5	TCC4/ TLC4	TCC3/ TLC3	TCC2/ TLC2	TCC1/ TLC1	TCC0/ TLC0	_
TCRF	TOLH	CKSH2	CKSH1	CKSH0	TOLL	CKSL2	CKSL1	CKSL0	Timer F
TCSRF	OVFH	CMFH	OVIEH	CCLRH	OVFL	CMFL	OVIEL	CCLRL	_
TCFH	TCFH7	TCFH6	TCFH5	TCFH4	TCFH3	TCFH2	TCFH1	TCFH0	=
TCFL	TCFL7	TCFL6	TCFL5	TCFL4	TCFL3	TCFL2	TCFL1	TCFL0	_
OCRFH	OCRFH7	OCRFH6	OCRFH5	OCRFH4	OCRFH3	OCRFH2	OCRFH1	OCRFH0	-
OCRFL	OCRFL7	OCRFL6	OCRFL5	OCRFL4	OCRFL3	OCRFL2	OCRFL1	OCRFL0	_
TMG	OVFH	OVFL	OVIE	IIEGS	CCLR1	CCLR0	CKS1	CKS0	Timer G
ICRGF	ICRGF7	ICRGF6	ICRGF5	ICRGF4	ICRGF3	ICRGF2	ICRGF1	ICRGF0	_
ICRGR	ICRGR7	ICRGR6	ICRGR5	ICRGR4	ICRGR3	ICRGR2	ICRGR1	ICRGR0	_
LPCR	DTS1	DTS0	CMX	_	SGS3	SGS2	SGS1	SGS0	LCD*3
LCR	_	PSW	ACT	DISP	CKS3	CKS2	CKS1	CKS0	_
LCR2	LCDAB	_	_	_	CDS3	CDS2	CDS1	CDS0	_
ADRRH	ADR9	ADR8	ADR7	ADR6	ADR5	ADR4	ADR3	ADR2	A/D
ADRRL	ADR1	ADR0	_	_	_	_	_	_	converter
AMR	CKS	TRGE	_	_	CH3	CH2	CH1	CH0	_
ADSR	ADSF	_	_	_	_	_	_	_	_
PMR1	IRQ3	IRQ2	IRQ1	IRQ4	TMIG	TMOFH	TMOFL	TMOW	I/O port
PMR2	EXCL	_	_	_	_	_	_	_	-
PMR3	AEVL	AEVH	WDCKS	NSC	IRQ0	_	UD	PWM	-
PMR5	WKP7	WKP6	WKP5	WKP4	WKP3	WKP2	WKP1	WKP0	_
PWCR	_	_	_	_	_	_	PWCR1	PWCR0	14-bit
PWDRU	_	_	PWDRU5	PWDRU4	PWDRU3	PWDRU2	PWDRU1	PWDRU0	PWM
PWDRL	PWDRL7	PWDRL6	PWDRL5	PWDRL4	PWDRL3	PWDRL2	PWDRL1	PWDRL0	_
PDR1	P17	P16	P15	P14	P13	P12	P11	P10	I/O port
PDR3	P37	P36	P35	P34	P33	P32	P31	P30	_
PDR4	_	_	_	_	P43	P42	P41	P40	_
PDR5	P57	P56	P55	P54	P53	P52	P51	P50	_
PDR6	P67	P66	P65	P64	P63	P62	P61	P60	=
PDR7	P77	P76	P75	P74	P73	P72	P71	P70	=
PDR8	P87	P86	P85	P84	P83	P82	P81	P80	=
PDRA	_	_	_	_	PA3	PA2	PA1	PA0	=



Register Abbreviation	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Module Name
PDRB	PB7	PB6	PB5	PB4	PB3	PB2	PB1	PB0	I/O port
PUCR1	PUCR17	PUCR16	PUCR15	PUCR14	PUCR13	PUCR12	PUCR11	PUCR10	- '
PUCR3	PUCR37	PUCR36	PUCR35	PUCR34	PUCR33	PUCR32	PUCR31	PUCR30	_'
PUCR5	PUCR57	PUCR56	PUCR55	PUCR54	PUCR53	PUCR52	PUCR51	PUCR50	_'
PUCR6	PUCR67	PUCR66	PUCR65	PUCR64	PUCR63	PUCR62	PUCR61	PUCR60	-
PCR1	PCR17	PCR16	PCR15	PCR14	PCR13	PCR12	PCR11	PCR10	_'
PCR3	PCR37	PCR36	PCR35	PCR34	PCR33	PCR32	PCR31	PCR30	-
PCR4	_	_	_	_	_	PCR42	PCR41	PCR40	- '
PCR5	PCR57	PCR56	PCR55	PCR54	PCR53	PCR52	PCR51	PCR50	- '
PCR6	PCR67	PCR66	PCR65	PCR64	PCR63	PCR62	PCR61	PCR60	- '
PCR7	PCR77	PCR76	PCR75	PCR74	PCR73	PCR72	PCR71	PCR70	_
PCR8	PCR87	PCR86	PCR85	PCR84	PCR83	PCR82	PCR81	PCR80	- '
PCRA	_	_	_	_	PCRA3	PCRA2	PCRA1	PCRA0	- '
SYSCR1	SSBY	STS2	STS1	STS0	LSON	_	MA1	MA0	SYSTEM
SYSCR2	_	_	_	NESEL	DTON	MSON	SA1	SA0	- '
IEGR	_	_	_	IEG4	IEG3	IEG2	IEG1	IEG0	Interrupts
IENR1	IENTA	_	IENWP	IEN4	IEN3	IEN2	IEN1	IEN0	_
IENR2	IENDT	IENAD	_	IENTG	IENTFH	IENTFL	IENTC	IENEC	- '
IRR1	IRRTA	_	_	IRRI4	IRRI3	IRRI2	IRRI1	IRRI0	_
IRR2	IRRDT	IRRAD	_	IRRTG	IRRTFH	IRRTFL	IRRTC	IRREC	-
IWPR	IWPF7	IWPF6	IWPF5	IWPF4	IWPF3	IWPF2	IWPF1	IWPF0	_
CKSTPR1	_	S31CKSTP	S32CKSTP	ADCKSTP	TGCKSTP	TFCKSTP	TCCKSTP	TACKSTP	SYSTEM
CKSTPR2	_	_	_	_	AECKSTP	WDCKSTP	PWCKSTP	LDCKSTP	-

Notes: 1. AEC: Asynchronous event counter

WDT: Watchdog timer
 LCD: LCD controller/driver

15.3 Register States in Each Operating Mode

Register Abbreviation	Reset	Active	Sleep	Watch	Subactive	Subsleep	Standby	Module
FLMCR1	Initialized	_	_	Initialized	Initialized	Initialized	Initialized	ROM
FLMCR2	Initialized	_	_	_	_	_	_	- '
FLPWCR	Initialized	_	_	_	_	_	_	- '
EBR	Initialized	_	_	Initialized	Initialized	Initialized	Initialized	- '
FENR	Initialized	_	_	_	_	_	_	_'
WEGR	Initialized	_	_	_	_	_	_	Interrupts
SPCR	Initialized	_	_	_	_	_	_	SCI3
CWOSR	Initialized	_	_	_	_	_	_	Timer A
ECCSR	Initialized	_	_	_	_	_	_	AEC*1
ECH	Initialized	_	_	_	_	_	_	- '
ECL	Initialized		_	_	_	_	_	_'
SMR31	Initialized	_	_	Initialized	_	_	Initialized	SCI31
BRR31	Initialized	_	_	Initialized	_	_	Initialized	- '
SCR31	Initialized	_	_	Initialized	_	_	Initialized	- '
TDR31	Initialized	_	_	Initialized	_	_	Initialized	- '
SSR31	Initialized	_	_	Initialized	_	_	Initialized	- '
RDR31	Initialized	_	_	Initialized	_	_	Initialized	_'
SMR32	Initialized	_	_	Initialized	_	_	Initialized	SCI32
BRR32	Initialized	_	_	Initialized	_	_	Initialized	_
SCR32	Initialized	_	_	Initialized	_	_	Initialized	- '
TDR32	Initialized	_	_	Initialized	_	_	Initialized	<u>-</u> '
SSR32	Initialized	_	_	Initialized	_	_	Initialized	_
RDR32	Initialized	_	_	Initialized	_	_	Initialized	<u>-</u> '
TMA	Initialized	_	_	_	_	_	_	Timer A
TCA	Initialized	_	_	_	_	_	_	-
TCSRW	Initialized	_	_	_	_	_	_	WDT*2
TCW	Initialized	_	_	_	_	_	_	-



Register Abbreviation	Reset	Active	Sleep	Watch	Subactive	Subsleep	Standby	Module
TMC	Initialized	_	_	_	_	_	_	Timer C
TCC	Initialized	_	_	_	_	_	_	_
TLC	Initialized	_	_	_	_	_	_	-
TCRF	Initialized	_	_	_	_	_	_	Timer F
TCSRF	Initialized	_	_	_	_	_	_	-
TCFH	Initialized	_	_	_	_	_	_	=
TCFL	Initialized	_	_	_	_	_	_	=
OCRFH	Initialized	_	_	_	_	_	_	-
OCRFL	Initialized	_	_	_	_	_	_	=
TMG	Initialized	_	_	_	_	_	_	Timer G
ICRGF	Initialized	_	_	_	_	_	_	_
ICRGR	Initialized	_	_	_	_	_	_	=
LPCR	Initialized	_	_	_	_	_	_	LCD*3
LCR	Initialized	_	_	_	_	_	_	-
LCR2	Initialized	_	_	_	_	_	_	=
ADRRH	_	_	_	_	_	_	_	A/D
ADRRL	_	_	_		_	_	_	converter
AMR	Initialized	_	_	_	_	_	_	=
ADSR	Initialized	_	_	Initialized	Initialized	Initialized	Initialized	=
PMR1	Initialized	_	_	_	_	_	_	I/O port
PMR2	Initialized	_	_	_	_	_	_	_
PMR3	Initialized	_	_	_	_	_	_	_
PMR5	Initialized	_	_	_	_	_	_	=
PWCR	Initialized	_	_	_	_	_	_	14-bit
PWDRU	Initialized	_	_	_	_	_	_	PWM
PWDRL	Initialized	_	_	_	_	_	_	-
PDR1	Initialized	_	_	_	_	_	_	I/O port
PDR3	Initialized	_	_	_	_	_	_	-
PDR4	Initialized	_	_	_	_	_	_	_
PDR5	Initialized	_	_	_	_	_	_	_
PDR6	Initialized	_	_	_	_	_	_	_
PDR7	Initialized	_	_	_	_	_	_	=
PDR8	Initialized	_	_	_	_	_	_	_
PDRA	Initialized	_	_	_	_	_	_	_

Register Abbreviation	Reset	Active	Sleep	Watch	Subactive	Subsleep	Standby	Module
PDRB	Initialized	_	_	_	_	_	_	I/O port
PUCR1	Initialized	_	_	_	_	_	_	_
PUCR3	Initialized	_	_	_	_	_	_	_
PUCR5	Initialized	_	_	_	_	_	_	_
PUCR6	Initialized	_	_	_	_	_	_	_
PCR1	Initialized	_	_	_	_	_	_	_
PCR3	Initialized	_	_	_	_	_	_	_
PCR4	Initialized	_	_	_	_	_	_	_
PCR5	Initialized	_	_	_	_	_	_	_
PCR6	Initialized	_	_	_	_	_	_	_
PCR7	Initialized	_	_	_	_	_	_	_
PCR8	Initialized	_	_	_	_	_	_	_
PCRA	Initialized	_	_	_	_	_	_	_
SYSCR1	Initialized	_	_	_	_	_	_	SYSTEM
SYSCR2	Initialized	_	_	_	_	_	_	_
IEGR	Initialized	_	_	_	_	_	_	Interrupts
IENR1	Initialized	_	_	_	_	_	_	=
IENR2	Initialized	_	_	_	_	_	_	_
IRR1	Initialized	_	_	_	_	_	_	_
IRR2	Initialized	_	_	_	_	_	_	_
IWPR	Initialized	_	_	_	_	_	_	_
CKSTPR1	Initialized	_	_	_	_	_	_	SYSTEM
CKSTPR2	Initialized	_		_	_	_	_	_

Notes: — is not initialized

1. AEC: Asynchronous event counter

2. WDT: Watchdog timer

3. LCD: LCD controller/driver



Section 16 Electrical Characteristics

16.1 Absolute Maximum Ratings

Table 16.1 lists the absolute maximum ratings.

Table 16.1 Absolute Maximum Ratings

Item		Symbol	Value	Unit	Note
Power supply vol	tage	V _{cc}	-0.3 to +7.0	V	*1
		CV _{cc}	-0.3 to +4.3	V	_
Analog power su	pply voltage	AV _{cc}	-0.3 to +7.0	V	
Input voltage	Other than port B	V _{in}	-0.3 to V _{cc} +0.3	V	
	Port B	AV_{in}	-0.3 to AV _{cc} +0.3	V	
Operating temper	rature	T _{opr}	-20 to +75*2 (regular specifications)	°C	_
			-40 to +85*2 (wide-range temperature specifications)	_	
Storage temperat	ture	T _{stg}	-55 to +125	°C	

Notes: 1. Permanent damage may result if maximum ratings are exceeded. Normal operation should be under the conditions specified in Electrical Characteristics. Exceeding these values can result in incorrect operation and reduced reliability.

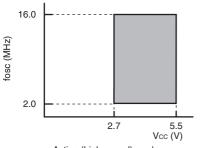
2. The operating temperature ranges from -20°C to $+75^{\circ}\text{C}$ when programming or erasing the flash memory.

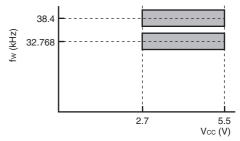
16.2 Electrical Characteristics

16.2.1 Power Supply Voltage and Operating Ranges

The power supply voltage and operating ranges (shaded portions) are shown below.

(1) Power Supply Voltage and Oscillation Frequency Range



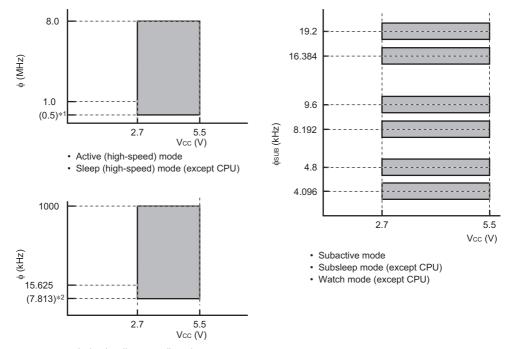


- · Active (high-speed) mode
- Sleep (high-speed) mode

· All operating modes

Note: fosc is the oscillator frequency. When an external clock is used 1 MHz is the minimum fosc value.

(2) Power Supply Voltage and Operating Frequency Range

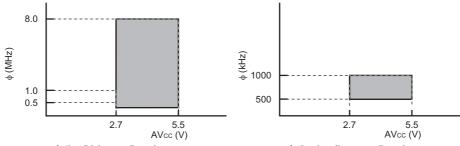


- · Active (medium-speed) mode
- Sleep (medium-speed) mode (except A/D converter)

Notes 1. The figure in parentheses () indicates the minimum operating frequency when an external clock is used. When the resonator is used the minimum operating frequency (φ) is 1 MHz.

2. The figure in parentheses () indicates the minimum operating frequency when an external clock is used. When the resonator is used the minimum operating frequency (\$\phi\$) is 15.625 kHz.

(3) Analog Power Supply Voltage and A/D Converter Operating Range



- · Active (high-speed) mode
- · Sleep (high-speed) mode

- Active (medium-speed) mode
- · Sleep (medium-speed) mode

16.2.2 DC Characteristics

Table 16.2 lists the DC characteristics.

Table 16.2 DC Characteristics

 V_{cc} = 2.7 V to 5.5 V, AV_{cc} = 2.7 V to 5.5 V, V_{ss} = AV_{ss} = 0.0 V, unless otherwise specified

				Valu	es			
Item	Symbol	Applicable Pins	Min	Тур	Max	Unit	Test Condition	Notes
Input high voltage	V _{IH}	RES, WKP ₀ to WKP ₇ , IRQ ₀ to IRQ ₄ , AEVL, AEVH,	V _{cc} × 0.8	_	V _{cc} + 0.3	V	$V_{cc} = 4.0 \text{ V to } 5.5 \text{ V}$	
		TMIC, TMIF, TMIG, ADTRG, SCK ₃₁ , SCK ₃₂	V _{cc} ×0.9	_	V _{cc} + 0.3	_	Other than above	_
		RXD ₃₁ , RXD ₃₂ ,	$V_{\rm CC} \times 0.7$	_	V _{cc} + 0.3	V	$V_{cc} = 4.0 \text{ V to } 5.5 \text{ V}$	
		UD	$V_{cc} \times 0.8$	_	V _{cc} + 0.3		Other than above	_
		OSC ₁	$V_{cc} \times 0.8$	_	V _{cc} + 0.3	V	$V_{cc} = 4.0 \text{ V to } 5.5 \text{ V}$	
			$V_{cc} \times 0.9$	_	V _{cc} + 0.3	_	Other than above	_
		P1 ₀ to P1 ₇ , P3 ₀ to P3 ₇ , P4 ₀ to P4 ₃ , P5 ₀ to P5 ₂ ,	V _{cc} × 0.7	_	V _{cc} + 0.3	V	$V_{cc} = 4.0 \text{ V to } 5.5 \text{ V}$	
		P6 ₀ to P6 ₇ , P7 ₀ to P7 ₇ , P8 ₀ to P8 ₇ , PA ₀ to PA ₃	V _{cc} × 0.8	_	V _{cc} + 0.3	_	Other than above	_
		PB _o to PB ₇	$V_{cc} \times 0.7$	_	AV _{CC} + 0.3	٧	$V_{cc} = 4.0 \text{ V to } 5.5 \text{ V}$	
			$V_{cc} \times 0.8$	_	AV _{CC} + 0.3	_	Other than above	_
		EXCL	$V_{cc} \times 0.9$	_	V _{cc} + 0.3	٧		

Note: Connect the TEST pin to $V_{\rm ss}$.

				Valu	es			
Item	Symbol	Applicable Pins	Min	Тур	Max	Unit	Test Condition	Notes
Input low voltage	V _{IL}	RES, WKP, to WKP, IRQ, to IRQ, AEVL, AEVH,	- 0.3	_	$V_{cc} \times 0.2$	V	$V_{cc} = 4.0 \text{ V to } 5.5 \text{ V}$	
		TMIC, TMIF, TMIG, ADTRG, SCK ₃₁ , SCK ₃₂	- 0.3	_	V _{cc} × 0.1		Other than above	
		RXD ₃₁ , RXD ₃₂ ,	- 0.3	_	$V_{\text{CC}} \times 0.3$	V	$V_{cc} = 4.0 \text{ V to } 5.5 \text{ V}$	_
		UD	- 0.3	_	$V_{\rm CC} \times 0.2$		Other than above	
		OSC,	- 0.3	_	$V_{\rm CC} \times 0.2$	V	$V_{cc} = 4.0 \text{ V to } 5.5 \text{ V}$	_
			- 0.3	_	$V_{\text{CC}} \times 0.1$		Other than above	
		EXCL	- 0.3	_	0.1 V _{cc}	V		
		P1 ₀ to P1 ₇ P3 ₀ to P3 ₇ , P4 ₀ to P4 ₃ , P5 ₀ to P5 ₇ , P6 ₀ to P6 ₇ ,	- 0.3	_	$V_{cc} \times 0.3$	V	V_{cc} = 4.0 V to 5.5 V	_
		P7 ₀ to P7 ₇ , P8 ₀ to P8 ₇ , PA ₀ to PA ₃ , PB ₀ to PB ₇	- 0.3	_	$V_{cc} \times 0.2$		Other than above	
Output high voltage	V _{OH}	P1 ₀ to P1 ₇ P3 ₀ to P3 ₇ , P4 ₀ to P4 ₂ ,	V _{cc} - 1.0	_	_	V	$V_{CC} = 4.0 \text{ V to } 5.5 \text{ V}$ $-I_{OH} = 1.0 \text{ mA}$	
vollage		P5 ₀ to P5 ₇ , P6 ₀ to P6 ₇ ,	V _{cc} - 0.5	_	_	_	$V_{cc} = 4.0 \text{ V to } 5.5 \text{ V}$ $-I_{OH} = 0.5 \text{ mA}$	_
		$P7_0$ to $P7_7$, $P8_0$ to $P8_7$, PA_0 to PA_3	V _{cc} - 0.3	_	_		$-I_{OH} = 0.1 \text{ mA}$	

				Value	s			
Item	Symbol	Applicable Pins	Min	Тур	Max	Unit	Test Condition	Notes
Output low voltage	V _{OL}	P1 ₀ to P1 ₇ , P4 ₀ to P4 ₂ , P5 ₀ to P5 ₇ ,	_	_	0.6	V	$V_{cc} = 4.0 \text{ V to } 5.5 \text{ V}$ $I_{OL} = 1.6 \text{ mA}$	
		P6 ₀ to P6 ₇ , P7 ₀ to P7 ₇ , P8 ₀ to P8 ₇ , PA ₀ to PA ₃	_	_	0.5		I _{OL} = 0.4 mA	_
		P3 ₀ to P3 ₇	_	_	1.0		$V_{cc} = 4.0 \text{ V to } 5.5 \text{ V}$ $I_{oL} = 10 \text{ mA}$	_
			_	_	0.6		$V_{cc} = 4.0 \text{ V to } 5.5 \text{ V}$ $I_{oL} = 1.6 \text{ mA}$	_
			_	_	0.5		$I_{OL} = 0.4 \text{ mA}$	_
Input/ output leakage current	I _{IL}	RES, P4 ₃ , P1 ₀ to P1 ₇ , OSC ₁ , X ₁ , P3 ₀ to P3 ₇ , P4 ₀ to P4 ₂ , P5 ₀ to P5 ₇ , P6 ₀ to P6 ₇ , P7 ₀ to P7 ₇ , P8 ₀ to P8 ₇ , PA ₀ to PA ₃	_	_	1.0	μΑ	$V_{IN} = 0.5 \text{ V to } V_{CC} - 0.5 \text{ V}$	
		PB ₀ to PB ₇	_	_	1.0		$V_{IN} = 0.5 \text{ V to AV}_{CC} - 0.5 \text{ V}$	_

Pull-up – MOS current	-I _p	P1 ₀ to P1 ₇ , P3 ₀ to P3 ₇ , P5 ₀ to P5 ₇ ,	20	_	200	μА	$V_{CC} = 5.0 \text{ V},$ $V_{IN} = 0.0 \text{ V}$	
		P6 ₀ to P6 ₇	_	40	_	<u> </u>	$V_{CC} = 2.7 \text{ V},$ $V_{IN} = 0.0 \text{ V}$	Refer- ence value
Input capaci- tance	C _{in}	All input pins except power supply pin	_	_	15.0	pF	f = 1 MHz, $V_{IN} = 0.0 \text{ V},$ $T_a = 25^{\circ}\text{C}$	



			Value	s			
Symbol	Applicable Pins	Min	Тур	Max	Unit	Test Condition	Notes
I _{OPE1}	V _{cc}	_	0.8	_	mA	Active (high-speed) mode $V_{cc} = 2.7 \text{ V},$ $f_{osc} = 2 \text{ MHz}$	*1 *3 *4 Approx. max. value = 1.1 × Typ.
		_	1.2	_			*2 *3 *4
							Approx. max. value = 1.1 × Typ.
		_	1.0	_	_	Active (high-speed) mode V _{CC} = 5 V, f _{OSC} = 2 MHz	*1 *3 *4 Approx. max. value = 1.1 × Typ.
		_	1.5	_	_		*2 *3 *4
							Approx. max. value = 1.1 × Typ.
		_	2.0	_	_	Active (high-speed)	*1 *3 *4
						mode $V_{CC} = 5 \text{ V},$ $f_{OSC} = 4 \text{ MHz}$	Approx. max. value = 1.1 × Typ.
		_	2.4	_	_		*2 *3 *4
		_	4.0	7.0	_	Active (high-speed)	*1 *3 *4
		_	4.9	7.0		mode V _{cc} = 5 V, f _{ees} = 10 MHz	*2 *3 *4
			* ''	Symbol Applicable Pins Min Typ I _{OPE1} V _{CC} — 0.8 — 1.2 — 1.0 — 1.5 — 2.0 — 4.0	Topes Voc	Name	Symbol Applicable Pins Min Typ Max Unit Test Condition I _{OPE1} V _{CC} — 0.8 — mA Active (high-speed) mode v _{CC} = 2.7 V, f _{OSC} = 2 MHz — 1.2 — — Active (high-speed) mode v _{CC} = 5 V, f _{OSC} = 2 MHz — 1.5 — Active (high-speed) mode v _{CC} = 5 V, f _{OSC} = 4 MHz — 2.4 — Active (high-speed) mode mode mode v _{CC} = 5 V, f _{OSC} = 4 MHz

				Value	s			
Item	Symbol	Applicable Pins	Min	Тур	Max	Unit	Test Condition	Notes
Active I, mode current consumption	I _{OPE2}	V _{cc}	_	0.4	_	mA	Active (medium- speed) mode $V_{cc} = 2.7 \text{ V},$ $f_{osc} = 2 \text{ MHz},$ $\phi_{osc}/128$	*1 *3 *4 Approx. max. value = 1.1 × Typ.
			_	0.7	_			*2 *3 *4
								Approx. max. value = 1.1 × Typ.
				0.5	_	_	Active (medium- speed) mode $V_{\rm CC} = 5 \text{ V},$ $f_{\rm OSC} = 2 \text{ MHz},$ $\phi_{\rm OSC}/128$ $Active \text{ (medium-speed) mode}$ $V_{\rm CC} = 5 \text{ V},$ $f_{\rm OSC} = 4 \text{ MHz},$ $\phi_{\rm OSC}/128$	*1*3*4 Approx. max. value = 1.1 × Typ.
			_	1.0	_			*2*3*4 Approx. max. value = 1.1 × Typ.
			_	0.8	_	_		*1 *3 *4 Approx. max. value = 1.1 × Typ.
			_	1.2	_			*2 *3 *4
			_	1.2	3.0	_	Active (medium-	*1 *3 *4
				_	1.7	3.0	_	speed) mode $V_{cc} = 5 \text{ V},$ $f_{osc} = 10 \text{ MHz},$ $\phi_{osc}/128$

	Symbol			Value	s			
Item		Applicable Pins	Min	Тур	Max	Unit	Test Condition	Notes
Sleep mode current consump- tion	I _{SLEEP}	V _{cc}	_	0.5	_	mA	$V_{CC} = 2.7 \text{ V},$ $f_{OSC} = 2 \text{ MHz}$	*1*3*4 Approx. max. value = 1.1 × Typ.
				8.0	_			*2 *3 *4
								Approx. max. value = 1.1 × Typ.
			_	0.7	_		$V_{CC} = 5 \text{ V},$ $f_{OSC} = 2 \text{ MHz}$	*1 *3 *4 Approx. max. value = 1.1 × Typ.
			_	1.2	_	_		*2 *3 *4
								Approx. max. value = 1.1 × Typ.
			_	1.1	_		$V_{\rm CC} = 5 \text{ V},$ $f_{\rm OSC} = 4 \text{ MHz}$	*1 *3 *4 Approx.
							osc	max. value = 1.1 × Typ.
			_	1.6	_			*2 *3 *4
				1.9	5.0		$V_{CC} = 5 \text{ V},$	*1 *3 *4
			_	2.6	5.0		$f_{OSC} = 10 \text{ MHz}$	*2 *3 *4
Subactive mode current	I _{SUB}	V _{cc}	_	12	_	μΑ	V _{cc} = 2.7 V, LCD on, 32-kHz crystal	*1 *3 *4 Reference value
consump- tion			_	15	_		resonator used	*2 *3 *4
แบบ							$(\phi_{SUB} = \phi_{W}/8)$	Reference value
			_	18	50		V _{CC} = 2.7 V,	*1 *3 *4
			_	30	50		LCD on, 32-kHz crystal resonator used $(\phi_{SUB} = \phi_W/2)$	*2 *3 *4

				Value	s			
Item	Symbol	Applicable Pins	Min	Тур	Max	Unit	Test Condition	Notes
Subsleep mode current consump- tion	SUBSP	V _{cc}	_	3.8	16	μА	$V_{CC} = 2.7 \text{ V},$ LCD on, 32-kHz crystal resonator used $(\phi_{SUB} = \phi_W/2)$	*3 *4
Watch I _{WATC} mode current consumption	I _{watch}	V _{cc}	_	1.8	_	μΑ	$V_{CC} = 2.7 \text{ V},$ $T_a = 25^{\circ}\text{C},$ 32-kHz crystal	*1 *3 *4 Reference value
			_	1.8	_		resonator used,	*2 *3 *4
uon							LCD not used	Reference value
			_	3.0	6.0	_	V _{cc} = 2.7 V, 32-kHz crystal resonator used, LCD not used	*3 *4
Standby mode current consump-	I _{STBY}	, V _{cc}	_	0.3	_	μΑ	$V_{CC} = 2.7 \text{ V},$ $T_a = 25^{\circ}\text{C},$ 32-kHz crystal resonator not used	*1 *3 *4 Reference value
tion			_	0.3	_		$V_{\rm CC}$ = 2.7 V, $T_{\rm a}$ = 25°C, 32-kHz crystal resonator not used	*2 *3 *4 Reference value
			_	0.4	_	_	$V_{cc} = 5.0 \text{ V},$ $T_a = 25^{\circ}\text{C},$ 32-kHz crystal	*1 *3 *4 Reference value
			_	0.5	_		resonator not used	*2 *3 *4
								Reference value
			_	1.0	5.0	_	32-kHz crystal resonator not used	*3 *4
RAM data retaining voltage	V _{RAM}	V _{cc}	2.0	_	_	V		*5
Allowable output low	I _{OL}	Output pins except port 3	_	_	2.0	mA	$V_{cc} = 4.0 \text{ V to } 5.5 \text{ V}$	
current (per pin)		Port 3	_	_	10.0		$V_{cc} = 4.0 \text{ V to } 5.5 \text{ V}$	_
(boi biii)		All output pins		_	0.5			_

				Value	s			
Item	Symbol	Applicable Pins	Min	Тур	Max	Unit	Test Condition	Notes
Allowable output low	ΣI_{OL}	Output pins except port 3	_	_	40.0	mA	$V_{cc} = 4.0 \text{ V to } 5.5 \text{ V}$	_
current (total)		Port 3	_	_	80.0		$V_{cc} = 4.0 \text{ V to } 5.5 \text{ V}$	_
		All output pins	_	_	20.0			_
Allowable	-I _{OH}	All output pins	_	_	2.0	mA	V _{cc} = 4.0 V to 5.5 V	
output high current (per pin)			_	_	0.2		Other than above	
Allowable	Σ -I _{OH}	All output pins	_	_	15.0	mA	$V_{cc} = 4.0 \text{ V to } 5.5 \text{ V}$	
output high current (total)			_	_	10.0		Other than above	
V _{cc} start voltage	V _{CCSTART}	V _{cc}	0	_	0.1	V		*2
V _{cc} rising gradient	SV _{cc}	V _{cc}	0.05	_	_	V/ms		*2

Notes: Connect the TEST pin to V_{ss}.

- 1. Applies to the mask-ROM version.
- 2. Applies to the flash memory version.
- 3. Pin states when current consumption is measured

Mode	RES Pin	Internal State	Other Pins	LCD Power	Oscillator Pins
Active (high-speed) mode (I _{OPE1}) Active (medium-speed) mode (I _{OPE2})	V _{cc}	Only CPU operates	V _{cc}	Supply Stops	System clock: crystal resonator Subclock: Pin X ₁ = GND
Sleep mode	V _{cc}	Only all on-chip timers operate	V _{cc}	Stops	
Subactive mode	V _{cc}	Only CPU operates	V _{cc}	Stops	System clock:
Subsleep mode	V _{cc}	Only all on-chip timers operate CPU stops	V _{cc}	Stops	crystal resonator Subclock: crystal resonator
Watch mode	V _{cc}	Only clock time base operates CPU stops	V _{cc}	Stops	_
Standby mode	V _{cc}	CPU and timers both stop	V _{cc}	Stops	System clock: crystal resonator Subclock: Pin X ₁ = GND

- 4. Except current which flows to the pull-up MOS or output buffer
- 5. Voltage maintained in standby mode

16.2.3 AC Characteristics

Table 16.3 lists the control signal timing and table 16.4 lists the serial interface timing.

Table 16.3 Control Signal Timing

 V_{cc} = 2.7 V to 5.5 V, AV_{cc} = 2.7 V to 5.5 V, V_{ss} = AV_{ss} = 0.0 V, unless otherwise specified

		Applicable	Values					Reference
Item	Symbol	Pins	Min	Тур	Max	Unit	Test Condition	Figure
System clock oscillation frequency	f _{osc}	OSC ₁ , OSC ₂	2.0	_	16.0	MHz		
OSC clock (ϕ_{osc}) cycle time	t _{osc}	OSC ₁ , OSC ₂	62.5	_	500 (1000)	ns		Figure 16.1*2
System clock (φ)	$t_{\rm cyc}$		2	_	128	$t_{\rm osc}$	_	
cycle time			_	_	128	μs		
Subclock oscillation frequency	f_w	X ₁ , X ₂ , EXCL	_	32.768 or 38.4	_	kHz		
Watch clock (ϕ_W) cycle time	t _w	X ₁ , X ₂ , EXCL	_	30.5 or 26.0	_	μs		Figure 16.1
Subclock (ϕ_{SUB}) cycle time	t _{subcyc}		2	_	4	t _w		*1
Instruction cycle time			2	_	_	t _{cyc} t _{subcyc}		
Oscillation stabilization time	t _{rc}	OSC ₁ , OSC ₂	_	20	45	μs	Ceramic resonator $(V_{cc} = 3.0 \text{ to } 5.5 \text{ V})$	Figure 16.9
			_	80	_		Ceramic resonator other than above	_
			_	0.8	2	ms	Crystal resonator	_
			_	_	50	_	Other than above	
	t _{rc}	X ₁ , X ₂	_	_	2.0	s		
External clock high width	t _{CPH}	OSC,	25	_	_	ns		Figure 16.1
		EXCL	_	15.26 or 13.02	_	μs		Figure 16.1
External clock low width	t _{CPL}	OSC ₁	25	_	_	ns		Figure 16.1
		EXCL	_	15.26 or 13.02	_	μs		Figure 16.1

		Applicable		Valu	es			Reference
Item	Symbol	Pins	Min	Тур	Max	Unit	Test Condition	Figure
External clock rise time	t _{CPr}	OSC ₁	_	_	6	ns		Figure 16.1
		EXCL	_	_	55.0			Figure 16.1
External clock fall time	t _{CPf}	OSC ₁	_	_	6	ns		Figure 16.1
		EXCL	_	_	55.0			Figure 16.1
RES pin low width	t _{REL}	RES	10	_	_	t _{cyc}		Figure 16.2
Input pin high width	t _{iH}	IRQ₀ to IRQ₄, WKP₀ to WKP7, ADTRG, TMIC, TMIF, TMIG	2	_	_	t _{cyc} t _{subcyc}		Figure 16.3
		AEVL, AEVH	32	_	_	ns	_	
Input pin low width	t _{il.}	IRQ₀ to IRQ₄, WKP₀ to WKP¬, ADTRG, TMIC, TMIF, TMIG	2	_	_	t _{cyc} t _{subcyc}		Figure 16.3
		AEVL, AEVH	32	_	_	ns	_	
UD pin minimum transition width	t _{UDH}	UD	4	_	_	t _{cyc} t _{subcyc}		Figure 16.4

Notes: 1. Determined by the SA1 and SA0 bits in the system control register 2 (SYSCR2).

2. The figure in parentheses () indicates the maximum fosc value when an external clock is used.

Table 16.4 Serial Interface (SCI3) Timing

 $V_{cc} = 2.7 \text{ V}$ to 5.5 V, $AV_{cc} = 2.7 \text{ V}$ to 5.5 V, $V_{ss} = AV_{ss} = 0.0 \text{ V}$, unless otherwise specified

				Values	3		Test	Reference
Item		Symbol	Min	Тур	Max	Unit	Condition	Figure
Input clock	Asynchronous	t _{scyc}	4	_	_	t _{cyc} or t _{subcyc}		Figure 16.5
cycle	Clocked synchronous	=	6	_	_	_		
Input clock	pulse width	t _{sckw}	0.4	_	0.6	t _{scyc}		Figure 16.5
Transmit da (clocked sy	ta delay time nchronous)	t _{TXD}	_	_	1	$\mathbf{t}_{\scriptscriptstyle{\mathrm{cyc}}}$ or $\mathbf{t}_{\scriptscriptstyle{\mathrm{subcyc}}}$		Figure 16.6
Receive dat (clocked sys	ta setup time nchronous)	t _{RXS}	200	_	_	ns		Figure 16.6
Receive data		t _{RXH}	200	_	_	ns		Figure 16.6

16.2.4 A/D Converter Characteristics

Table 16.5 shows the A/D converter characteristics.

Table 16.5 A/D Converter Characteristics

 $V_{cc} = 2.7 \text{ V}$ to 5.5 V, $AV_{cc} = 2.7 \text{ V}$ to 5.5 V, $V_{ss} = AV_{ss} = 0.0 \text{ V}$, unless otherwise specified

		Applicable		Valu	ies		Test	Reference
Item	Symbol	Pins	Min	Тур	Max	Unit	Condition	Figure
Analog power supply voltage	AV _{cc}	AV _{cc}	2.7	_	5.5	٧		*1
Analog input voltage	AV _{IN}	AN _o to AN ₃	- 0.3	_	AV _{cc} + 0.3	V		
Analog power supply	Al _{OPE}	AV _{cc}	_	_	1.5	mA	AV _{CC} = 5.0 V	
current	AI _{STOP1}	AV _{cc}		600	_	μΑ		*2
								Reference value
	Al _{STOP2}	AV _{cc}	_	_	5.0	μΑ		*3
Analog input capacitance	C _{AIN}	AN _o to AN ₇	_	_	15.0	pF		
Allowable signal source impedance	R _{AIN}		_	_	10.0	kΩ		
Resolution (data length)			_	_	10	bit		



		Applicable		Valu	ies		Test	Reference
Item	Symbol	• •	Min	Тур	Max	Unit	Condition	Figure
Nonlinearity error			_	_	±3.5	LSB	AV _{cc} = 4.0 V to 5.5 V	
			_	_	±7.5		AV _{cc} = 2.7 V to 5.5 V	-
Quantization error			_	_	±0.5	LSB		
Absolute accuracy			_	±2.0	±4.0	LSB	AV _{cc} = 4.0 V to 5.5 V	
			_	±2.0	±8.0		AV _{cc} = 2.7 V to 5.5 V	-
Conversion time			7.8	_	124	μs		

Notes: 1. Set $AV_{cc} = V_{cc}$ when the A/D converter is not used.

- 2. AI_{STOP1} is the current in active and sleep modes while the A/D converter is idle.
- 3. Al_{STOP2} is the current at reset and in standby, watch, subactive, and subsleep modes while the A/D converter is idle.

16.2.5 LCD Characteristics

Table 16.6 shows the LCD characteristics.

Table 16.6 LCD Characteristics

 $V_{cc} = 2.7 \text{ V}$ to 5.5 V, $AV_{cc} = 2.7 \text{ V}$ to 5.5 V, $V_{ss} = AV_{ss} = 0.0 \text{ V}$, unless otherwise specified

		Applicable		Value	s			Reference
Item	Symbol	Pins	Min	Тур	Max	Unit	Test Condition	Figure
Segment driver step-down voltage	V _{DS}	SEG, to SEG ₃₂	_	_	0.6	V	$I_D = 2 \mu A$ V1 = 2.7 V to 5.5 V	*1
Common driver step-down voltage	V _{DC}	COM, to COM,	_	_	0.3	V	$I_D = 2 \mu A$ V1 = 2.7 V to 5.5 V	*1
LCD power supply split-resistance	R _{LCD}		1.5	3.0	7.0	ΜΩ	Between V1 and V _{ss}	
Liquid crystal display voltage	V _{LCD}	V ₁	2.7	_	5.5	V		*2

Notes: 1. The voltage step-down from power supply pins V1, V2, V3, and V_{ss} to each segment pin or common pin.

2. When the liquid crystal display voltage is supplied from an external power supply, ensure that the following relationship is maintained: V1 \geq V2 \geq V3 \geq V_{ss}.



16.2.6 Flash Memory Characteristics

Table 16.7 Flash Memory Characteristics

Condition: $AV_{cc} = 2.7 \text{ V}$ to 5.5 V, $V_{ss} = AV_{ss} = 0.0 \text{ V}$, $V_{cc} = 2.7 \text{ V}$ to 5.5 V (range of operating voltage when reading), $V_{cc} = 3.0 \text{ V}$ to 5.5 V (range of operating voltage when programming/erasing), $T_a = -20^{\circ}\text{C}$ to $+75^{\circ}\text{C}$ (range of operating temperature when programming/erasing: product with regular specifications, product with wide-range

temperature specifications)

				Values			Test
Item		Symbol	Min	Тур	Max	Unit	Conditions
Programming t		t _p	_	7	200	ms/128 bytes	
Erase time*1*3*	* 5	t _E	_	100	1200	ms/block	
Reprogrammin	g count	N_{wec}	1000*8	10000*9	_	times	
Data retain per	riod	t_{DRP}	10*10	_	_	year	
Programming	Wait time after SWE-bit setting*1	X	1	_	_	μs	
	Wait time after PSU-bit setting*1	у	50	_	_	μs	
	Wait time after	z1	28	30	32	μs	1 ≤ n ≤ 6
	P-bit setting*1*4	z2	198	200	202	μs	7 ≤ n ≤ 1000
		z3	8	10	12	μs	Additional programming
	Wait time after P-bit clear*1	α	5	_	_	μs	
	Wait time after PSU-bit clear*1	β	5	_	_	μs	
	Wait time after PV-bit setting*1	γ	4	_	_	μs	
	Wait time after dummy write*1	ε	2	_	_	μs	
	Wait time after PV-bit clear*1	η	2	_	_	μs	
	Wait time after SWE-bit clear*1	θ	100	_	_	μs	
	Maximum programming count*1*4*5	N	_	_	1000	times	

				Value	s		Test
Item		Symbol	Min	Тур	Max	Unit	Conditions
Erase	Wait time after SWE-bit setting*1	Х	1	_	_	μs	
	Wait time after ESU-bit setting*1	у	100	_	_	μs	
	Wait time after E-bit setting*1*6	Z	10	_	100	ms	
	Wait time after E-bit clear*1	α	10	_	_	μs	
	Wait time after ESU-bit clear*1	β	10	_	_	μs	
	Wait time after EV-bit setting*1	γ	20	_	_	μs	
	Wait time after dummy write*1	ε	2	_	_	μs	
	Wait time after EV-bit clear*1	η	4	_	_	μs	
	Wait time after SWE-bit clear*1	θ	100	_	_	μs	
	Maximum erase count****	N			120	times	

Notes:

- 1. Set the times according to the program/erase algorithms.
- 2. Programming time per 128 bytes (Shows the total period for which the P bit in FLMCR1 is set. It does not include the programming verification time.)
- 3. Block erase time (Shows the total period for which the E bit in FLMCR1 is set. It does not include the erase verification time.)
- 4. Maximum programming time (t_n (max))
 - t_P (max) = Wait time after P-bit setting (z) × maximum number of writes (N)
- The maximum number of writes (N) should be set according to the actual set value of z1, z2, and z3 to allow programming within the maximum programming time (t_p (max)).

The wait time after P-bit setting (z1 and z2) should be alternated according to the number of writes (n) as follows:

$$1 \le n \le 6$$
 $z1 = 30 \ \mu s$
 $7 \le n \le 1000$ $z2 = 200 \ \mu s$

- 6. Maximum erase time (t_E (max))
 - t_{E} (max) = Wait time after E-bit setting (z) × maximum erase count (N)
- 7. The maximum number of erases (N) should be set according to the actual set value of z to allow erasing within the maximum erase time (t_F (max)).
- 8. This minimum value guarantees all characteristics after reprogramming (the guaranteed range is from 1 to the minimum value).
- Reference value when the temperature is 25°C (normally reprogramming will be performed by this
 count).
- 10. This is a data retain characteristic when reprogramming is performed within the specification range including this minimum value.



16.3 Operation Timing

Figures 16.1 to 16.6 show timing diagrams.

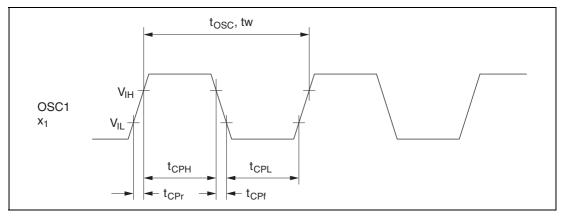


Figure 16.1 Clock Input Timing

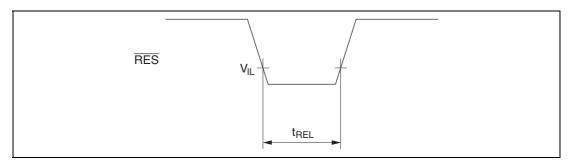


Figure 16.2 RES Low Width

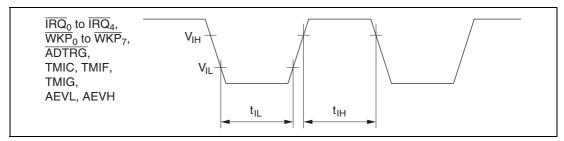


Figure 16.3 Input Timing

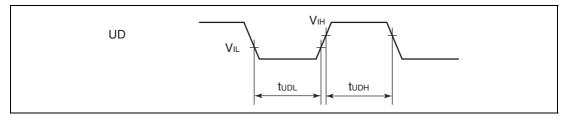


Figure 16.4 UD Pin Minimum Modulation Width Timing

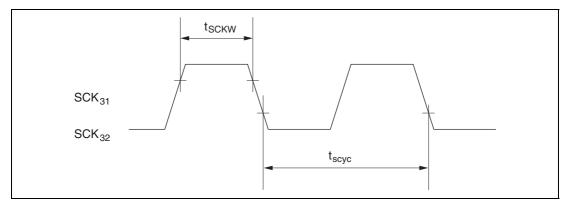


Figure 16.5 SCK3 Input Clock Timing

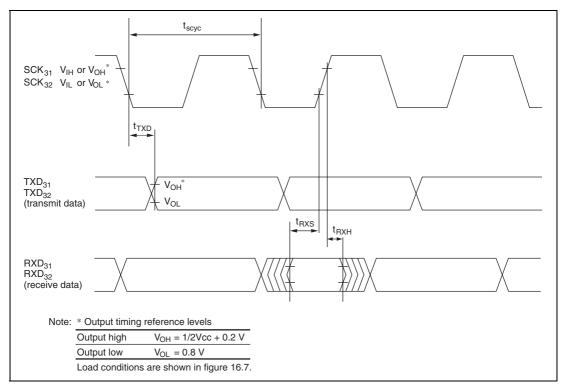


Figure 16.6 SCI3 Synchronous Mode Input/Output Timing

16.4 Output Load Circuit

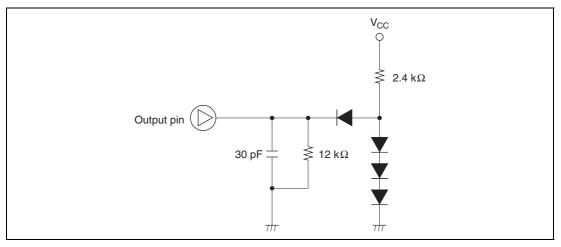


Figure 16.7 Output Load Condition

16.5 Resonator

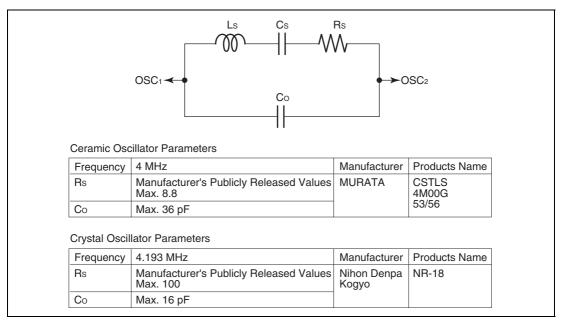


Figure 16.8 Resonator Equivalent Circuit

Resonating Frequency	Manufacturer	Model	C ₁ , C ₂
4 MHz	Nihon Denpa Kogyo	NR-18	12pF ± 20%
10 MHz			

Ceramic resonator

Resonating Frequency	Manufacturer	Model	C ₁ , C ₂
2 MHz	MURATA	CSTCC2M00G53-B0	15pF ± 20%
		CSTCC2M00G56-B0	47pF ± 20%
4 MHz		CSTLS4M00G53-B0	15pF ± 20%
		CSTLS4M00G56-B0	47pF ± 20%
10 MHz		CSTLS10M0G53-B0	15pF ± 20%
		CSTLS10M0G56-B0	47pF ± 20%

Figure 16.9 Resonator Equivalent Circuit

16.6 Usage Note

Each of the products covered in this manual satisfy the electrical characteristics indicated. However, the actual electrical characteristics, operating margin and noise margin may differ from the indicated values due to differences in the manufacturing process, built-in ROM, layout pattern and other factors.

If a system evaluation test is conducted with the flash memory version, when switching to a mask ROM version, perform the same evaluation test with the mask ROM version.

Appendix

A. Instruction Set

A.1 Instruction List

Condition Code

Symbol	Description
Rd	General destination register
Rs	General source register
Rn	General register
ERd	General destination register (address register or 32-bit register)
ERs	General source register (address register or 32-bit register)
ERn	General register (32-bit register)
(EAd)	Destination operand
(EAs)	Source operand
PC	Program counter
SP	Stack pointer
CCR	Condition-code register
N	N (negative) flag in CCR
Z	Z (zero) flag in CCR
V	V (overflow) flag in CCR
С	C (carry) flag in CCR
disp	Displacement
\rightarrow	Transfer from the operand on the left to the operand on the right, or transition from the state on the left to the state on the right
+	Addition of the operands on both sides
_	Subtraction of the operand on the right from the operand on the left
×	Multiplication of the operands on both sides
÷	Division of the operand on the left by the operand on the right
٨	Logical AND of the operands on both sides
V	Logical OR of the operands on both sides
\oplus	Logical exclusive OR of the operands on both sides

Symbol	Description
7	NOT (logical complement)
(), <>	Contents of operand

Note: General registers include 8-bit registers (R0H to R7H and R0L to R7L) and 16-bit registers (R0 to R7 and E0 to E7).

Condition Code Notation (cont)

Symbol	Description
_	Changed according to execution result
*	Undetermined (no guaranteed value)
0	Cleared to 0
1	Set to 1
_	Not affected by execution of the instruction
Δ	Varies depending on conditions, described in notes



Table A.1 Instruction Set

1. Data Transfer Instructions

				A Inst			_		le ai)								No. Stat	of es*1
	Mnemonic	Operand Size	*xx#	Rn	@ERn	@(d, ERn)	@-ERn/@ERn+	@aa	@(d, PC)	@ @aa	ı	Operation	ı	Con	ditio	n Co	v	С	Normal	Advanced
MOV	MOV.B #xx:8, Rd	В	2									#xx:8 → Rd8	_	_	1	1	0	_	2	2
	MOV.B Rs, Rd	В		2								Rs8 → Rd8	_	_	1	1	0	_	2	2
	MOV.B @ERs, Rd	В			2							@ERs → Rd8	_	<u> </u>	1	1	0	_		1
	MOV.B @(d:16, ERs), Rd	В				4						@(d:16, ERs) → Rd8	_	_	1	1	0	_	6	3
	MOV.B @(d:24, ERs), Rd	В				8						@(d:24, ERs) → Rd8	_	_	1	1	0	_	1	0
	MOV.B @ERs+, Rd	В					2					@ERs → Rd8 ERs32+1 → ERs32	-	-	1	\$	0	_	6	3
	MOV.B @aa:8, Rd	В						2				@aa:8 → Rd8	_	<u> </u>	1	1	0	_		1
	MOV.B @aa:16, Rd	В						4				@aa:16 → Rd8	_	_	1	1	0	_	6	6
	MOV.B @aa:24, Rd	В						6				@aa:24 → Rd8	_	_	1	1	0	_	8	3
	MOV.B Rs, @ERd	В			2							Rs8 → @ERd	_	<u> </u>	1	1	0	_		1
	MOV.B Rs, @(d:16, ERd)	В				4						Rs8 → @(d:16, ERd)	_	_	1	1	0	_	6	3
	MOV.B Rs, @(d:24, ERd)	В				8						Rs8 → @(d:24, ERd)	_	_	1	1	0	_	1	0
	MOV.B Rs, @-ERd	В					2					ERd32-1 → ERd32 Rs8 → @ERd	_	_	1	\$	0	_	6	;
	MOV.B Rs, @aa:8	В						2				Rs8 → @aa:8	_	_	1	1	0	_	4	1
	MOV.B Rs, @aa:16	В						4				Rs8 → @aa:16	_	_	1	1	0	_	6	3
	MOV.B Rs, @aa:24	В						6				Rs8 → @aa:24	_	_	1	1	0	_	8	3
	MOV.W #xx:16, Rd	W	4									#xx:16 → Rd16	_	<u> </u>	1	1	0	_	4	1
	MOV.W Rs, Rd	W		2								Rs16 → Rd16	_	_	1	1	0	_	2	2
	MOV.W @ERs, Rd	W			2							@ERs → Rd16	_	_	1	1	0	_	4	1
	MOV.W @(d:16, ERs), Rd	W				4						@(d:16, ERs) → Rd16	_	_	1	1	0	_	6	3
	MOV.W @(d:24, ERs), Rd	W				8						@(d:24, ERs) → Rd16	_	_	1	1	0	_	1	0
	MOV.W @ERs+, Rd	W					2					@ERs → Rd16 ERs32+2 → @ERd32	_	_	1	\$	0	_	6	;
	MOV.W @aa:16, Rd	W						4				@aa:16 → Rd16	_	<u> </u>	1	1	0	_	6	3
	MOV.W @aa:24, Rd	W						6				@aa:24 → Rd16	_	_	1	1	0	_	8	3
	MOV.W Rs, @ERd	W			2							Rs16 → @ERd	_	_	1	1	0	_		1
	MOV.W Rs, @(d:16, ERd)	W				4						Rs16 → @(d:16, ERd)	_	-	1	1	0	_	6	3
	MOV.W Rs, @(d:24, ERd)	W				8						Rs16 → @(d:24, ERd)	_	_	1	\$	0	_	1	0

				A Inst					le a)								No. Stat	of es*1
	Mnemonic	Operand Size	#xx	Rn	@ERn	@(d, ERn)	@-ERn/@ERn+	@aa	@(d, PC)	@ @ aa	ı	Operation	ı	Con	ditio	n Co	ode	С	Normal	Advanced
MOV	MOV.W Rs, @-ERd	W					2					ERd32-2 → ERd32 Rs16 → @ERd	-	-	\$	\$	0	-	6	5
	MOV.W Rs, @aa:16	W						4				Rs16 → @aa:16	_	_	1	1	0	_	6	3
	MOV.W Rs, @aa:24	W						6				Rs16 → @aa:24	_	_	1	1	0	_	8	3
	MOV.L #xx:32, ERd	L	6									#xx:32 → ERd32	_	_	1	1	0	_	6	3
	MOV.L ERs, ERd	L		2								ERs32 → ERd32	_	_	1	1	0	_	2	2
	MOV.L @ERs, ERd	L			4							@ERs → ERd32	<u> </u>	_	1	1	0	<u> </u>	8	3
	MOV.L @(d:16, ERs), ERd	L				6						@(d:16, ERs) → ERd32	_	_	1	1	0	_	1	0
	MOV.L @(d:24, ERs), ERd	L				10						@(d:24, ERs) → ERd32	_	_	1	1	0	_	1	4
	MOV.L @ERs+, ERd	L					4					@ERs → ERd32 ERs32+4 → ERs32	-	-	\$	\$	0	-	1	0
	MOV.L @aa:16, ERd	L						6				@aa:16 → ERd32	_	_	1	1	0	_	1	0
	MOV.L @aa:24, ERd	L						8				@aa:24 → ERd32	_	_	1	1	0	_	1	2
	MOV.L ERs, @ERd	L			4							ERs32 → @ ERd	_	_	1	1	0	_	8	3
	MOV.L ERs, @(d:16, ERd)	L				6						ERs32 → @(d:16, ERd)	_	_	1	1	0	_	1	0
	MOV.L ERs, @(d:24, ERd)	L				10						ERs32 → @(d:24, ERd)	_	_	1	1	0	_	1	4
	MOV.L ERs, @-ERd	L					4					ERd32–4 → ERd32 ERs32 → @ ERd	_	_	\$	1	0	-	1	0
	MOV.L ERs, @aa:16	L						6				ERs32 → @aa:16	_	_	1	1	0	_	1	0
	MOV.L ERs, @aa:24	L						8				ERs32 → @aa:24	_	_	1	1	0	_	1	2
POP	POP.W Rn	W									2		_	_	\$	\$	0	_	6	6
	POP.L ERn	L									4		_	_	\$	\$	0	_	1	0
PUSH	PUSH.W Rn	W									2	$SP-2 \rightarrow SP$ $Rn16 \rightarrow @SP$	-	-	\$	\$	0	-	6	6
	PUSH.L ERn	L									4	$SP-4 \rightarrow SP$ $ERn32 \rightarrow @SP$	-	-	\$	\$	0	-	1	0
MOVFPE	MOVFPE @aa:16, Rd	В						4				Cannot be used in this LSI		anno s LS	t be	use	ed ir	ì	1	
MOVTPE	MOVTPE Rs, @aa:16	В						4				Cannot be used in this LSI		anno s LS	t be	use	ed ir	1		



2. Arithmetic Instructions

						essi				nd /tes)									. of tes*1
	Mnemonic	Operand Size	*xx	Rn	@ERn	@(d, ERn)	@-ERn/@ERn+	@aa	@(d, PC)	@ @ aa		Operation	1	Con	ditio	on Co	ode	С	Normal	Advanced
ADD	ADD.B #xx:8, Rd	В	2	_			_				<u>'</u>	Rd8+#xx:8 → Rd8	<u>'</u>	т	N	1	v	1		
ADD	ADD.B Rs, Rd	В	_	2								Rd8+Rs8 → Rd8		1	1	1	1	1		 2
	ADD.W #xx:16, Rd	w	4									Rd16+#xx:16 → Rd16	-	(1)	1	1	1	1		4
	ADD.W Rs, Rd	w		2								Rd16+Rs16 → Rd16	_	(1)	1	1	1	1	- 2	2
	ADD.L #xx:32, ERd	L	6									ERd32+#xx:32 → ERd32	-	(2)	\$	\$	\$	\$	(6
	ADD.L ERs, ERd	L		2								ERd32+ERs32 → ERd32	_	(2)	\$	\$	\$	\$	2	2
ADDX	ADDX.B #xx:8, Rd	В	2									$Rd8+#xx:8+C \rightarrow Rd8$	-	1	1	(3)	1	1	2	2
	ADDX.B Rs, Rd	В		2								Rd8+Rs8 +C → Rd8	_	1	1	(3)	1	1	2	2
ADDS	ADDS.L #1, ERd	L		2								ERd32+1 → ERd32	_	_	_	_	_	_	2	2
	ADDS.L #2, ERd	L		2								ERd32+2 → ERd32	_	_	_	_	_	_	2	2
	ADDS.L #4, ERd	L		2								ERd32+4 → ERd32	_	_	_	_	_	_	2	2
INC	INC.B Rd	В		2								Rd8+1 → Rd8	<u> </u>	_	1	1	1	_	2	2
	INC.W #1, Rd	W		2								Rd16+1 → Rd16	_	_	1	1	1	_	2	2
	INC.W #2, Rd	W		2								Rd16+2 → Rd16	_	_	1	1	1	_	2	2
	INC.L #1, ERd	L		2								ERd32+1 → ERd32	_	_	1	1	1	_	2	2
	INC.L #2, ERd	L		2								ERd32+2 → ERd32	_	_	1	1	1	_	2	2
DAA	DAA Rd	В		2								Rd8 decimal adjust → Rd8	-	*	\$	1	*	_	2	2
SUB	SUB.B Rs, Rd	В		2								Rd8–Rs8 → Rd8	_	1	1	1	1	1	2	2
	SUB.W #xx:16, Rd	W	4									Rd16–#xx:16 → Rd16	_	(1)	1	1	1	1		4
	SUB.W Rs, Rd	W		2								Rd16–Rs16 → Rd16	<u> </u>	(1)	1	1	1	1	2	2
	SUB.L #xx:32, ERd	L	6									ERd32-#xx:32 → ERd32	_	(2)	1	1	1	1	(6
	SUB.L ERs, ERd	L		2								ERd32–ERs32 → ERd32	_	(2)	1	1	1	1	2	2
SUBX	SUBX.B #xx:8, Rd	В	2									Rd8-#xx:8-C → Rd8	_	1	1	(3)	1	1	2	2
	SUBX.B Rs, Rd	В		2								Rd8–Rs8–C → Rd8	-	1	1	(3)	1	1	2	2
SUBS	SUBS.L #1, ERd	L		2								ERd32−1 → ERd32	_	_	_	<u> </u>	_	_	2	2
	SUBS.L #2, ERd	L		2								ERd32–2 → ERd32	_	_	_	<u> </u>	_	_	2	2
	SUBS.L #4, ERd	L		2								ERd32–4 → ERd32	_	_	_	_	_	_	2	2
DEC	DEC.B Rd	В		2								Rd8−1 → Rd8	-	_	1	1	1	_	2	2
	DEC.W #1, Rd	W		2								Rd16–1 → Rd16	_	_	1	1	1	_	2	2
	DEC.W #2, Rd	W		2								Rd16–2 → Rd16	_	_	1	1	1	_	2	2

									de a)								No Stat	of es*1
	Mnemonic	Operand Size	*xx	Rn	@ERn	@(d, ERn)	@-ERn/@ERn+	@aa	@(d, PC)	@ @ aa	ı	Operation				n Co			Vormal	Advanced
DEC	DEC.L #1, ERd	L	#	2	•	•	•	0	•	•	-	ERd32–1 → ERd32	ı	Н	N	Z	v	С	_	2
DEC	DEC.L #2, ERd	ᆫ		2							_	ERd32-2 \rightarrow ERd32	_	\vdash	1	↓	↓	_		2
DAS	DAS.Rd	В		2								Rd8 decimal adjust → Rd8	_	*	\$	1	*	_		2
MULXU	MULXU. B Rs, Rd	В		2								Rd8 × Rs8 → Rd16 (unsigned multiplication)	_	_	_	_	_	_	1	4
	MULXU. W Rs, ERd	W		2								Rd16 × Rs16 → ERd32 (unsigned multiplication)	_	_	_	_	_	_	2	2
MULXS	MULXS. B Rs, Rd	В		4								Rd8 × Rs8 → Rd16 (signed multiplication)	_	_	1	1	_	_	1	6
	MULXS. W Rs, ERd	W		4								Rd16 × Rs16 → ERd32 (signed multiplication)	_	-	1	1	_	_	2	!4
DIVXU	DIVXU. B Rs, Rd	В		2								Rd16 ÷ Rs8 → Rd16 (RdH: remainder, RdL: quotient) (unsigned division)	_	_	(6)	(7)	_	_	1	4
	DIVXU. W Rs, ERd	W		2								ERd32 ÷ Rs16 → ERd32 (Ed: remainder, Rd: quotient) (unsigned division)	_	_	(6)	(7)	_	_	2	2
DIVXS	DIVXS. B Rs, Rd	В		4								Rd16 ÷ Rs8 → Rd16 (RdH: remainder, RdL: quotient) (signed division)	_	_	(8)	(7)	_	_	1	6
	DIVXS. W Rs, ERd	W		4								ERd32 ÷ Rs16 → ERd32 (Ed: remainder, Rd: quotient) (signed division)	_	_	(8)	(7)	_	_	2	4
CMP	CMP.B #xx:8, Rd	В	2									Rd8-#xx:8	_	1	1	1	1	1	2	2
	CMP.B Rs, Rd	В		2								Rd8-Rs8	_	1	1	1	1	1	2	2
	CMP.W #xx:16, Rd	W	4									Rd16-#xx:16	_	(1)	1	\$	1	1		4
	CMP.W Rs, Rd	W		2								Rd16-Rs16	_	(1)	1	1	1	1	2	2
	CMP.L #xx:32, ERd	L	6									ERd32-#xx:32	_	(2)	1	1	1	1	- 4	4
	CMP.L ERs, ERd	L		2								ERd32-ERs32	_	(2)	1	1	1	1	2	2



					ddr					nd ⁄tes)								No Stat	of tes*1
	Mnemonic	Operand Size	_		@ERn	@(d, ERn)	-ERn/@ERn+	@aa	@(d, PC)	@aa		Operation		Con	ditic	n Co	ode		Normal	Advanced
		g	XX#	Rn	<u>@</u>	0	<u>@</u>	0	0	(9)	1		1	н	N	z	٧	С	ž	Ad
NEG	NEG.B Rd	В		2								0–Rd8 → Rd8	-	1	1	1	1	1	2	2
	NEG.W Rd	W		2								0–Rd16 → Rd16	-	1	1	1	1	1	2	2
	NEG.L ERd	L		2								0–ERd32 → ERd32	-	1	1	1	1	1	2	2
EXTU	EXTU.W Rd	W		2								0 → (<bits 15="" 8="" to=""> of Rd16)</bits>	_	_	0	\$	0	_	2	2
	EXTU.L ERd	L		2								0 → (<bits 16="" 31="" to=""> of ERd32)</bits>	_	-	0	1	0	_	2	2
EXTS	EXTS.W Rd	W		2								(<bit 7=""> of Rd16) → (<bits 15="" 8="" to=""> of Rd16)</bits></bit>	-	-	\$	1	0	_	2	2
	EXTS.L ERd	L		2								(<bit 15=""> of ERd32) → (<bits 16="" 31="" to=""> of ERd32)</bits></bit>	_	_	\$	\$	0	_	2	2

3. Logic Instructions

								Mod)									of es*1
	Mnemonic	Operand Size	#xx	Rn	@ERn	@(d, ERn)	-ERn/@ERn+	@aa	@(d, PC)	@aa		Operation		Con	ditio		1		Normal	Advanced
		<u> </u>	-	~	(a)	(a)	(9)	0	(a)	@	ı		ı	Н	N	Z	٧	С		
AND	AND.B #xx:8, Rd	В	2									Rd8∧#xx:8 → Rd8	_	_	1	1	0	_	_	2
	AND.B Rs, Rd	В		2								Rd8∧Rs8 → Rd8	_	_	1	1	0	_		2
	AND.W #xx:16, Rd	W	4									Rd16∧#xx:16 → Rd16	-	_	1	1	0	_	4	4
	AND.W Rs, Rd	W		2								Rd16∧Rs16 → Rd16	_	_	1	1	0	_	2	2
	AND.L #xx:32, ERd	L	6									$ERd32 \wedge \#xx: 32 \to ERd32$	_	_	1	1	0	_	6	ŝ
	AND.L ERs, ERd	L		4								ERd32∧ERs32 → ERd32	_	-	1	1	0	—	4	4
OR	OR.B #xx:8, Rd	В	2									Rd8∨#xx:8 → Rd8	_	_	1	1	0	—	2	2
	OR.B Rs, Rd	В		2								Rd8∨Rs8 → Rd8	_	_	1	1	0	—	2	2
	OR.W #xx:16, Rd	W	4									Rd16∨#xx:16 → Rd16	_	_	1	1	0	—	4	4
	OR.W Rs, Rd	W		2								Rd16∨Rs16 → Rd16	_	_	1	1	0	_	2	2
	OR.L #xx:32, ERd	L	6									ERd32∨#xx:32 → ERd32	_	_	1	1	0	_	6	ŝ
	OR.L ERs, ERd	L		4								ERd32∨ERs32 → ERd32	_	_	1	1	0	_	4	4
XOR	XOR.B #xx:8, Rd	В	2									Rd8⊕#xx:8 → Rd8	_	_	1	1	0	_	2	2
	XOR.B Rs, Rd	В		2								Rd8⊕Rs8 → Rd8	_	_	1	1	0	_	2	2
	XOR.W #xx:16, Rd	W	4									Rd16⊕#xx:16 → Rd16	_	_	1	1	0	_	4	4
	XOR.W Rs, Rd	W		2								Rd16⊕Rs16 → Rd16	_	_	1	1	0	_	2	2
	XOR.L #xx:32, ERd	L	6									ERd32⊕#xx:32 → ERd32	_	_	1	1	0	_	6	<u> </u>
	XOR.L ERs, ERd	L		4								ERd32⊕ERs32 → ERd32	_	_	1	1	0	_		4
NOT	NOT.B Rd	В		2								¬ Rd8 → Rd8	_	_	1	1	0	_	2	2
	NOT.W Rd	W		2								¬ Rd16 → Rd16	_	_	1	1	0	_	2	2
	NOT.L ERd	L		2								¬ Rd32 → Rd32	_	_	1	1	0	_	2	2

4. Shift Instructions

							ing l			nd ⁄tes)								No. Stat	
	Mnemonic	Operand Size	×	_	@ERn	@(d, ERn)	@-ERn/@ERn+	Фаа	@(d, PC)	@aa		Operation	,	Con	ditio	n Co	ode		Normal	Advanced
		ō	XX#	Ru	@	@	@	@	@	(9)	I		1	Н	N	z	٧	С	ž	¥
SHAL	SHAL.B Rd	В		2								□ +0	_	_	1	1	1	1	2	2
	SHAL.W Rd	W		2									_	_	1	1	1	1	2	2
	SHAL.L ERd	L		2								MSB LSB	_	_	1	1	1	1	2	2
SHAR	SHAR.B Rd	В		2									_	_	1	1	0	1	2	2
	SHAR.W Rd	W		2									_	_	1	1	0	1	2	2
	SHAR.L ERd	L		2								MSB LSB	_	_	1	1	0	1	2	2
SHLL	SHLL.B Rd	В		2								-0	_	_	\$	1	0	\$	2	2
	SHLL.W Rd	W		2									_	_	1	1	0	1	2	2
	SHLL.L ERd	L		2								MSB LSB	_	_	1	1	0	1	2	2
SHLR	SHLR.B Rd	В		2								0	_	_	1	1	0	1	2	2
	SHLR.W Rd	W		2								0-1	_	_	1	1	0	1	2	2
	SHLR.L ERd	L		2								MSB LSB	_	_	\$	1	0	\$	2	2
ROTXL	ROTXL.B Rd	В		2									_	_	1	1	0	1	2	2
	ROTXL.W Rd	W		2									_	_	1	1	0	1	2	2
	ROTXL.L ERd	L		2								MSB ≺ LSB	_	_	\$	1	0	\$	2	2
ROTXR	ROTXR.B Rd	В		2									_	_	\$	1	0	1	2	2
	ROTXR.W Rd	W		2									_	_	1	1	0	1	2	2
	ROTXR.L ERd	L		2								MSB → LSB	_	_	1	1	0	1	2	2
ROTL	ROTL.B Rd	В		2									_	_	\$	1	0	\$	2	2
	ROTL.W Rd	W		2									_	_	\$	1	0	\$	2	2
	ROTL.L ERd	L		2								MSB ← LSB	_	_	\$	1	0	\$	2	2
ROTR	ROTR.B Rd	В		2								r ra	_		1	1	0	1	2	2
	ROTR.W Rd	W		2									_	_	1	1	0	1	2	2
	ROTR.L ERd	L		2								MSB → LSB	_		1	1	0	1	2	2

5. Bit-Manipulation Instructions

				A Inst			ng I Ler)								No Stat	of es*1
	Mnemonic	Operand Size	#xx	Rn	@ERn	@(d, ERn)	@-ERn/@ERn+	@aa	@(d, PC)	@ @ aa	1	Operation	1	Con	ditio	n Co	ode	С	Normal	Advanced
BSET	BSET #xx:3, Rd	В		2								(#xx:3 of Rd8) ← 1	_	_	_	_	_	_	2	2
	BSET #xx:3, @ERd	В			4							(#xx:3 of @ERd) ← 1	_	_	_	_	_	_	8	3
	BSET #xx:3, @aa:8	В						4				(#xx:3 of @aa:8) ← 1	_	_	_	-	_	_	3	3
	BSET Rn, Rd	В		2								(Rn8 of Rd8) ← 1	_	_	_	_	_	_	2	2
	BSET Rn, @ERd	В			4							(Rn8 of @ERd) ← 1	<u> </u>	_	_	_	_	_	8	3
	BSET Rn, @aa:8	В						4				(Rn8 of @aa:8) ← 1	_	_	_	-	_	_	3	3
BCLR	BCLR #xx:3, Rd	В		2								(#xx:3 of Rd8) ← 0	_	_	_	_	_	_	2	2
	BCLR #xx:3, @ERd	В			4							(#xx:3 of @ERd) ← 0	_	_	_	_	_	_	8	3
	BCLR #xx:3, @aa:8	В						4				(#xx:3 of @aa:8) ← 0	_	_	_	_	_	_	8	3
	BCLR Rn, Rd	В		2								(Rn8 of Rd8) ← 0	_	_	_	_	_	_	2	2
	BCLR Rn, @ERd	В			4							(Rn8 of @ERd) ← 0	_	_	_	_	_	_	8	3
	BCLR Rn, @aa:8	В						4				(Rn8 of @aa:8) ← 0	_	_	_	_	_	_	8	3
BNOT	BNOT #xx:3, Rd	В		2								(#xx:3 of Rd8) ← ¬ (#xx:3 of Rd8)	-	_	_	_	-	_	2	2
	BNOT #xx:3, @ERd	В			4							(#xx:3 of @ERd) ← ¬ (#xx:3 of @ERd)	_	_	_	_	_	_	8	3
	BNOT #xx:3, @aa:8	В						4				(#xx:3 of @aa:8) ← ¬ (#xx:3 of @aa:8)	-	-	_	-	-	_	8	3
	BNOT Rn, Rd	В		2								(Rn8 of Rd8) ← ¬ (Rn8 of Rd8)	_	_	_	_	_	_	2	2
	BNOT Rn, @ERd	В			4							(Rn8 of @ERd) ← ¬ (Rn8 of @ERd)	_	_	_	_	_	_	8	3
	BNOT Rn, @aa:8	В						4				(Rn8 of @aa:8) ← ¬ (Rn8 of @aa:8)	_	_	_	_	_	_	8	3
BTST	BTST #xx:3, Rd	В		2								¬ (#xx:3 of Rd8) → Z	_	_	_	1	_	-	2	2
	BTST #xx:3, @ERd	В			4							¬ (#xx:3 of @ERd) → Z	_	_	_	1	_	_	6	3
	BTST #xx:3, @aa:8	В						4				¬ (#xx:3 of @aa:8) → Z	_	_	_	1	_	_	6	<u> </u>
	BTST Rn, Rd	В		2								¬ (Rn8 of @Rd8) \rightarrow Z	_	_	_	1	_	_	2	2
	BTST Rn, @ERd	В			4							¬ (Rn8 of @ERd) → Z	_	_	_	1	_	_	6	<u> </u>
	BTST Rn, @aa:8	В						4				¬ (Rn8 of @aa:8) → Z	_	_	_	1	_	_	6	3
BLD	BLD #xx:3, Rd	В		2								(#xx:3 of Rd8) → C	_	_	_	_	_	1	2	2



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	Mnemonic	Operand Size	xx#	Rn	@ERn	@(d, ERn)	@-ERn/@ERn+	@aa	@(d, PC)	Ф Фаа		Operation			ditio		T		Normal	Advanced
	I		#	~		Ø	Ø	ø	ø	0	ı		1	Н	N	z	٧	C		
BLD	BLD #xx:3, @ERd	В			4							(#xx:3 of @ERd) → C	_	_	_	_	_	1	6	
	BLD #xx:3, @aa:8	В						4				(#xx:3 of @aa:8) → C	_		_	_	_	1	6	
BILD	BILD #xx:3, Rd	В		2								¬ (#xx:3 of Rd8) → C	_	_	_	_	_	1	2	2
	BILD #xx:3, @ERd	В			4							¬ (#xx:3 of @ERd) → C	_	_	_	_	_	1	6	j
	BILD #xx:3, @aa:8	В						4				¬ (#xx:3 of @aa:8) → C	_	_	_	_	_	1	6	;
BST	BST #xx:3, Rd	В		2								$C \rightarrow (\#xx:3 \text{ of Rd8})$	_	_	—	_	_	_	2	2
	BST #xx:3, @ERd	В			4							$C \rightarrow (\#xx:3 \text{ of } @ERd24)$	_	_	_	_	—	_	8	3
	BST #xx:3, @aa:8	В						4				C → (#xx:3 of @aa:8)	_	_	_	_	_	_	8	3
BIST	BIST #xx:3, Rd	В		2								$\neg C \rightarrow (\#xx:3 \text{ of Rd8})$	_	_	_	_	_	_	2	2
	BIST #xx:3, @ERd	В			4							$\neg C \rightarrow (\#xx:3 \text{ of } @ERd24)$	_	_	_	_	_	_	8	3
	BIST #xx:3, @aa:8	В						4				¬ C → (#xx:3 of @aa:8)	_	_	_	_	_	_	8	3
BAND	BAND #xx:3, Rd	В		2								$C \land (\#xx:3 \text{ of Rd8}) \rightarrow C$	_	_	_	_	_	1	2	2
	BAND #xx:3, @ERd	В			4							C∧(#xx:3 of @ERd24) → C	_	_	_	_	_	1	6	;
	BAND #xx:3, @aa:8	В						4				C∧(#xx:3 of @aa:8) → C	_	_	_	_	_	1	6	;
BIAND	BIAND #xx:3, Rd	В		2								$C \land \neg \text{ (#xx:3 of Rd8)} \rightarrow C$	_	_	_	_	_	1	2	2
	BIAND #xx:3, @ERd	В			4							C∧¬ (#xx:3 of @ERd24) → C	_	_	_	_	_	1	6	3
	BIAND #xx:3, @aa:8	В						4				C∧¬ (#xx:3 of @aa:8) → C	_	_	_	_	_	1	6	3
BOR	BOR #xx:3, Rd	В		2								C∨(#xx:3 of Rd8) → C	_	_	_	_	_	1	2	2
	BOR #xx:3, @ERd	В			4							C√(#xx:3 of @ERd24) → C	_	_	_	_	_	1	6	3
	BOR #xx:3, @aa:8	В						4				C√(#xx:3 of @aa:8) → C	_	_	_	_	_	1	6	3
BIOR	BIOR #xx:3, Rd	В		2								C∨¬ (#xx:3 of Rd8) → C	_	_	_	_	_	1	2	2
	BIOR #xx:3, @ERd	В			4							C∨¬ (#xx:3 of @ERd24) → C	_	_	_	_	_	1	6	3
	BIOR #xx:3, @aa:8	В						4				C∨¬ (#xx:3 of @aa:8) → C	_	_	_	_	_	1	6	,
BXOR	BXOR #xx:3, Rd	В		2								$C⊕(\#xx:3 \text{ of Rd8}) \rightarrow C$	_	<u> </u>	_	_	_	1		2
	BXOR #xx:3, @ERd	В			4							C⊕(#xx:3 of @ERd24) → C	_		_	_	_	1	6	
	BXOR #xx:3, @aa:8	В						4				C⊕(#xx:3 of @aa:8) → C	_		_		_	1	6	
BIXOR	BIXOR #xx:3, Rd	В		2				·				$C \oplus \neg (\#xx:3 \text{ of Rd8}) \rightarrow C$	_		_	_	_	1		2
5.7.011	BIXOR #xx:3, @ERd	В		_	4							$C \oplus \neg$ (#xx:3 of @ERd24) → C	_	\vdash			_	1	6	
	BIXOR #xx:3, @aa:8	В			Ė			4				$C \oplus \neg (\#xx:3 \text{ of } @aa:8) \rightarrow C$	_					1	6	
	2							L .	<u> </u>			Co (MARIO OI GUALO) -> O			<u> </u>			*		

6. Branching Instructions

									le ai)									No Stat	. of es ^{*1}
	Mnemonic	Operand Size	#xx	Rn	@ERn	@(d, ERn)	@-ERn/@ERn+	@aa	@(d, PC)	@ @ aa		Oper	ation Branch			ditio			ı	Normal	Advanced
_	DDA 40 (DT 40)	0	#	~	0	0	0	0	2	0	I	16 100	Condition	1	Н	N	Z	٧	С		
Bcc	BRA d:8 (BT d:8)	-	_						4			If condition is true then	Always	_	_	_	_	_	_		1 3
	BRA d:16 (BT d:16)	+-							2			PC ← PC+d	Never	_	_	_	_		_	_	1
	BRN d:8 (BF d:8)	-	_						4			else next;	Never	_	_	_	<u> </u>	_	_		
	BRN d:16 (BF d:16)	+-	_										0.7.0	_	_	_	<u> </u>	_	_	6	
	BHI d:8	+-							2			+	C∨ Z = 0	_	_	_	_	_	-		1
	BHI d:16	-							4			+	0.7.4	_	_	_	-	_	-		3
	BLS d:8	_	_						2				C∨ Z = 1	_	_	_	_	_	_	4	
	BLS d:16	-							4			-	0 0	_	_	_	_	_	_	(
	BCC d:8 (BHS d:8)	_							2				C = 0	_	_	_	_	_	_		1
	BCC d:16 (BHS d:16)	<u> -</u>							4					_	_	_	_	_	_	(
	BCS d:8 (BLO d:8)	-	_						2				C = 1	_	_	_	_	_	_	4	•
	BCS d:16 (BLO d:16)	<u> -</u>							4					_	_	_		_	_		5
	BNE d:8	_							2				Z = 0	_	_	_	_	_	_		1
	BNE d:16	-							4					_	_	_	_	_	_	6	
	BEQ d:8								2				Z = 1	_	_	_	_	_	_		1
	BEQ d:16	-							4					_	_	_	_	_	_	6	
	BVC d:8	-							2				V = 0	_	_	_	_	_	_	4	1
	BVC d:16	_							4					_	_	_	_	_	-	6	3
	BVS d:8	_							2				V = 1	_	_	_	_	_	_	4	1
	BVS d:16	_							4						_	_	_	_	_	6	3
	BPL d:8	_							2				N = 0	_	_	_	_	_	_	4	1
	BPL d:16	-							4					_	_	_	_	_	_	6	6
	BMI d:8	_							2				N = 1	_		_		_		4	1
	BMI d:16								4										Ŀ	6	6
	BGE d:8								2				N⊕V = 0	_	_	_	_	_	_	4	1
	BGE d:16								4					_					Ŀ	6	3
	BLT d:8	_							2				N⊕V = 1	_	_	_	_	_	_	4	1
	BLT d:16	_							4					_	_	_	_	_	_	6	3
	BGT d:8	_							2				Z∨ (N⊕V) = 0	_	_	_	_	_	_	4	1
	BGT d:16	-							4					_	_	_	_	_	_	6	6
	BLE d:8	<u> </u>							2			1	Z∨ (N⊕V) = 1	_	_	_	_	_	_		1
	BLE d:16	-							4			1		_	_	_	_	_	_	6	6

				A Inst			ng l Ler)									of es*1
	Mnemonic	Operand Size			@ERn	@(d, ERn)	@-ERn/@ERn+	la	@(d, PC)	@aa		Operation		Con	ditio	n Co	ode		Normal	Advanced
		o	XX#	R	@	e	9	@aa	<u>@</u>	<u>@</u>	ı		1	н	N	z	٧	С	2	Ad
JMP	JMP @ERn	_			2							PC ← ERn	_	_	_	_	_	_	4	4
	JMP @aa:24	_						4				PC ← aa:24	_	<u> </u>	_	_	_	_	(3
	JMP @ @aa:8	_								2		PC ← @aa:8	_	<u> </u>	_	_	_	_	8	10
BSR	BSR d:8	_							2			$PC \rightarrow @-SP$ $PC \leftarrow PC+d:8$	_	-	_	-	_	_	6	8
	BSR d:16	-							4			PC → @-SP PC ← PC+d:16	_	-	-	_	_	_	8	10
JSR	JSR @ERn	_			2							$\begin{array}{l} PC \to @-SP \\ PC \leftarrow ERn \end{array}$	_	-	_	_	_	_	6	8
	JSR @aa:24	_						4				PC → @-SP PC ← aa:24	_	_	_	_	_	_	8	10
	JSR @ @aa:8	_								2		PC → @-SP PC ← @aa:8	_	-	_	_	-	-	8	12
RTS	RTS	 —									2	PC ← @SP+	_	-	_	_	_	_	8	10

7. System Control Instructions

							ng l Ler			nd /tes)								No. Stat	of es ^{*1}
	Mnemonic	Operand Size	×		ERn	@(d, ERn)	@-ERn/@ERn+	@аа	@(d, PC)	@aa		Operation		Con	ditio	n Co	ode		Normal	Advanced
		ŏ	XX#	Rn	@	0	9	0	0	@	I		1	Н	N	z	٧	С	ž	Ad
RTE	RTE	-										CCR ← @SP+ PC ← @SP+	1	1	1	1	1	1	1	0
SLEEP	SLEEP	_										Transition to power- down state							2	2
LDC	LDC #xx:8, CCR	В	2									#xx:8 → CCR	1	1	1	1	1	1	2	2
	LDC Rs, CCR	В		2								Rs8 → CCR	1	1	1	1	1	1	2	2
	LDC @ERs, CCR	W			4							@ERs → CCR	1	1	1	1	1	1	6	6
	LDC @(d:16, ERs), CCR	W				6						@(d:16, ERs) → CCR	1	1	1	1	1	1	8	3
	LDC @(d:24, ERs), CCR	W				10						@(d:24, ERs) → CCR	1	1	1	1	1	1	1:	2
	LDC @ERs+, CCR	W					4					@ERs → CCR ERs32+2 → ERs32	\$	\$	\$	\$	1	\$	8	3
	LDC @aa:16, CCR	W						6				@aa:16 → CCR	1	1	1	1	1	1	8	3
	LDC @aa:24, CCR	W						8				@aa:24 → CCR	1	1	1	1	1	1	1	0
STC	STC CCR, Rd	В		2								CCR → Rd8							2	2
	STC CCR, @ERd	W			4							CCR → @ERd							6	3
	STC CCR, @(d:16, ERd)	W				6						CCR → @(d:16, ERd)							8	3
	STC CCR, @(d:24, ERd)	W				10						CCR → @(d:24, ERd)							1:	2
	STC CCR, @-ERd	W					4					ERd32–2 \rightarrow ERd32 CCR \rightarrow @ERd			8	3				
	STC CCR, @aa:16	W						6				CCR → @aa:16			8	3				
	STC CCR, @aa:24	W						8				CCR → @aa:24							1	0
ANDC	ANDC #xx:8, CCR	В	2									CCR∧#xx:8 → CCR	1	1	1	1	1	1	2	2
ORC	ORC #xx:8, CCR	В	2									CCR√#xx:8 → CCR	1	1	1	1	1	1	2	2
XORC	XORC #xx:8, CCR	В	2									CCR⊕#xx:8 → CCR	1	1	\$	1	1	1	2	2
NOP	NOP	_									2	PC ← PC+2							2	2

8. Block Transfer Instructions

					ddr)								No. State	
	Mnemonic	Operand Size	×		@ERn	@(d, ERn)	-ERn/@ERn+	33	@(d, PC)	@aa		Operation		Con	ditio	n Co	ode		Normal	Advanced
		o	XX#	R	<u>@</u>	0	9	@aa	0	0	1		Т	н	N	z	٧	С	2	Ad
EEPMOV	ЕЕРМОV. В	_									4	if R4L \neq 0 then repeat @R5 \rightarrow @R6 R5+1 \rightarrow R5 R6+1 \rightarrow R6 R4L-1 \rightarrow R4L until R4L=0 else next	_	_	_	_		_	8+ 4n*2	
	EEPMOV. W										4	$\begin{array}{l} \text{if R4} \neq 0 \text{ then} \\ \text{repeat} & @R5 \rightarrow @R6 \\ & R5+1 \rightarrow R5 \\ & R6+1 \rightarrow R6 \\ & R4-1 \rightarrow R4 \\ \text{until} & R4=0 \\ \text{else next} \end{array}$	_	_	_	_		_	8+ 4n*2	

Notes: 1. The number of states in cases where the instruction code and its operands are located in on-chip memory is shown here. For other cases, see appendix A.3, Number of Execution States.

- 2. n is the value set in register R4L or R4.
- (1) Set to 1 when a carry or borrow occurs at bit 11; otherwise cleared to 0.
- (2) Set to 1 when a carry or borrow occurs at bit 27; otherwise cleared to 0.
- (3) Retains its previous value when the result is zero; otherwise cleared to 0.
- (4) Set to 1 when the adjustment produces a carry; otherwise retains its previous value.
- (5) The number of states required for execution of an instruction that transfers data in synchronization with the E clock is variable.
- (6) Set to 1 when the divisor is negative; otherwise cleared to 0.
- (7) Set to 1 when the divisor is zero; otherwise cleared to 0.
- (8) Set to 1 when the quotient is negative; otherwise cleared to 0.

A.2 Operation Code Map

Table A.2 Operation Code Map (1)



Table A.2 Operation Code Map (2)

AH AL	0	-	5	е	4	2	9	7	80	0	∢	В	O	۵	ш	ш
01	MOV				LDC/STC				SLEEP				Table A-2 Table A-2 (3)	Table A-2 (3)		Table A-2 (3)
0A	INC											AD	ADD			
0B	ADDS					NC		NC	ADİ	ADDS				INC		NC
HO.	DAA											MOV	20			
10	SHLL			SHLL					SH	SHAL		SHAL				
11	SHLR	- Н		SHLR					HS.	SHAR		SHAR				
12	ROTXL	Z		ROTXL					RO	ROTL		ROTL				
13	ROTXR	XR		ROTXR					RO	ROTR		ROTR				
17	NOT	Т		NOT		EXTU		EXTU	NE	NEG		NEG		EXTS		EXTS
1A	DEC											SUB	JB			
1B	SUBS					DEC		DEC	SUB	BI				DEC		DEC
1F	DAS											CN	CMP			
58	BRA	BRN	BHI	BLS	BCC	BCS	BNE	ВЕО	BVC	BVS	BPL	BMI	BGE	BLT	ВСТ	BLE
79	MOV	ADD	CMP	SUB	OR	XOR	AND									
7A	MOV	ADD	CMP	SUB	HO	XOR	AND									

Instruction code: 1st byte

BH BL

AH

Table A.2 **Operation Code Map (3)**

Instruction code:	1 2 3 4 8 9 A B C D E F	LDC LDC LDC STC LDC STC	NAULXS NAULXS	DIVXS DIVXS	OR XOR AND	BIST	BTST BOR BXOR BIAND BILD	ET BNOT BCLR BCLR BST BST	ET BNOT BCLR	BIST	BTST BOR BXOR BIAND BILD	ET BNOT BCLR BCLR BST BST	ET BNOT BOLR
e: 1st AH	-			DIVXS				BNOT	BNOT			BNOT	BNOT
on code	0		MULXS					BSET	BSET			BSET	BSET
Instruct	ALBH ALBH BLCH	01406	01C05	01D05	01F06	7Cr06*1	7Cr07*1	7Dr06*1	7Dr07*1	7Eaa6*2	7Eaa7*2	7Faa6*2	7Faa7*2

Notes: 1. r is the register designation field. 2. aa is the absolute address field.

A.3 Number of Execution States

The status of execution for each instruction of the H8/300H CPU and the method of calculating the number of states required for instruction execution are shown below. Table A.4 shows the number of cycles of each type occurring in each instruction, such as instruction fetch and data read/write. Table A.3 shows the number of states required for each cycle. The total number of states required for execution of an instruction can be calculated by the following expression:

Execution states =
$$I \times S_1 + J \times S_2 + K \times S_K + L \times S_L + M \times S_M + N \times S_N$$

Examples: When instruction is fetched from on-chip ROM, and an on-chip RAM is accessed.

BSET #0, @FF00

From table A.4:

$$I = L = 2$$
, $J = K = M = N = 0$

From table A.3:

$$S_{I} = 2$$
, $S_{L} = 2$

Number of states required for execution = $2 \times 2 + 2 \times 2 = 8$

When instruction is fetched from on-chip ROM, branch address is read from on-chip ROM, and on-chip RAM is used for stack area.

JSR @@ 30

From table A.4:

$$I = 2$$
, $J = K = 1$, $L = M = N = 0$

From table A.3:

$$S_I = S_J = S_K = 2$$

Number of states required for execution = $2 \times 2 + 1 \times 2 + 1 \times 2 = 8$

Table A.3 Number of Cycles in Each Instruction

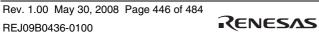
Execution Status		Ad	ccess Location
(Instruction Cycle)		On-Chip Memory	On-Chip Peripheral Module
Instruction fetch	S	2	_
Branch address read	S _J	_	
Stack operation	S _K	_	
Byte data access	S _L	_	2 or 3*
Word data access	S _M	_	_
Internal operation	S _N		1

Note: * Depends on which on-chip peripheral module is accessed. See section 15.1, Register Addresses (Address Order).

Table A.4 Number of Cycles in Each Instruction

Instruction	Mnemonic	Instruction Fetch I	Branch Addr. Read J	Stack Operation K	Byte Data Access L	Word Data Access M	Internal Operation N
ADD	ADD.B #xx:8, Rd	1					
	ADD.B Rs, Rd	1					
	ADD.W #xx:16, Rd	2					
	ADD.W Rs, Rd	1					
	ADD.L #xx:32, ERd	3					
	ADD.L ERs, ERd	1					
ADDS	ADDS #1/2/4, ERd	1					
ADDX	ADDX #xx:8, Rd	1					
	ADDX Rs, Rd	1					
AND	AND.B #xx:8, Rd	1					
	AND.B Rs, Rd	1					
	AND.W #xx:16, Rd	2					
	AND.W Rs, Rd	1					
	AND.L #xx:32, ERd	3					
	AND.L ERs, ERd	2					
ANDC	ANDC #xx:8, CCR	1					
BAND	BAND #xx:3, Rd	1					
	BAND #xx:3, @ERd	2			1		
	BAND #xx:3, @aa:8	2			1		
Bcc	BRA d:8 (BT d:8)	2					
	BRN d:8 (BF d:8)	2					
	BHI d:8	2					
	BLS d:8	2					
	BCC d:8 (BHS d:8)	2					
	BCS d:8 (BLO d:8)	2					
	BNE d:8	2					
	BEQ d:8	2					
	BVC d:8	2					
	BVS d:8	2					
	BPL d:8	2					
	BMI d:8	2					
	BGE d:8	2					

Instruction	Mnemonic	Instruction Fetch I	Branch Addr. Read J	Stack Operation K	Byte Data Access L	Word Data Access M	Internal Operation N
Всс	BLT d:8	2					
	BGT d:8	2					
	BLE d:8	2					
	BRA d:16(BT d:16)	2					2
	BRN d:16(BF d:16)	2					2
	BHI d:16	2					2
	BLS d:16	2					2
	BCC d:16(BHS d:16)	2					2
	BCS d:16(BLO d:16)	2					2
	BNE d:16	2					2
	BEQ d:16	2					2
	BVC d:16	2					2
	BVS d:16	2					2
	BPL d:16	2					2
	BMI d:16	2					2
	BGE d:16	2					2
	BLT d:16	2					2
	BGT d:16	2					2
	BLE d:16	2					2
BCLR	BCLR #xx:3, Rd	1					
	BCLR #xx:3, @ERd	2			2		
	BCLR #xx:3, @aa:8	2			2		
	BCLR Rn, Rd	1					
	BCLR Rn, @ERd	2			2		
	BCLR Rn, @aa:8	2			2		
BIAND	BIAND #xx:3, Rd	1					
	BIAND #xx:3, @ERd	2			1		
	BIAND #xx:3, @aa:8	2			1		
BILD	BILD #xx:3, Rd	1					
	BILD #xx:3, @ERd	2			1		
	BILD #xx:3, @aa:8	2			1		



Instruction	Mnemonic	Instruction Fetch	Branch Addr. Read J	Stack Operation K	Byte Data Access L	Word Data Access M	Internal Operation N
BIOR	BIOR #xx:3, Rd	1					
	BIOR #xx:3, @ERd	2			1		
	BIOR #xx:3, @aa:8	2			1		
BIST	BIST #xx:3, Rd	1					
	BIST #xx:3, @ERd	2			2		
	BIST #xx:3, @aa:8	2			2		
BIXOR	BIXOR #xx:3, Rd	1					
	BIXOR #xx:3, @ERd	2			1		
	BIXOR #xx:3, @aa:8	2			1		
BLD	BLD #xx:3, Rd	1					
	BLD #xx:3, @ERd	2			1		
	BLD #xx:3, @aa:8	2			1		
BNOT	BNOT #xx:3, Rd	1					
	BNOT #xx:3, @ERd	2			2		
	BNOT #xx:3, @aa:8	2			2		
	BNOT Rn, Rd	1					
	BNOT Rn, @ERd	2			2		
	BNOT Rn, @aa:8	2			2		
BOR	BOR #xx:3, Rd	1					
	BOR #xx:3, @ERd	2			1		
	BOR #xx:3, @aa:8	2			1		
BSET	BSET #xx:3, Rd	1					
	BSET #xx:3, @ERd	2			2		
	BSET #xx:3, @aa:8	2			2		
	BSET Rn, Rd	1					
	BSET Rn, @ERd	2			2		
	BSET Rn, @aa:8	2			2		
BSR	BSR d:8	2		1			
	BSR d:16	2		1			2
BST	BST #xx:3, Rd	1					
	BST #xx:3, @ERd	2			2		
	BST #xx:3, @aa:8	2			2		

Instruction	Mnemonic	Instruction Fetch I	Branch Addr. Read J	Stack Operation K	Byte Data Access L	Word Data Access M	Internal Operation N
BTST	BTST #xx:3, Rd	1					
	BTST #xx:3, @ERd	2			1		
	BTST #xx:3, @aa:8	2			1		
	BTST Rn, Rd	1					
	BTST Rn, @ERd	2			1		
	BTST Rn, @aa:8	2			1		
BXOR	BXOR #xx:3, Rd	1					
	BXOR #xx:3, @ERd	2			1		
	BXOR #xx:3, @aa:8	2			1		
CMP	CMP.B #xx:8, Rd	1					
	CMP.B Rs, Rd	1					
	CMP.W #xx:16, Rd	2					
	CMP.W Rs, Rd	1					
	CMP.L #xx:32, ERd	3					
	CMP.L ERs, ERd	1					
DAA	DAA Rd	1					
DAS	DAS Rd	1					
DEC	DEC.B Rd	1					
	DEC.W #1/2, Rd	1					
	DEC.L #1/2, ERd	1					
DUVXS	DIVXS.B Rs, Rd	2					12
	DIVXS.W Rs, ERd	2					20
DIVXU	DIVXU.B Rs, Rd	1					12
	DIVXU.W Rs, ERd	1					20
EEPMOV	EEPMOV.B	2			2n+2*1		
	EEPMOV.W	2			2n+2*1		
EXTS	EXTS.W Rd	1					
	EXTS.L ERd	1					
EXTU	EXTU.W Rd	1					
	EXTU.L ERd	1					

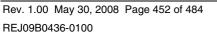


Instruction	Mnemonic	Instruction Fetch	Branch Addr. Read J	Stack Operation K	Byte Data Access L	Word Data Access M	Internal Operation N
INC	INC.B Rd	1					
	INC.W #1/2, Rd	1					
	INC.L #1/2, ERd	1					
JMP	JMP @ERn	2					
	JMP @aa:24	2					2
	JMP @@aa:8	2	1				2
JSR	JSR @ERn	2		1			
	JSR @aa:24	2		1			2
	JSR @@aa:8	2	1	1			
LDC	LDC #xx:8, CCR	1					
	LDC Rs, CCR	1					
	LDC@ERs, CCR	2				1	
	LDC@(d:16, ERs), CCR	3				1	
	LDC@(d:24,ERs), CCR	5				1	
	LDC@ERs+, CCR	2				1	2
	LDC@aa:16, CCR	3				1	
	LDC@aa:24, CCR	4				1	
MOV	MOV.B #xx:8, Rd	1					
	MOV.B Rs, Rd	1					
	MOV.B @ERs, Rd	1			1		
	MOV.B @(d:16, ERs), Rd	2			1		
	MOV.B @(d:24, ERs), Rd	4			1		
	MOV.B @ERs+, Rd	1			1		2
	MOV.B @aa:8, Rd	1			1		
	MOV.B @aa:16, Rd	2			1		
	MOV.B @aa:24, Rd	3			1		
	MOV.B Rs, @Erd	1			1		
	MOV.B Rs, @(d:16, ERd)	2			1		
	MOV.B Rs, @(d:24, ERd)	4			1		
	MOV.B Rs, @-ERd	1			1		2
	MOV.B Rs, @aa:8	1			1		

Instruction	Mnemonic	Instruction Fetch I	Branch Addr. Read J	Stack Operation K	Byte Data Access L	Word Data Access M	Internal Operation N
MOV	MOV.B Rs, @aa:16	2			1		
	MOV.B Rs, @aa:24	3			1		
	MOV.W #xx:16, Rd	2					
	MOV.W Rs, Rd	1					
	MOV.W @ERs, Rd	1				1	
	MOV.W @(d:16,ERs), Rd	2				1	
	MOV.W @(d:24,ERs), Rd	4				1	
	MOV.W @ERs+, Rd	1				1	2
	MOV.W @aa:16, Rd	2				1	
	MOV.W @aa:24, Rd	3				1	
	MOV.W Rs, @ERd	1				1	
	MOV.W Rs, @(d:16,ERd)	2				1	
	MOV.W Rs, @(d:24,ERd)	4				1	
MOV	MOV.W Rs, @-ERd	1				1	2
	MOV.W Rs, @aa:16	2				1	
	MOV.W Rs, @aa:24	3				1	
	MOV.L #xx:32, ERd	3					
	MOV.L ERs, ERd	1					
	MOV.L @ERs, ERd	2				2	
	MOV.L @(d:16,ERs), ERd	3				2	
	MOV.L @(d:24,ERs), ERd	5				2	
	MOV.L @ERs+, ERd	2				2	2
	MOV.L @aa:16, ERd	3				2	
	MOV.L @aa:24, ERd	4				2	
	MOV.L ERs,@ERd	2				2	
	MOV.L ERs, @(d:16,ERd)	3				2	
	MOV.L ERs, @(d:24,ERd)	5				2	
	MOV.L ERs, @-ERd	2				2	2
	MOV.L ERs, @aa:16	3				2	
	MOV.L ERs, @aa:24	4				2	
MOVFPE	MOVFPE @aa:16, Rd*2	2			1		
MOVTPE	MOVTPE Rs,@aa:16*2	2			1		

Instruction	Mnemonic	Instruction Fetch I	Branch Addr. Read J	Stack Operation K	Byte Data Access L	Word Data Access M	Internal Operation N
MULXS	MULXS.B Rs, Rd	2					12
	MULXS.W Rs, ERd	2					20
MULXU	MULXU.B Rs, Rd	1					12
	MULXU.W Rs, ERd	1					20
NEG	NEG.B Rd	1					
	NEG.W Rd	1					
	NEG.L ERd	1					
NOP	NOP	1					
NOT	NOT.B Rd	1					
	NOT.W Rd	1					
	NOT.L ERd	1					
OR	OR.B #xx:8, Rd	1					
	OR.B Rs, Rd	1					
	OR.W #xx:16, Rd	2					
	OR.W Rs, Rd	1					
	OR.L #xx:32, ERd	3					
	OR.L ERs, ERd	2					
ORC	ORC #xx:8, CCR	1					
POP	POP.W Rn	1				1	2
	POP.L ERn	2				2	2
PUSH	PUSH.W Rn	1				1	2
	PUSH.L ERn	2				2	2
ROTL	ROTL.B Rd	1					
	ROTL.W Rd	1					
	ROTL.L ERd	1					
ROTR	ROTR.B Rd	1					
	ROTR.W Rd	1					
	ROTR.L ERd	1					
ROTXL	ROTXL.B Rd	1					
	ROTXL.W Rd	1					
	ROTXL.L ERd	1					

Instruction	Mnemonic	Instruction Fetch I	Branch Addr. Read J	Stack Operation K	Byte Data Access L	Word Data Access M	Internal Operation N
ROTXR	ROTXR.B Rd	1					
	ROTXR.W Rd	1					
	ROTXR.L ERd	1					
RTE	RTE	2		2			2
RTS	RTS	2		1			2
SHAL	SHAL.B Rd	1					
	SHAL.W Rd	1					
	SHAL.L ERd	1					
SHAR	SHAR.B Rd	1					
	SHAR.W Rd	1					
	SHAR.L ERd	1					
SHLL	SHLL.B Rd	1					
	SHLL.W Rd	1					
	SHLL.L ERd	1					
SHLR	SHLR.B Rd	1					
	SHLR.W Rd	1					
	SHLR.L ERd	1					
SLEEP	SLEEP	1					
STC	STC CCR, Rd	1					
	STC CCR, @ERd	2				1	
	STC CCR, @(d:16,ERd)	3				1	
	STC CCR, @(d:24,ERd)	5				1	
	STC CCR,@-ERd	2				1	2
	STC CCR, @aa:16	3				1	
	STC CCR, @aa:24	4				1	
SUB	SUB.B Rs, Rd	1					
	SUB.W #xx:16, Rd	2					
	SUB.W Rs, Rd	1					
	SUB.L #xx:32, ERd	3					
	SUB.L ERs, ERd	1					
SUBS	SUBS #1/2/4, ERd	1					





Instruction	Mnemonic	Instruction Fetch I	Branch Addr. Read J	Stack Operation K	Byte Data Access L	Word Data Access M	Internal Operation N
SUBX	SUBX #xx:8, Rd	1					
	SUBX. Rs, Rd	1					
XOR	XOR.B #xx:8, Rd	1					
	XOR.B Rs, Rd	1					
	XOR.W #xx:16, Rd	2					
	XOR.W Rs, Rd	1					
	XOR.L #xx:32, ERd	3					
	XOR.L ERs, ERd	2					
XORC	XORC #xx:8, CCR	1					

Notes: 1. n: Specified value in R4L. The source and destination operands are accessed n+1 times respectively.

2. It cannot be used in this LSI.

A.4 Combinations of Instructions and Addressing Modes

Table A.5 Combinations of Instructions and Addressing Modes

		Addressing Mode												
Functions	Instructions	xx#	Rn	@ERn	@(d:16.ERn)	@(d:24.ERn)	@ERn+/@ERn	@aa:8	@aa:16	@aa:24	@(d:8.PC)	@(d:16.PC)	@ @ aa:8	1
Data	MOV	BWL	BWL	BWL	BWL	BWL	BWL	В	BWL	BWL	_	_	_	
transfer	POP, PUSH	_	_	_	_	_	_	_	_	_	_	_	_	WL
instructions	MOVFPE, MOVTPE	_	_	_	_	_	_	_	_	_	_	_	_	_
Arithmetic	ADD, CMP	BWL	BWL	_	_	_	_	_	_	_	_	_	_	_
operations	SUB	WL	BWL	_	_	_	_	_	_	_	_	_	_	_
	ADDX, SUBX	В	В	_	_	_	_	_	_	_	_	_	_	
	ADDS, SUBS	_	L	_	_	_	_	_	_	_	_	—	_	-
	INC, DEC	_	BWL	_	_	—	_	_	_	_	_	—	_	-
	DAA, DAS	_	В	_	_	—	_	_	_	_	_	—	_	-
	MULXU,	_	BW	_	_	_	_	_	_	_	_	_	_	
	MULXS,													
	DIVXU,													
	DIVXS													
	NEG		BWL					_	_	_				
	EXTU, EXTS	_	WL					_	_	_				
Logical	AND, OR, XOR	_	BWL					_	_					
operations	NOT	_	BWL	_			_	_	_	_				
Shift operation		_	BWL					_	_		_			
Bit manipulat		_	В	В			_	В	_	_	_	_		
Branching	BCC, BSR	_						_	_					
instructions	JMP, JSR	_		0			_	_	_	_	0	0		
	RTS	_	_	_	_	_	_	_	_	0	_	_	0	
System	RTE	_	_	_	_	_	_	_	_	_	_	_	_	0
control	SLEEP	_	_	_	_	_	_	_	_	_	_	_	_	0
instructions	LDC	В	В	W	W	W	W	_	W	W	_			
	STC	_	В	W	W	W	W	_	W	W	_		_	_
	ANDC, ORC, XORC	В	_	_	_	_	_	_	_	_	_	-	—	-
	NOP	_	_	_	_	_	_	_	_	_	_	_	_	
Block data tra	ansfer instructions	_	_	_	_	_	_	_	_	_	_	_	_	BW

B. I/O Port Block Diagrams

B.1 Block Diagrams of Port 1

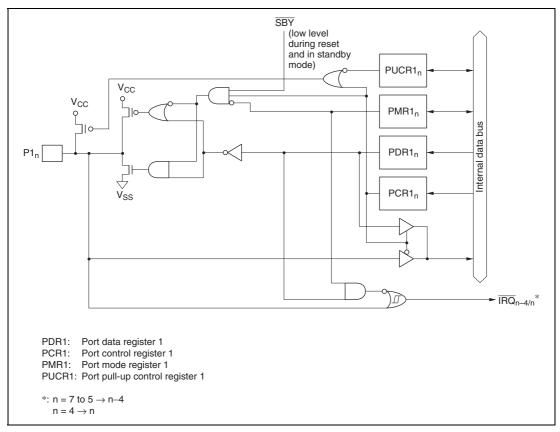


Figure B.1 (a) Port 1 Block Diagram (Pins P1, to P1,

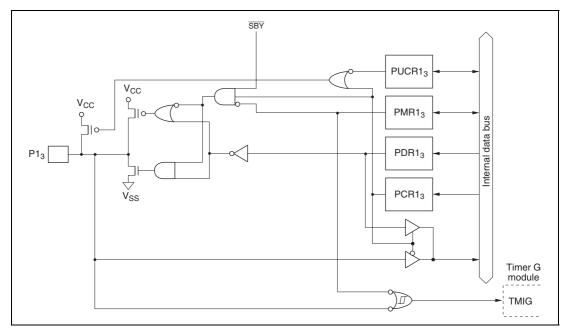
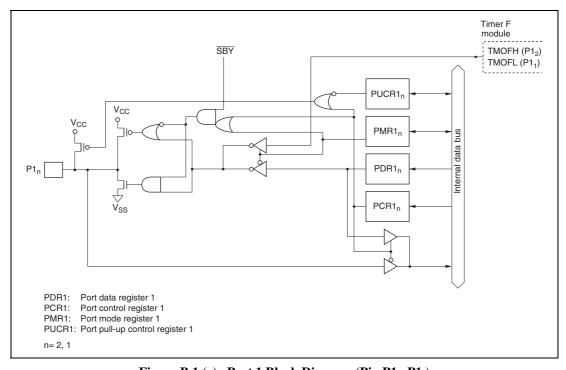


Figure B.1 (b) Port 1 Block Diagram (Pin P1₃)



 $Figure\ B.1\ (c)\quad Port\ 1\ Block\ Diagram\ (Pin\ P1_2,\ P1_1)$

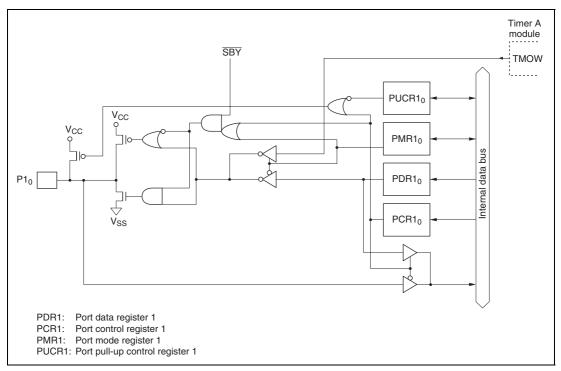


Figure B.1 (d) Port 1 Block Diagram (Pin P1₀)

B.2 Block Diagrams of Port 3

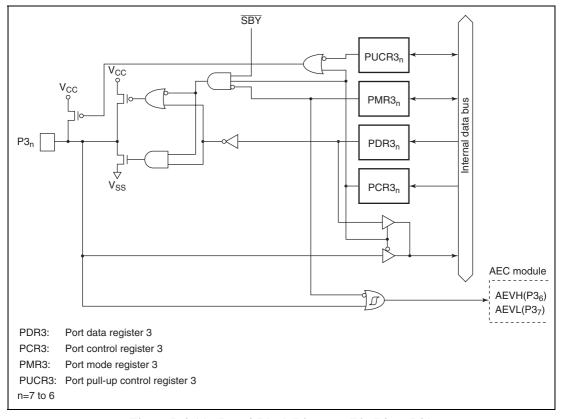


Figure B.2 (a) Port 3 Block Diagram (Pin $P3_7$ to $P3_6$)

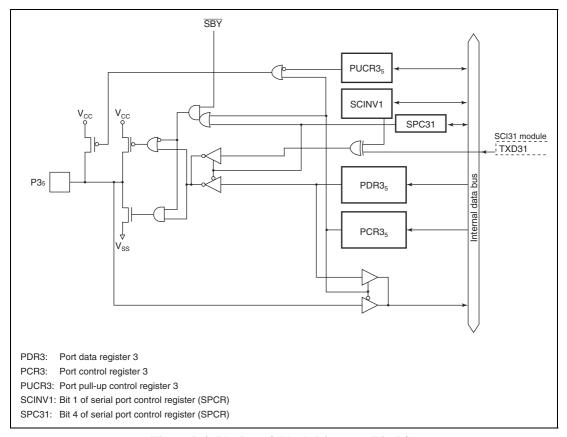


Figure B.2 (b) Port 3 Block Diagram (Pin $P3_5$)

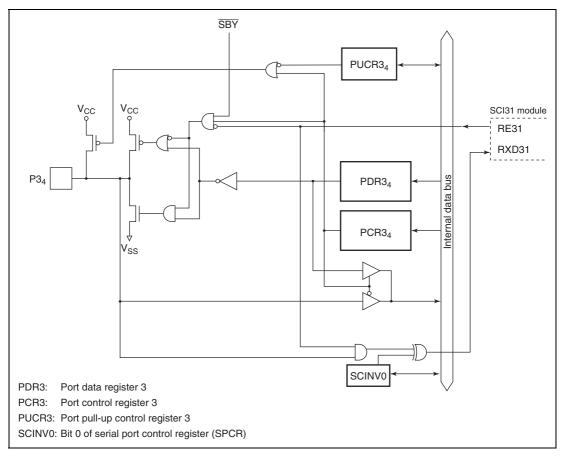


Figure B.2 (c) Port 3 Block Diagram (Pin P3₄)

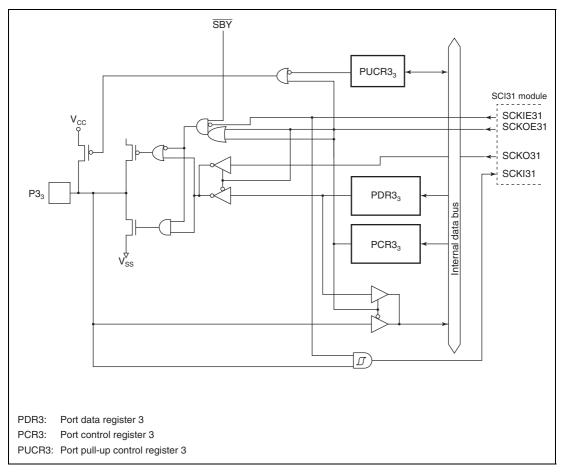


Figure B.2 (d) Port 3 Block Diagram (Pin P3₃)

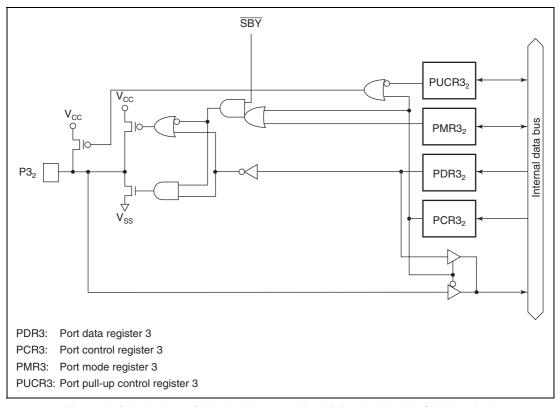


Figure B.2 (e-1) Port 3 Block Diagram (Pin P3₂ in the Mask ROM Version)

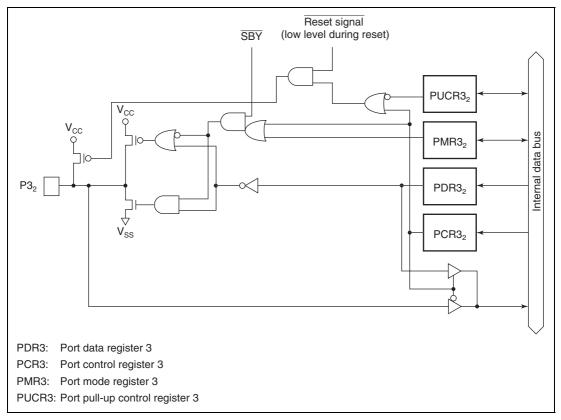


Figure B.2 (e-2) Port 3 Block Diagram (Pin P3, in the Flash Memory Version)

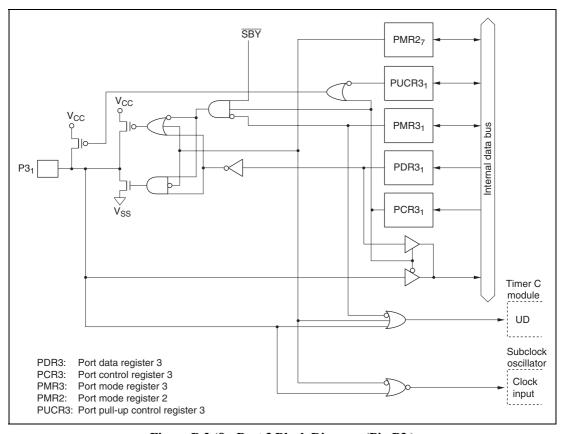


Figure B.2 (f) Port 3 Block Diagram (Pin P3,)

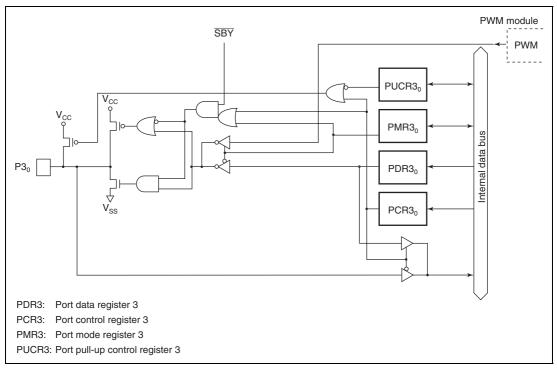


Figure B.2 (g) Port 3 Block Diagram (Pin P3₀)

B.3 Block Diagrams of Port 4

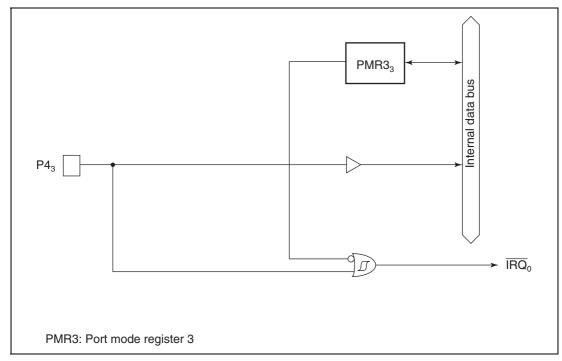
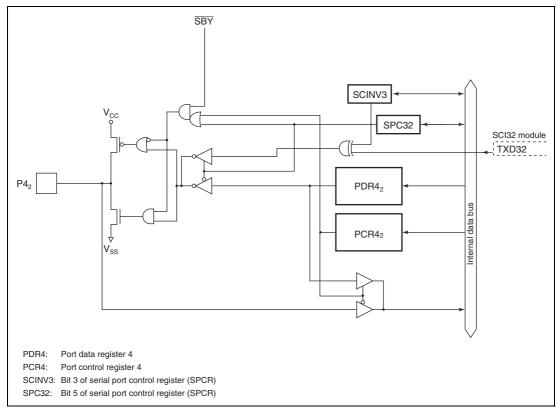


Figure B.3 (a) Port 4 Block Diagram (Pin P4₃)



 $Figure\ B.3\ (b)\quad Port\ 4\ Block\ Diagram\ (Pin\ P4_2)$

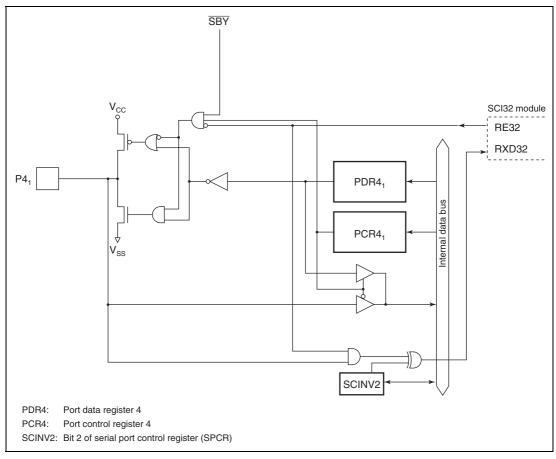


Figure B.3 (c) Port 4 Block Diagram (Pin P4₁)

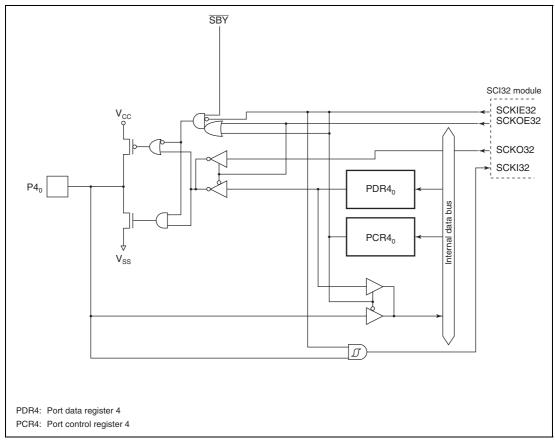


Figure B.3 (d) Port 4 Block Diagram (Pin P4₀)

B.4 Block Diagram of Port 5

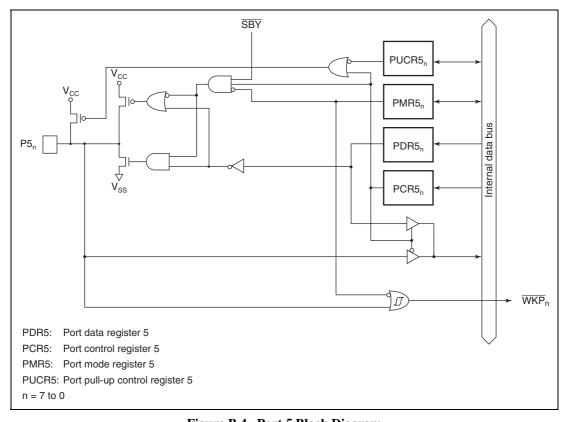


Figure B.4 Port 5 Block Diagram

B.5 Block Diagram of Port 6

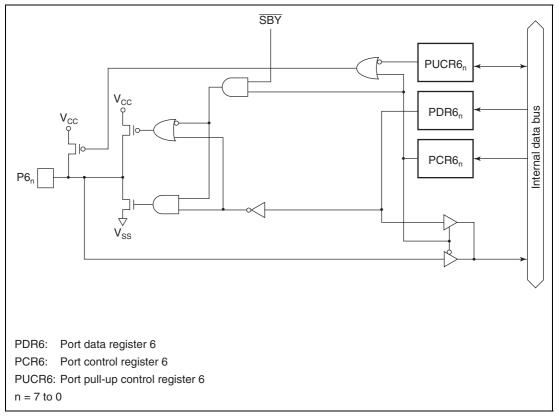


Figure B.5 Port 6 Block Diagram



B.6 Block Diagram of Port 7

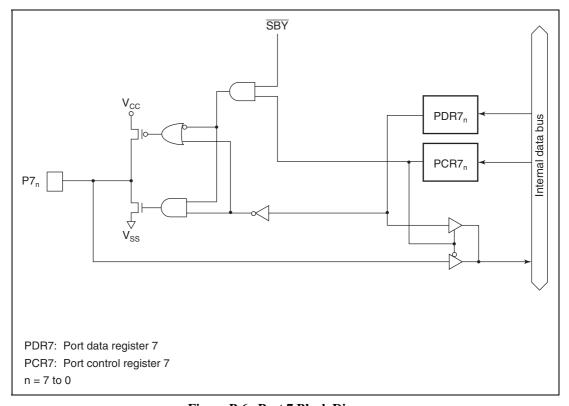


Figure B.6 Port 7 Block Diagram

B.7 Block Diagrams of Port 8

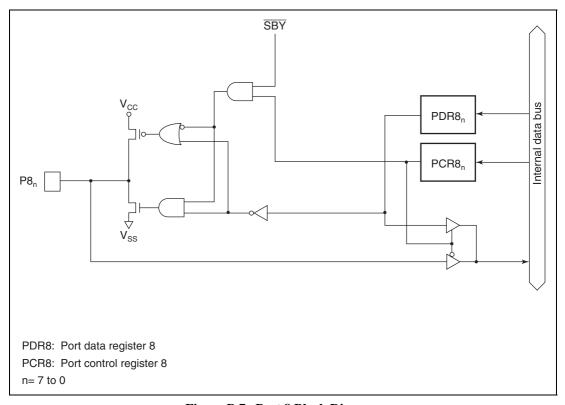


Figure B.7 Port 8 Block Diagram

B.8 Block Diagram of Port A

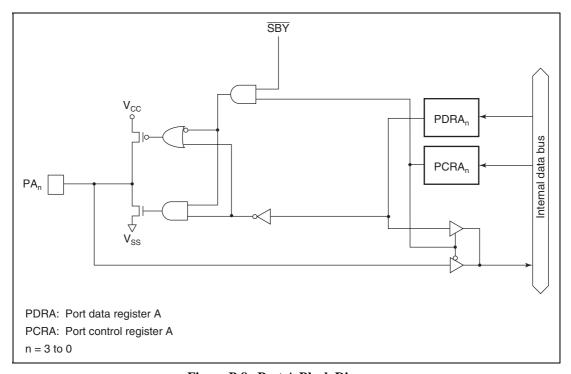


Figure B.8 Port A Block Diagram

B.9 Block Diagram of Port B

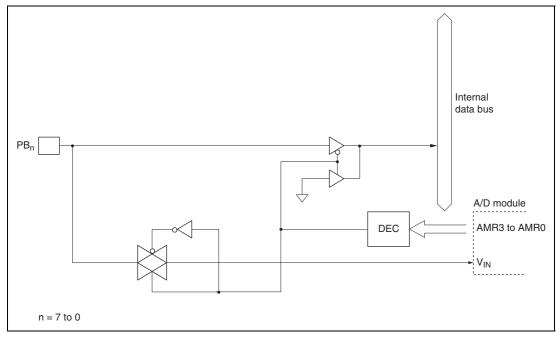


Figure B.9 Port B Block Diagram

C. Port States in the Different Processing States

Table C.1 Port States Overview

Port	Reset	Sleep	Subsleep	Standby	Watch	Subactive	Active
P1, to P1 ₀	High impedance	Retained	Retained	High impedance*1	Retained	Functions	Functions
P3, to P3 ₀	High impedance*2	Retained	Retained	High impedance*1	Retained	Functions	Functions
P4 ₃ to	High impedance	Retained	Retained	High impedance	Retained	Functions	Functions
P5 ₇ to P5 ₀	High impedance	Retained	Retained	High impedance*1	Retained	Functions	Functions
P6 ₇ to P6 ₀	High impedance	Retained	Retained	High impedance	Retained	Functions	Functions
P7, to P7 ₀	High impedance	Retained	Retained	High impedance	Retained	Functions	Functions
P8, to P80	High impedance	Retained	Retained	High impedance	Retained	Functions	Functions
PA ₃ to PA ₀	High impedance	Retained	Retained	High impedance	Retained	Functions	Functions
PB ₇ to PB ₀	High impedance	High impedance	High impedance	High impedance	High impedance	High impedance	High impedance

Notes: 1. High level output when MOS pull-up is in on state.

^{2.} On-chip pull-up MOS turns on for pin P32 only (Flash Memory Version).

D. List of Product Codes

Table D.1 Product Code Lineup

Product T	ype			Product Code	Mark Code	Package (Package Code)
H8/38537	H8/38532	Mask ROM	Regular	HD64338532H	38532H	80-pin QFP (FP-80A)
Group		versions	products	HD64338532W	38532W	80-pin TQFP (TFP-80C)
			Wide-range	HD64338532HW	38532H	80-pin QFP (FP-80A)
			specification products	HD64338532WW	38532W	80-pin TQFP (TFP-80C)
	H8/38533	Mask ROM	Regular	HD64338533H	38533H	80-pin QFP (FP-80A)
		versions	products	HD64338533W	38533W	80-pin TQFP (TFP-80C)
			Wide-range specification products	HD64338533HW	38533H	80-pin QFP (FP-80A)
				HD64338533WW	38533W	80-pin TQFP (TFP-80C)
	H8/38534		Regular products	HD64338534H	38534H	80-pin QFP (FP-80A)
				HD64338534W	38534W	80-pin TQFP (TFP-80C)
			Wide-range	HD64338534HW	38534H	80-pin QFP (FP-80A)
			specification products	HD64338534WW	38534W	80-pin TQFP (TFP-80C)
		Flash	Regular	HD64F38534H	F38534H	80-pin QFP (FP-80A)
		memory versions	products	HD64F38534W	F38534W	80-pin TQFP (TFP-80C)
			Wide-range	HD64F38534HW	F38534H	80-pin QFP (FP-80A)
	•		specification products	HD64F38534WW	F38534W	80-pin TQFP (TFP-80C)



Product T	ype			Product Code	Mark Code	Package (Package Code)
H8/38537	H8/38535	Mask ROM	Regular	HD64338535H	38535H	80-pin QFP (FP-80A)
Group		versions	products	HD64338535W	38535W	80-pin TQFP (TFP-80C)
			Wide-range	HD64338535HW	38535H	80-pin QFP (FP-80A)
			specification products	HD64338535WW	38535W	80-pin TQFP (TFP-80C)
	H8/38536	Mask ROM	Regular	HD64338536H	38536H	80-pin QFP (FP-80A)
	V	versions	products	HD64338536W	38536W	80-pin TQFP (TFP-80C)
			Wide-range specification products	HD64338536HW	38536H	80-pin QFP (FP-80A)
				HD64338536WW	38536W	80-pin TQFP (TFP-80C)
	H8/38537	Mask ROM	Regular	HD64338537H	38537H	80-pin QFP (FP-80A)
		versions	products	HD64338537W	38537W	80-pin TQFP (TFP-80C)
			Wide-range	HD64338537HW	38537H	80-pin QFP (FP-80A)
			specification products	HD64338537WW	38537W	80-pin TQFP (TFP-80C)
		Flash	Regular	HD64F38537H	F38537H	80-pin QFP (FP-80A)
		memory versions	products	HD64F38537W	F38537W	80-pin TQFP (TFP-80C)
			Wide-range	HD64F38537HW	F38537H	80-pin QFP (FP-80A)
			specification products	HD64F38537WW	F38537W	80-pin TQFP (TFP-80C)

E. Package Dimensions

Dimensional drawings of the packages FP-80A and TFP-80C are shown in figures E.1 and E.2, below.

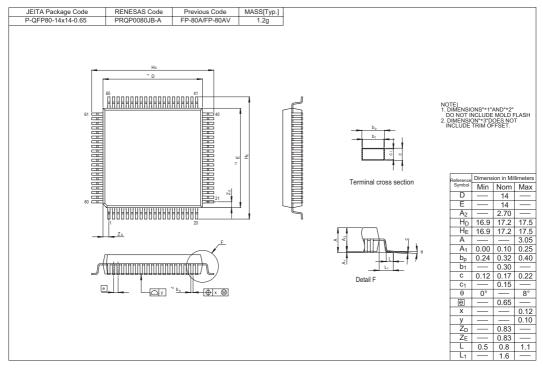


Figure E.1 FP-80A Package Dimensions

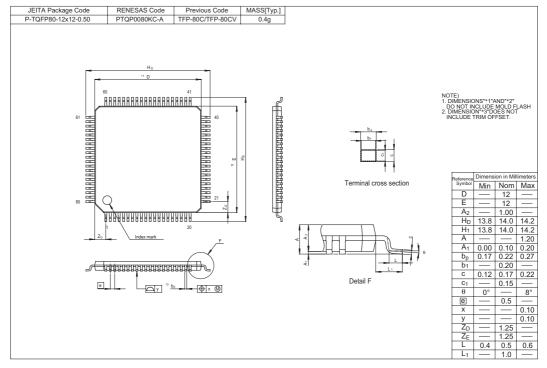


Figure E.2 TFP-80C Package Dimensions



Index

A	IRR1	68
ADRRH	IRR2	69
ADRRL352	IWPR	72
ADSR		
AMR352		
	\mathbf{L}	
	LCR	370
В	LCR2	
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C	0	
CKSTPR1212, 220, 234, 251, 308, 355	OCRF	228
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C W O O C W O O C W O C	O CITA E	220
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H8/38537 Group Hardware Manual



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